

LISP-Click

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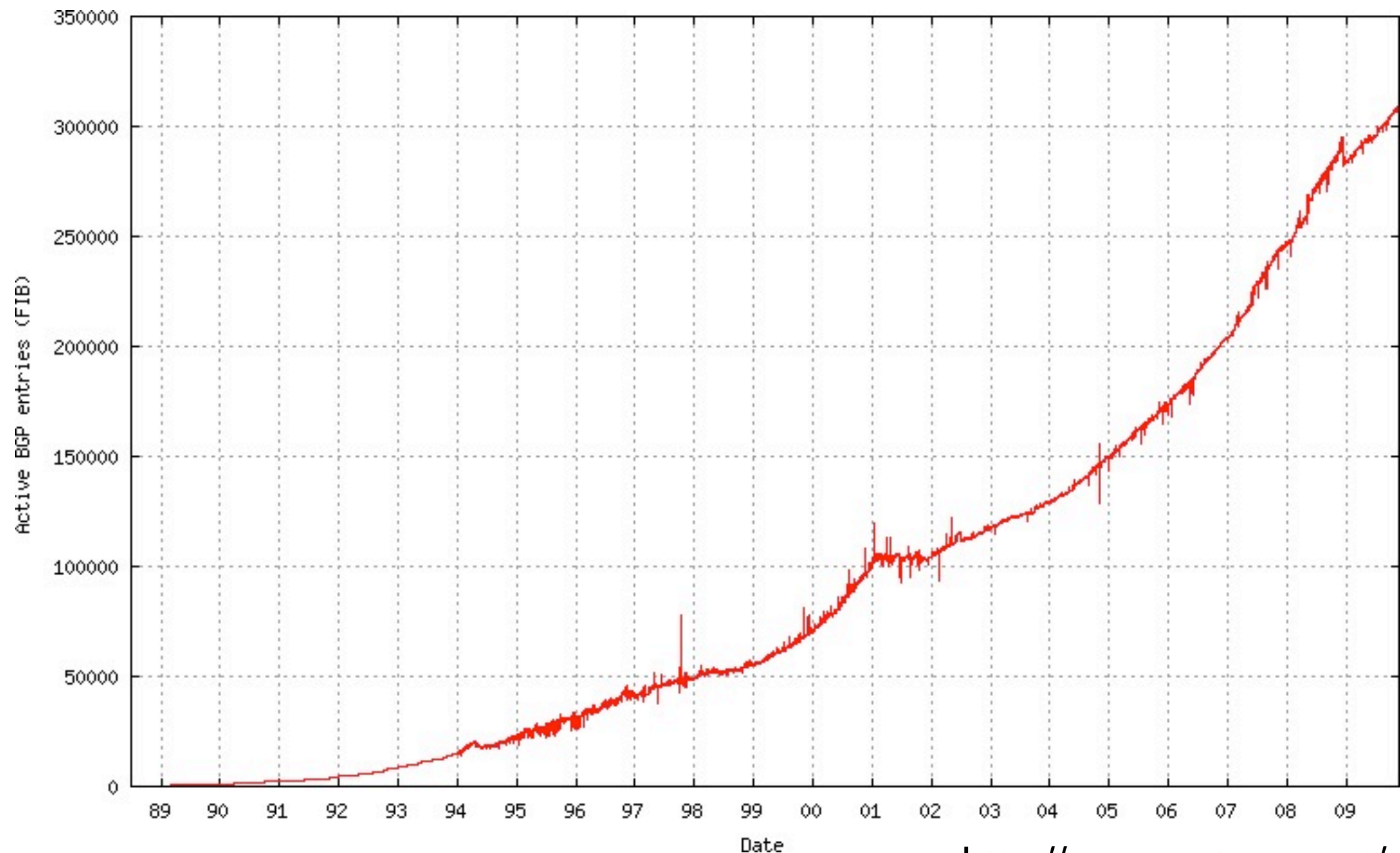
<http://inl.info.ucl.ac.be>

A Click implementation of the Locator/ID Separation Protocol

The Internet is Broken!

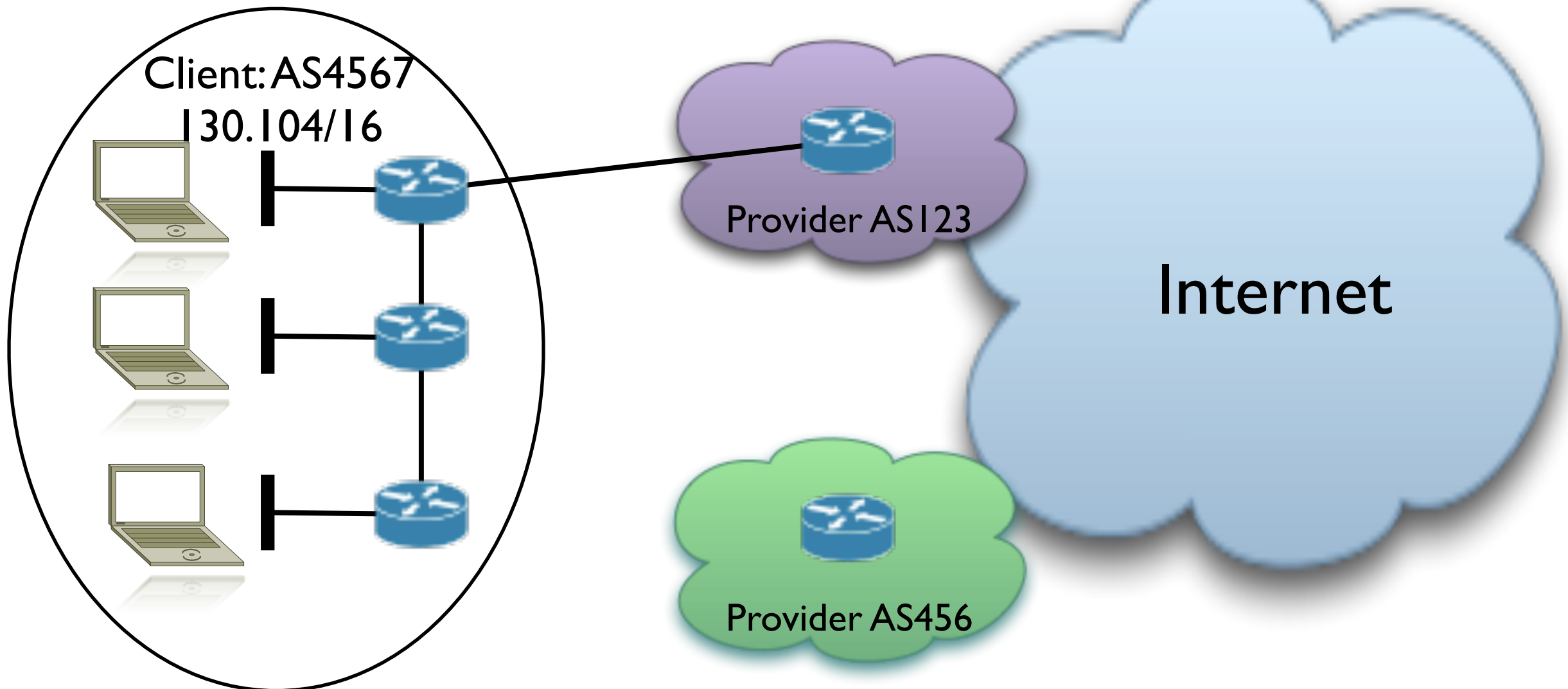
The Internet is broken!

Growth of BGP routing table



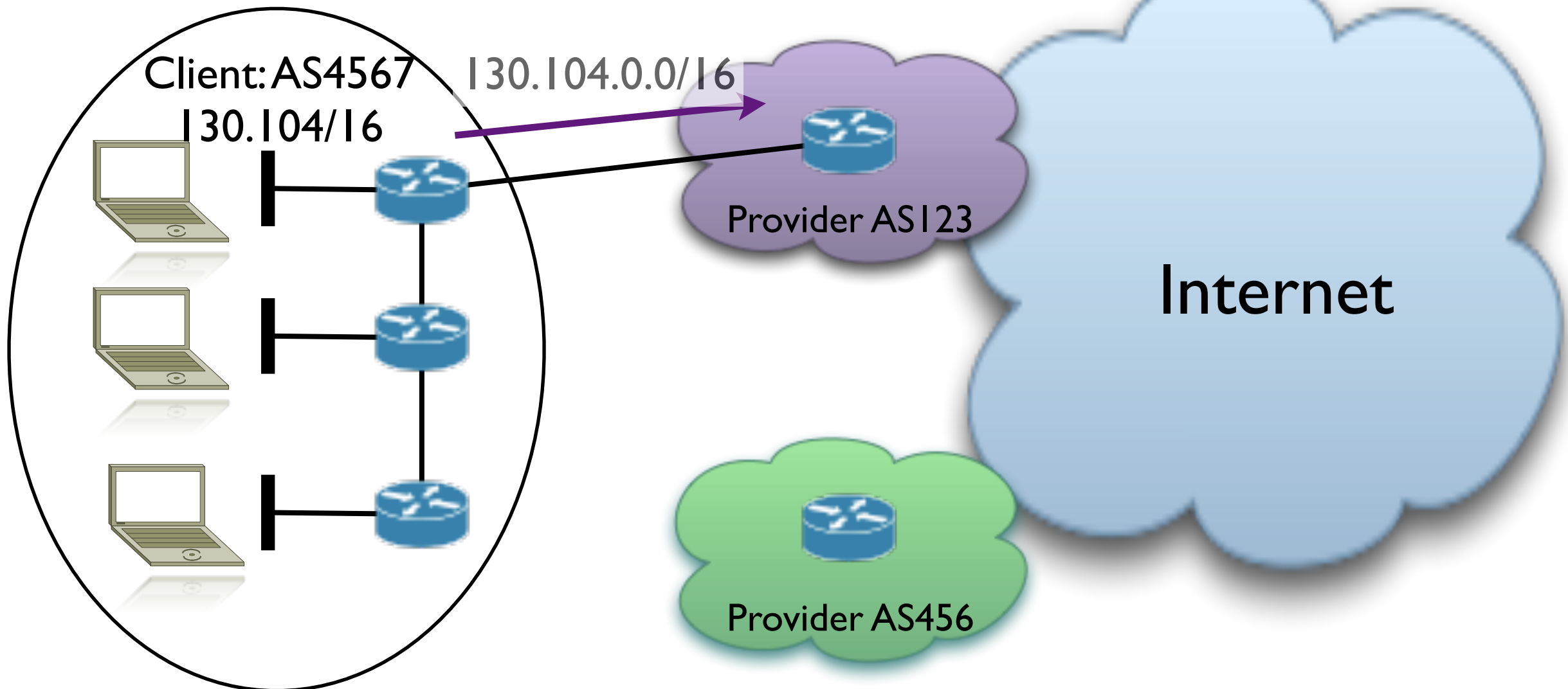
Reasons of the BGP growth?

- More and more networks are internally fragmented. Internet is cheaper than internal links



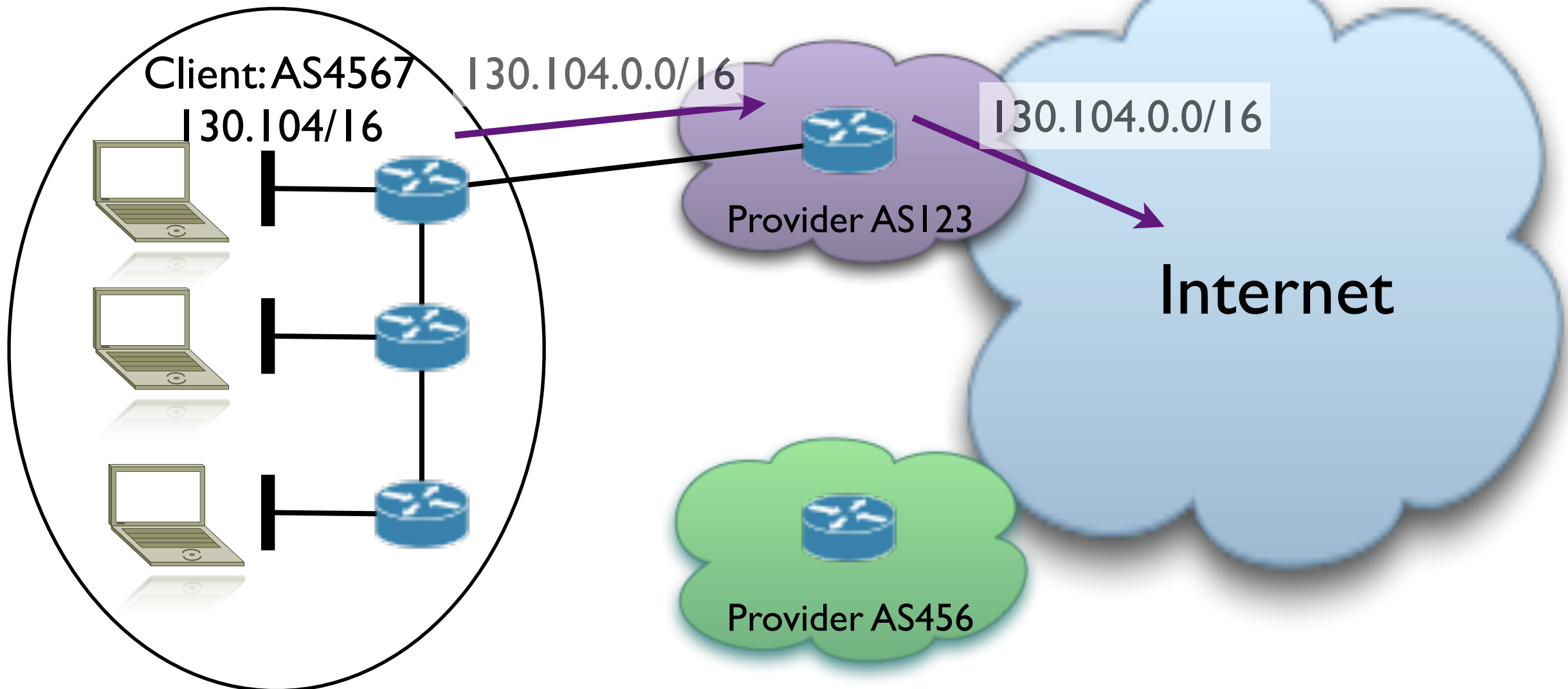
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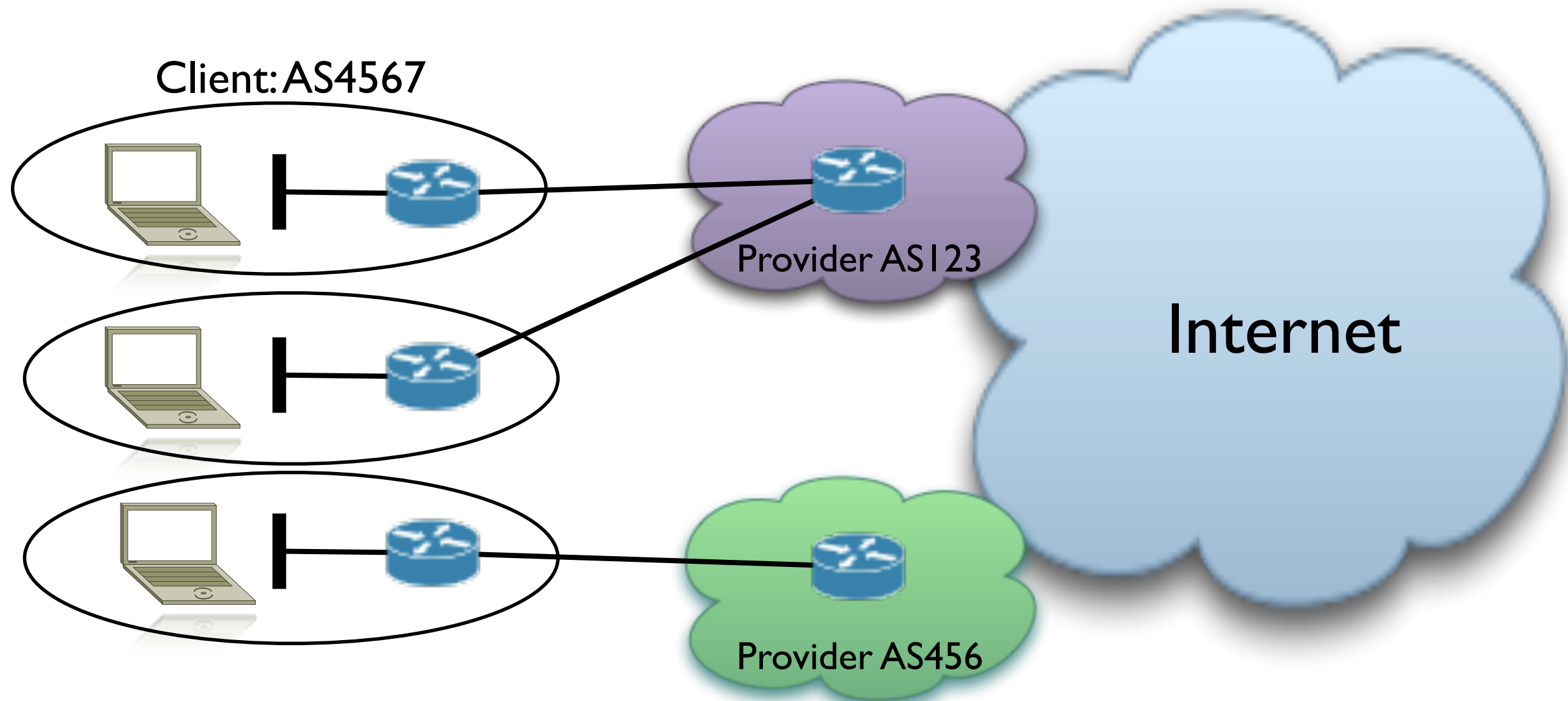
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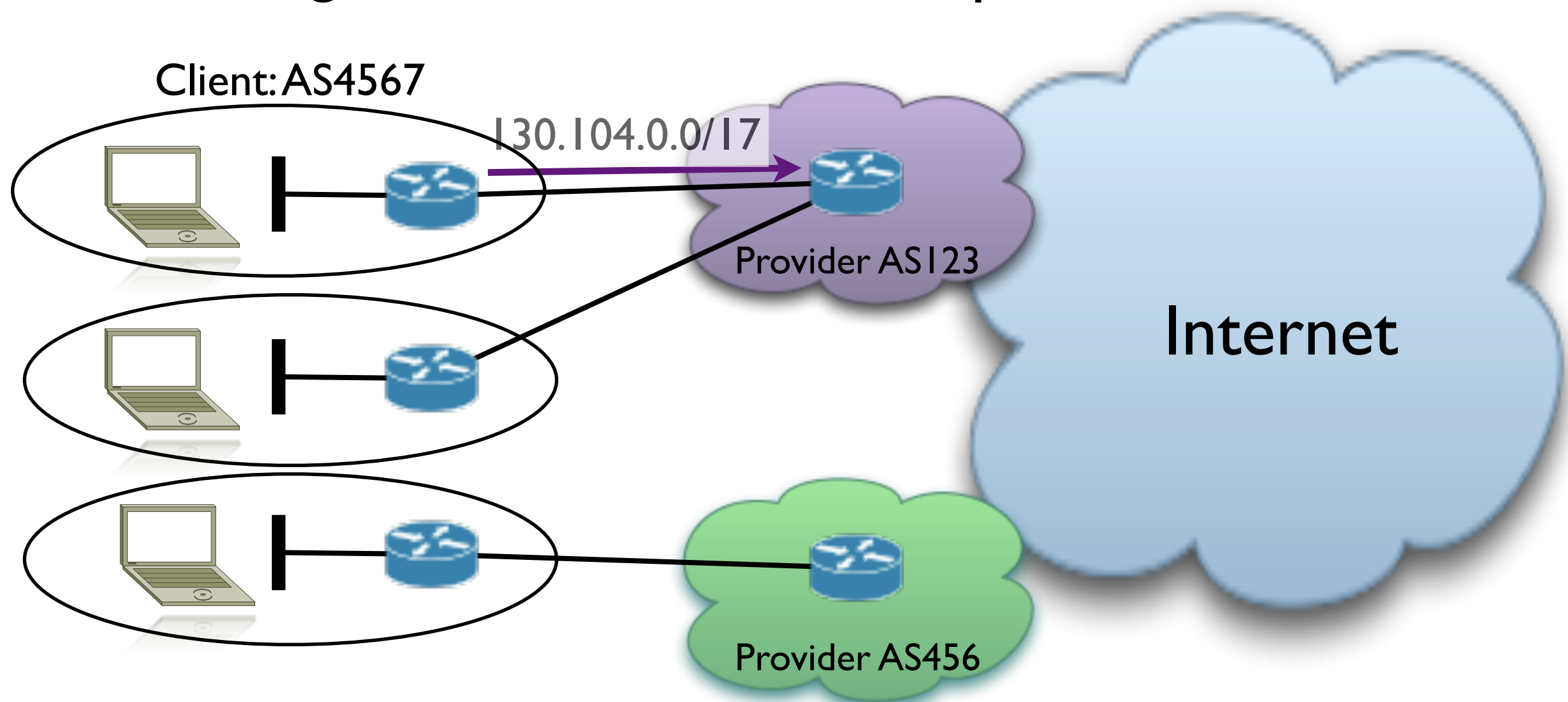
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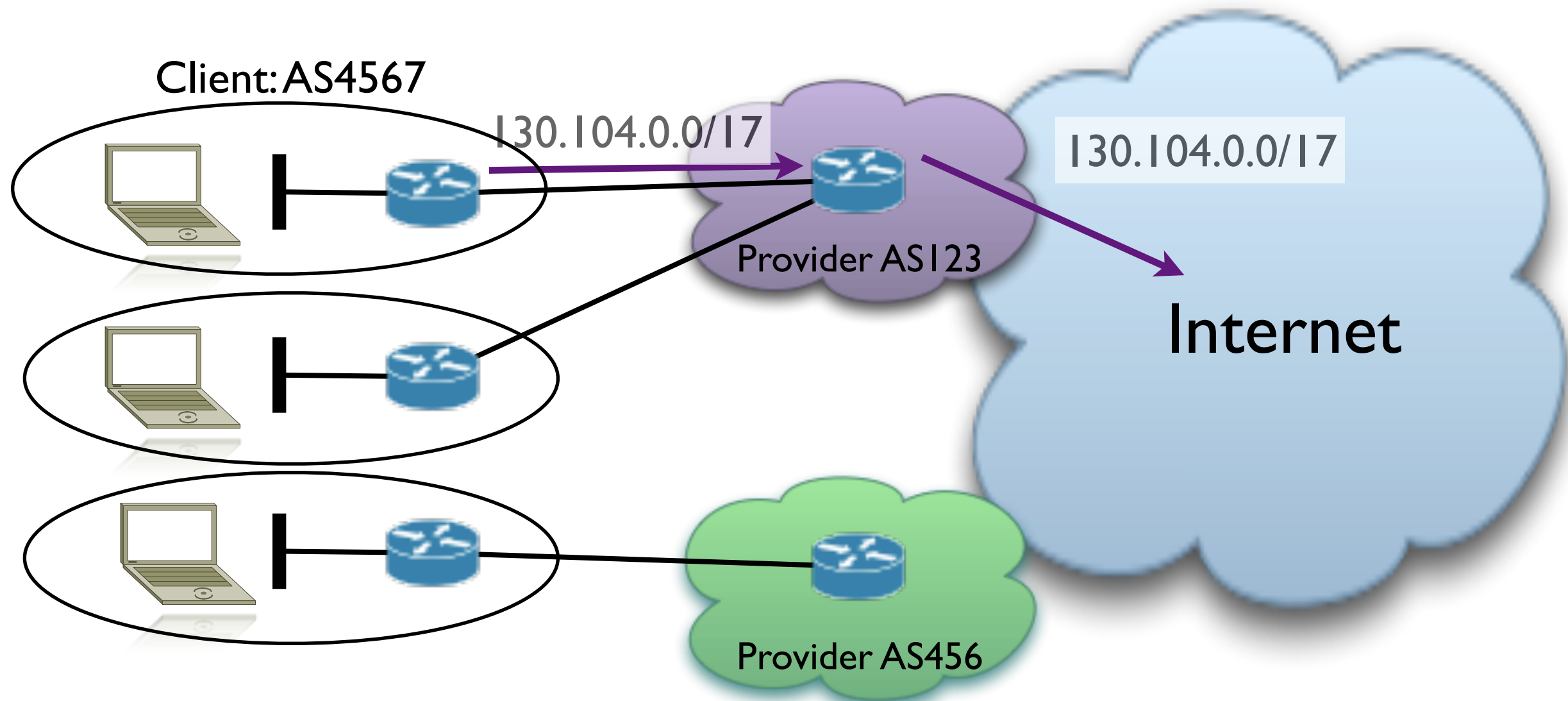
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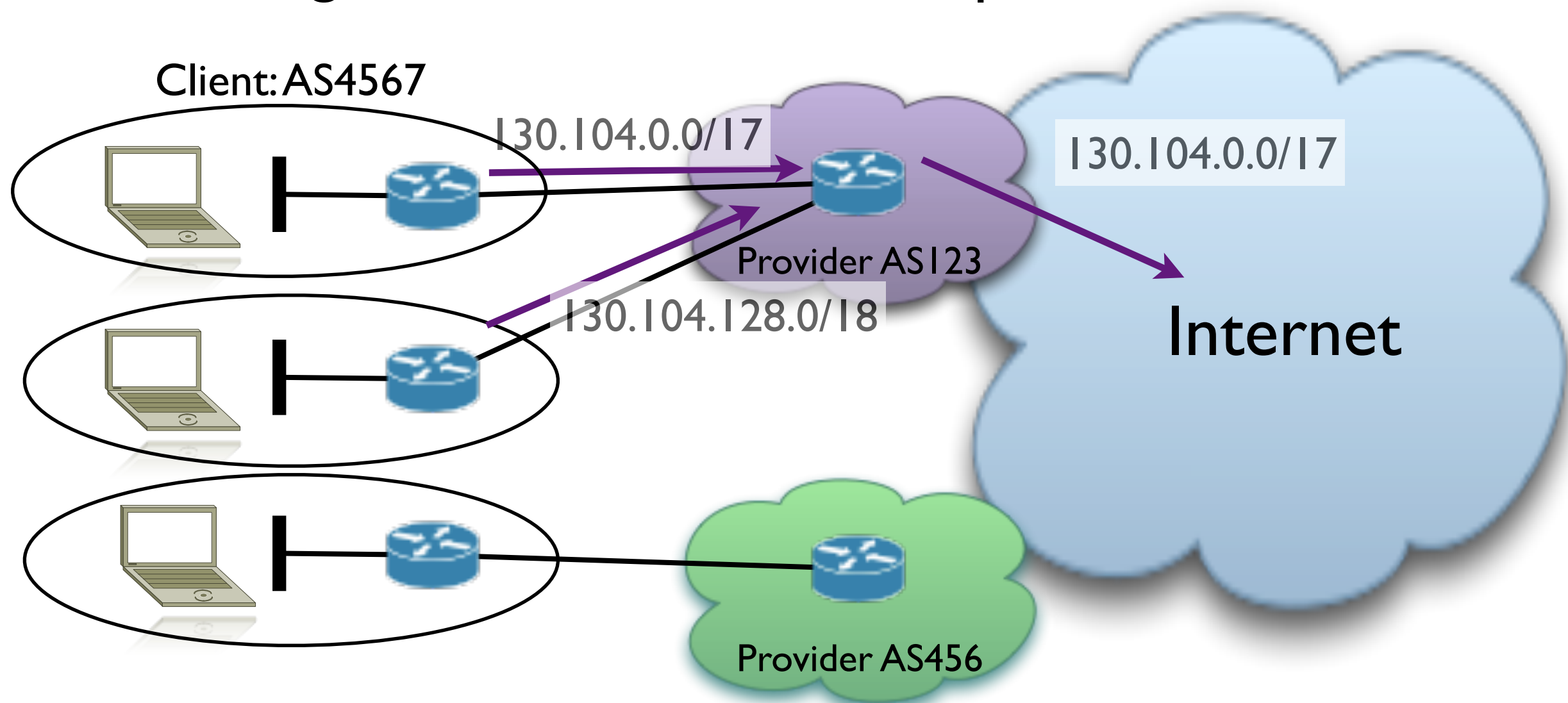
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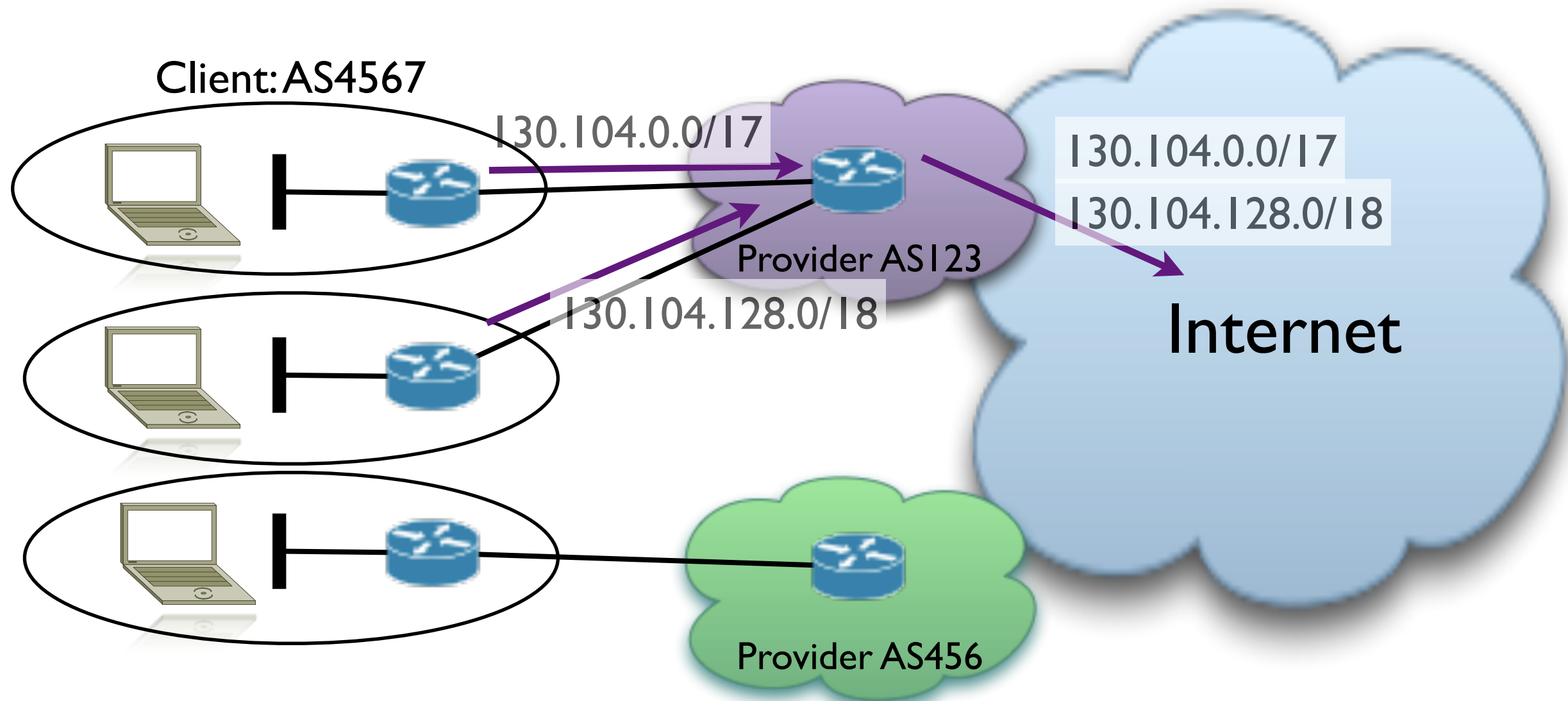
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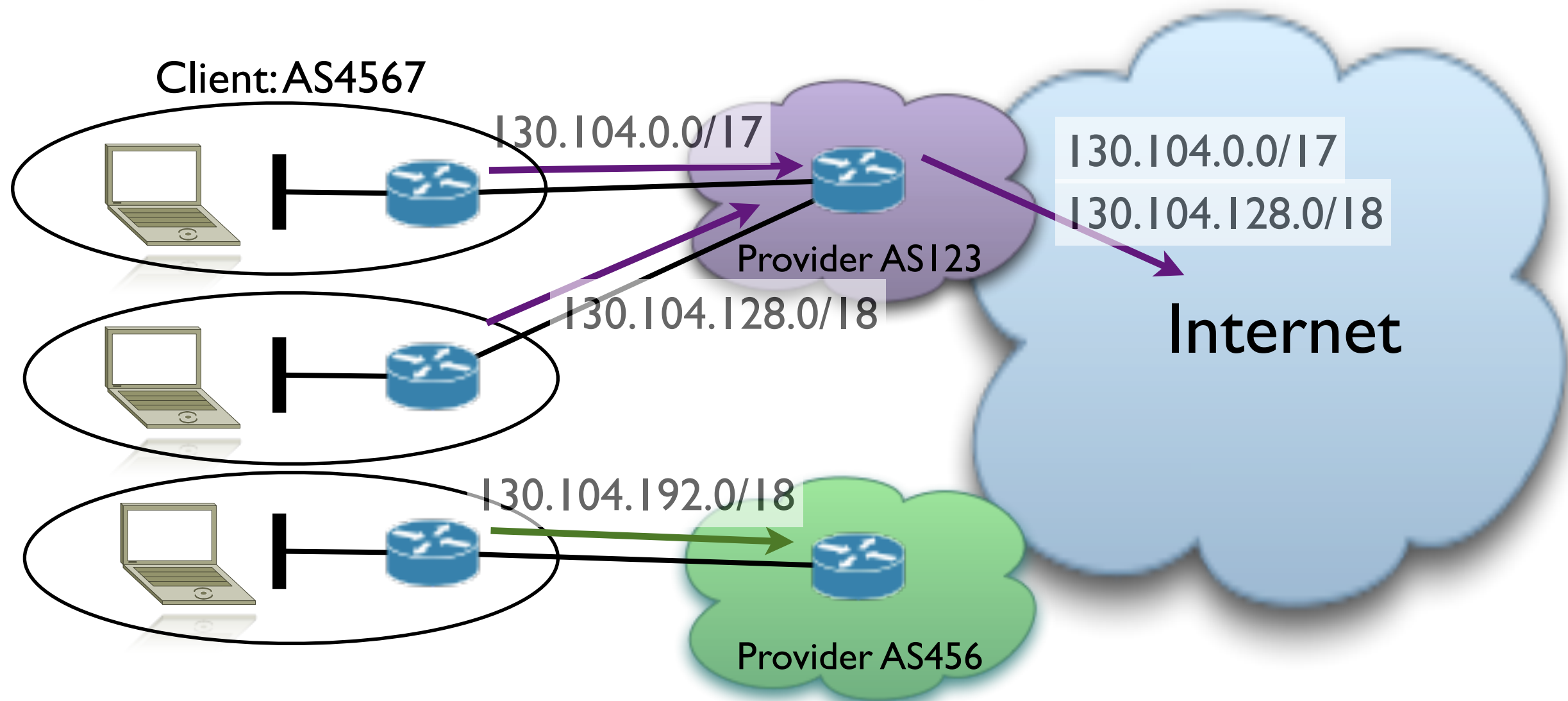
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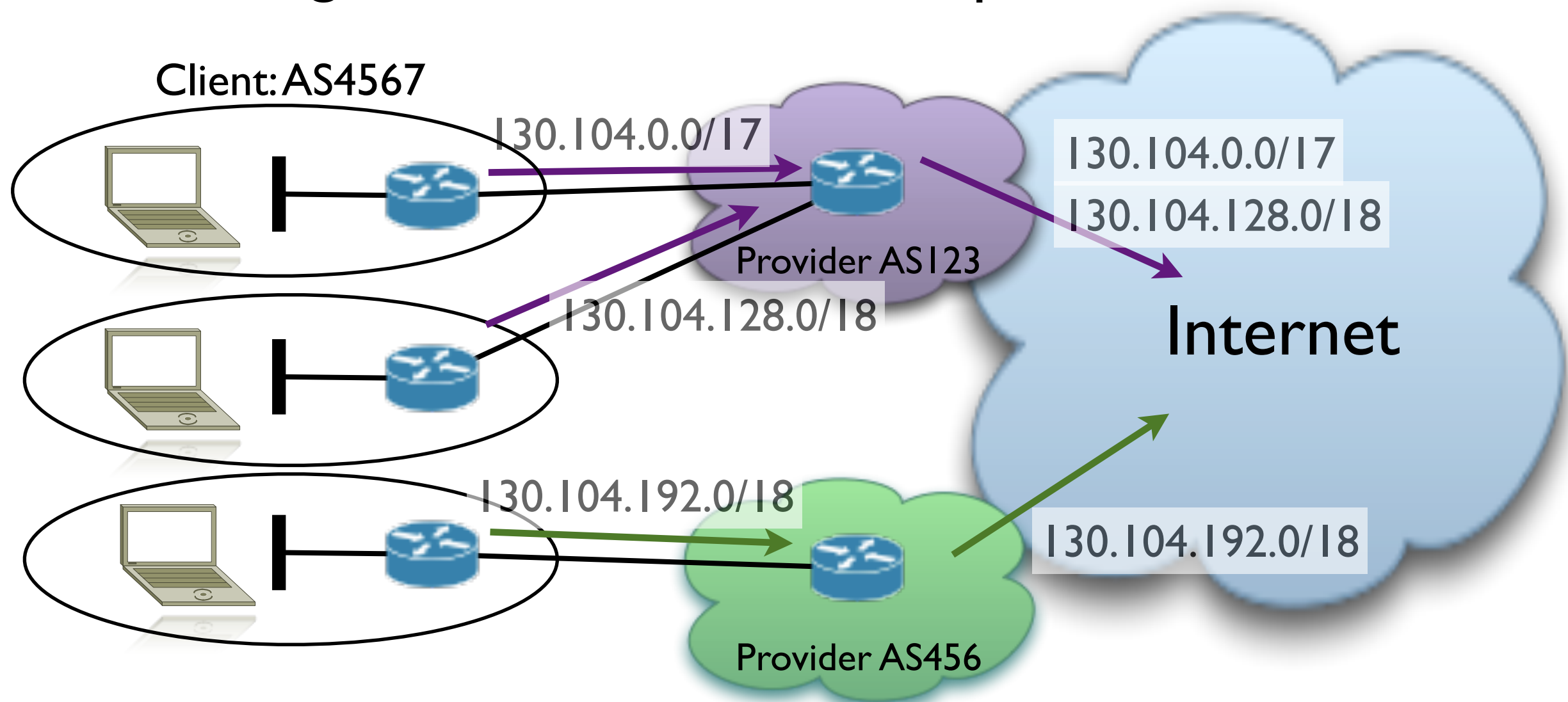
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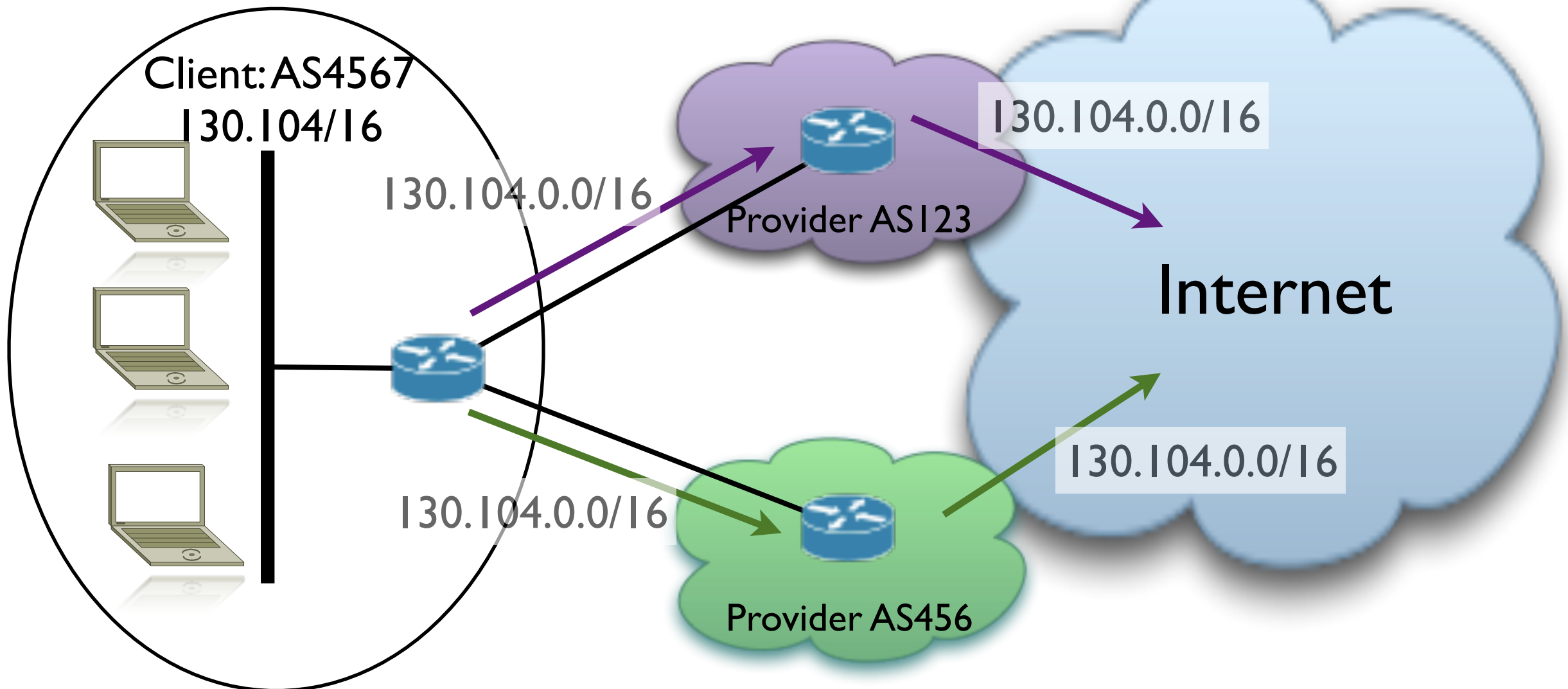
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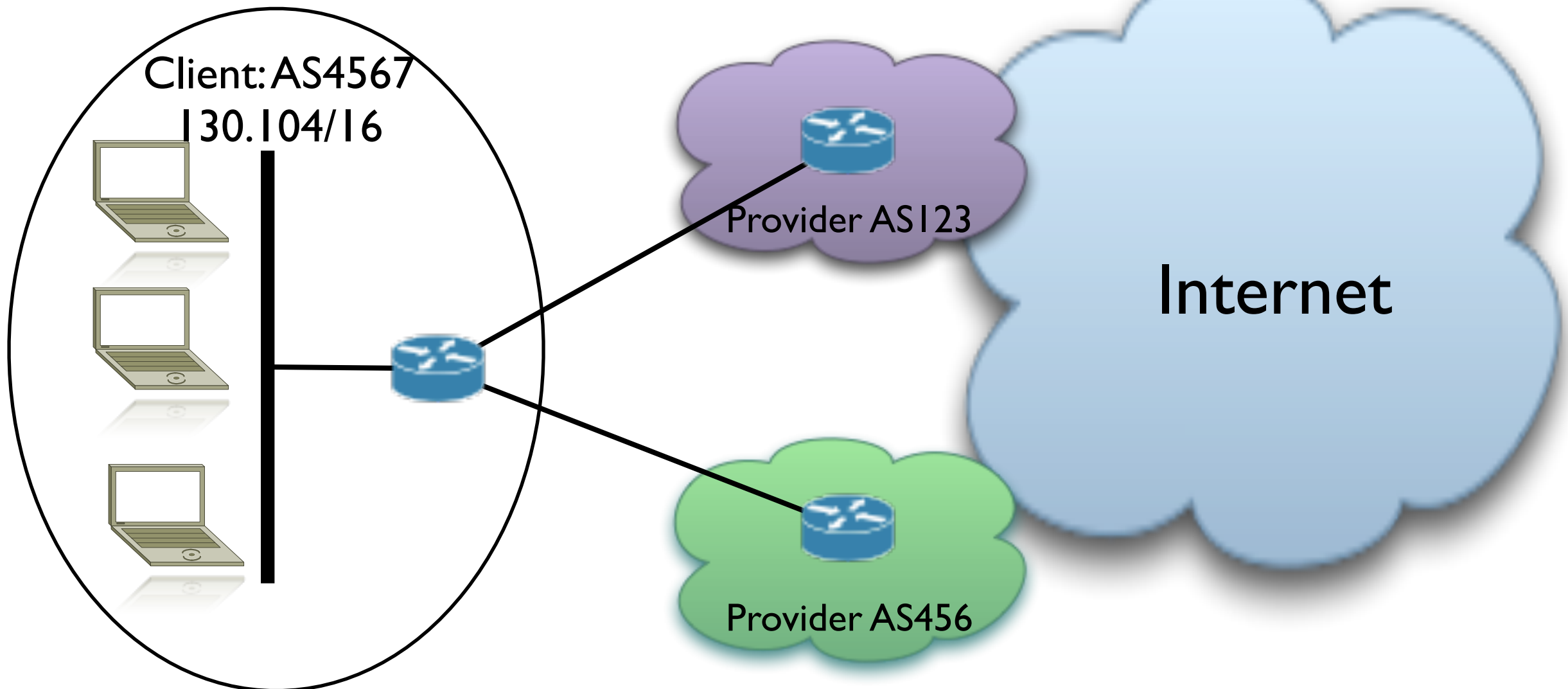
Reasons of the BGP growth?

- Multihoming



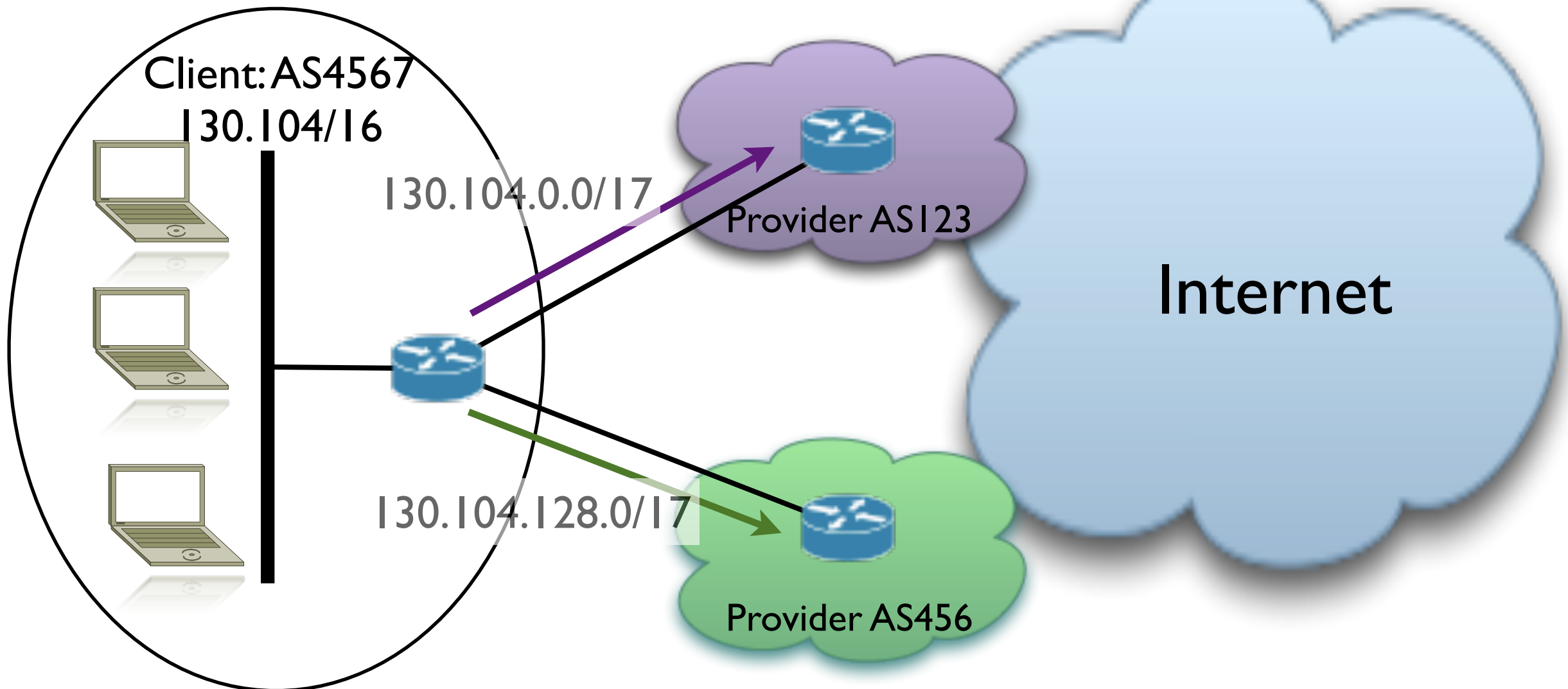
Reasons of the BGP growth?

- Traffic Engineering



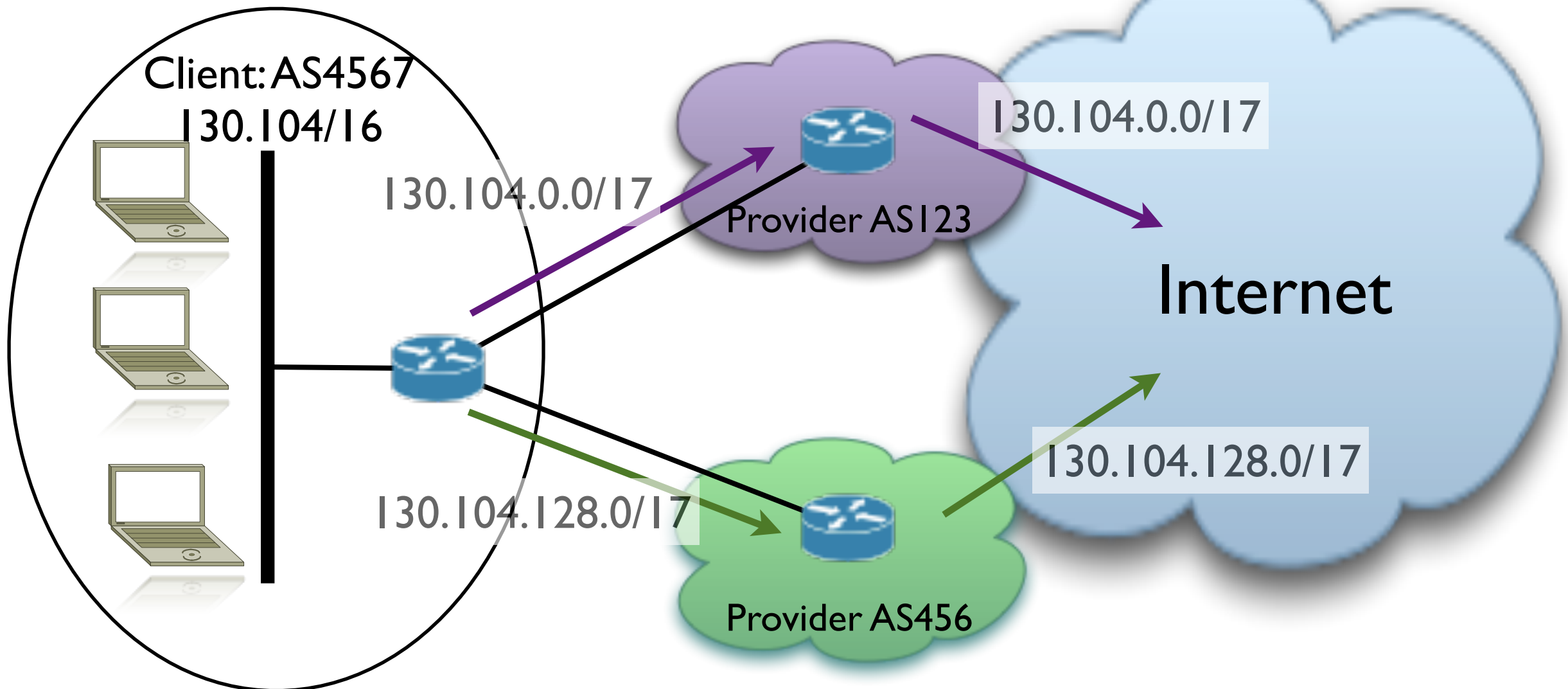
Reasons of the BGP growth?

- Traffic Engineering



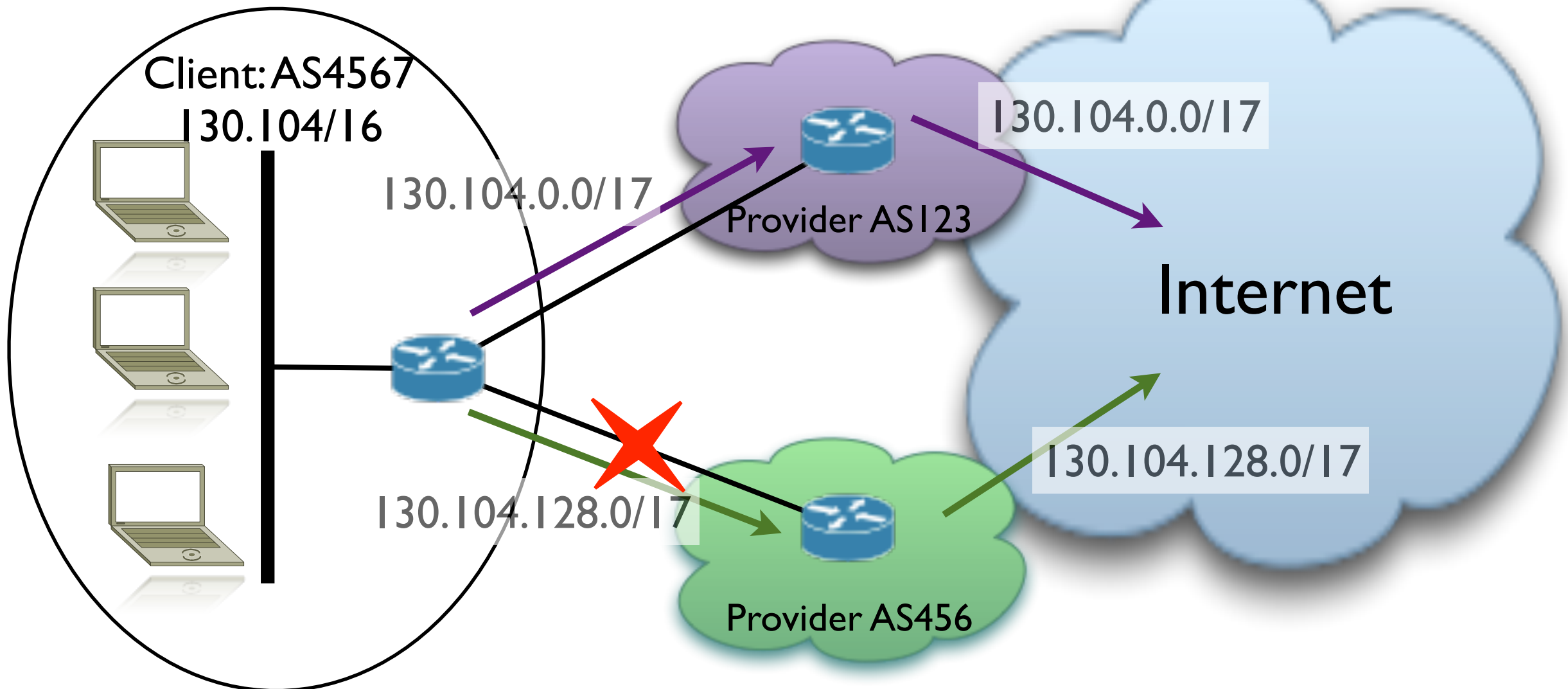
Reasons of the BGP growth?

- Traffic Engineering



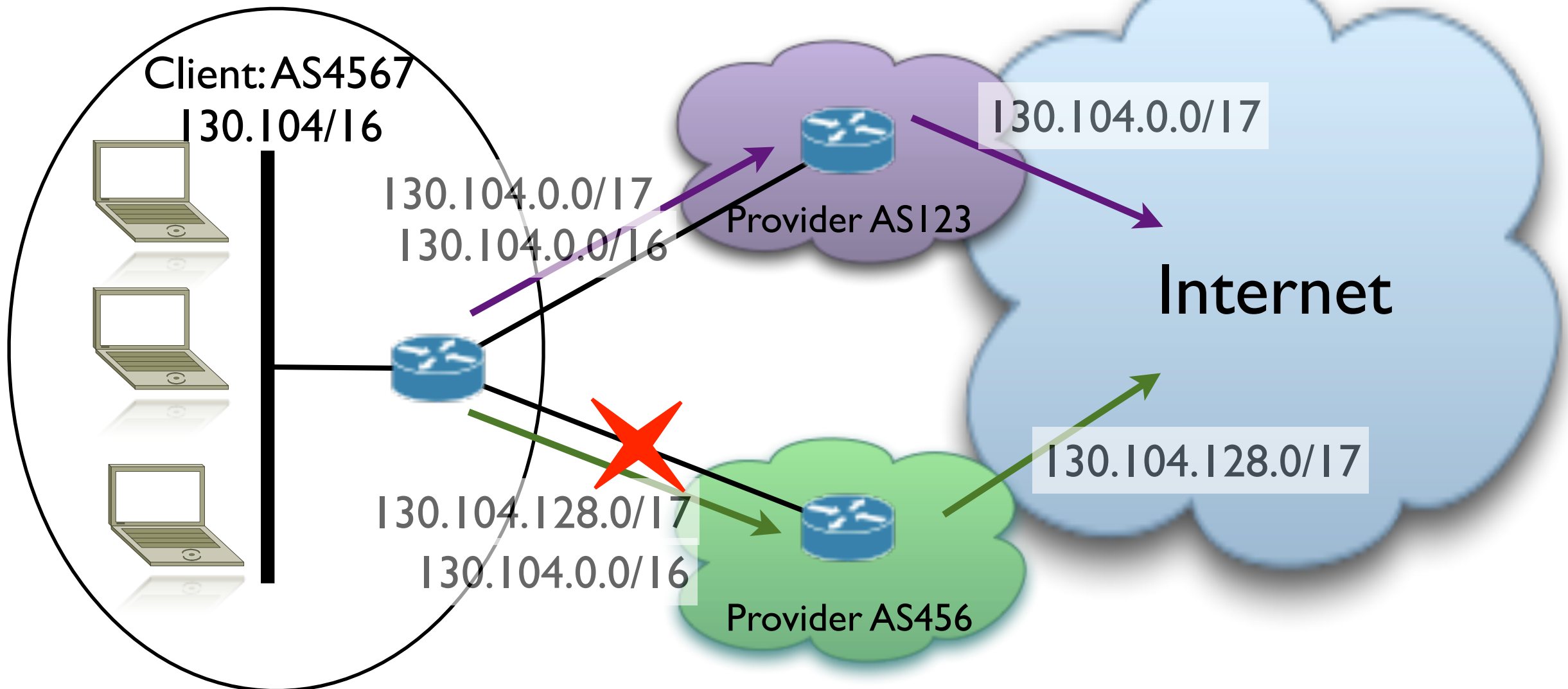
Reasons of the BGP growth?

- Traffic Engineering



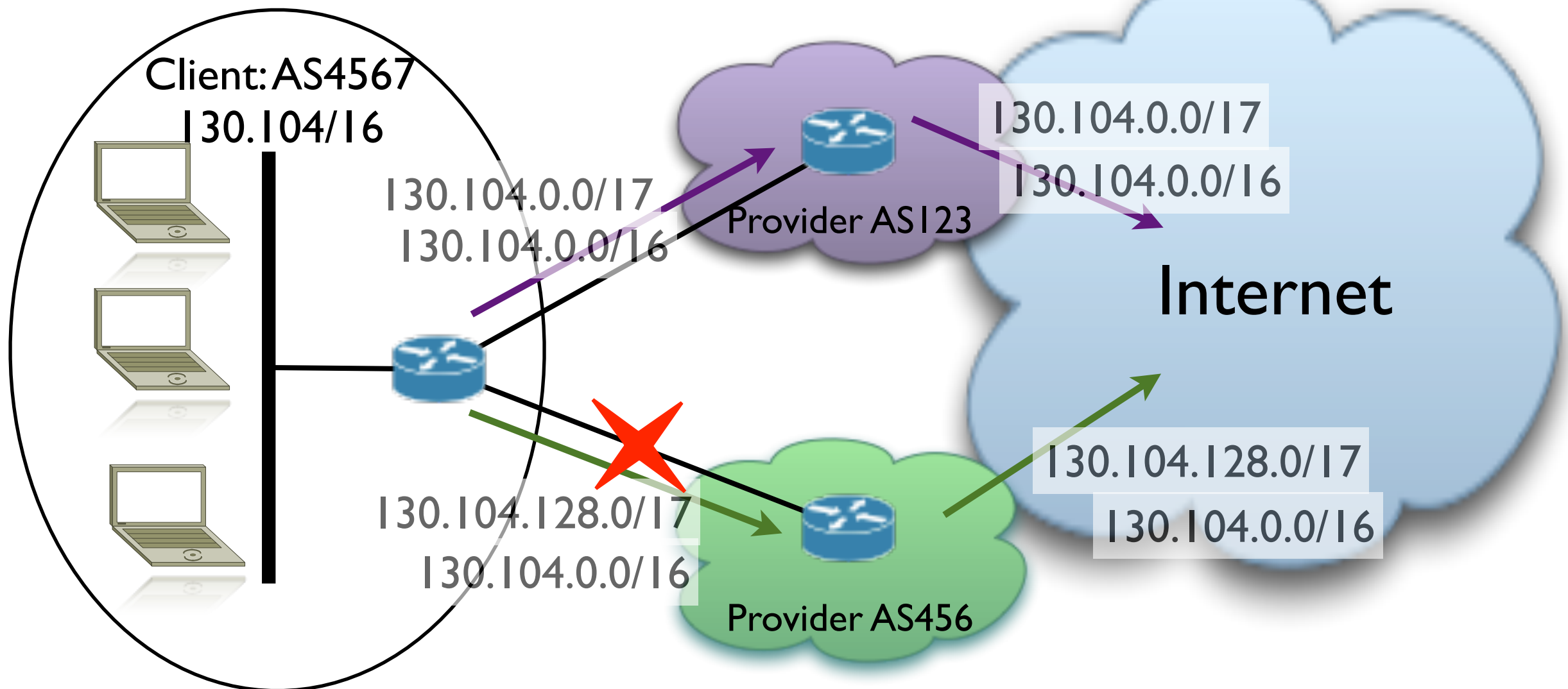
Reasons of the BGP growth?

- Traffic Engineering



Reasons of the BGP growth?

- Traffic Engineering



The Internet is broken!

- The IP addresses currently used by end-hosts play two complementary roles
 - **Identifier**: the IP address identifies (with port) the endpoint of transport flows
 - **Locator**: the IP address indicates the paths used to reach the end-host

The Cure?

The Locator/Identifier Separation Protocol (LISP)

`draft-ietf-lisp-05.txt`

LISP Main Design Goals

- Minimize required changes to Internet
- No end-systems (hosts) changes
- Be incrementally deployable
- No router hardware changes
- Minimize router software changes

The Locator Identifier Separation Protocol (I/2)

- Define a router-based solution where current IP addresses are separated in two different spaces:
 - **EndPoint Identifiers (EID)**
 - identify end-hosts
 - non-globally routable
 - hosts in a given site are expected to use EIDs in the same prefix
 - **Routing Locators (RLOC)**
 - attached to routers (router interfaces)
 - globally routable

The Locator Identifier Separation Protocol (2/2)

- Follows the Map-and-Encap principle
 - A **mapping system** maps EID prefixes onto site routers RLOCs
 - Routers **encapsulate** the packets received from hosts before sending them towards the destination RLOC
 - Routers **decapsulate** the packets received from the Internet before sending them towards the destination hosts

Example

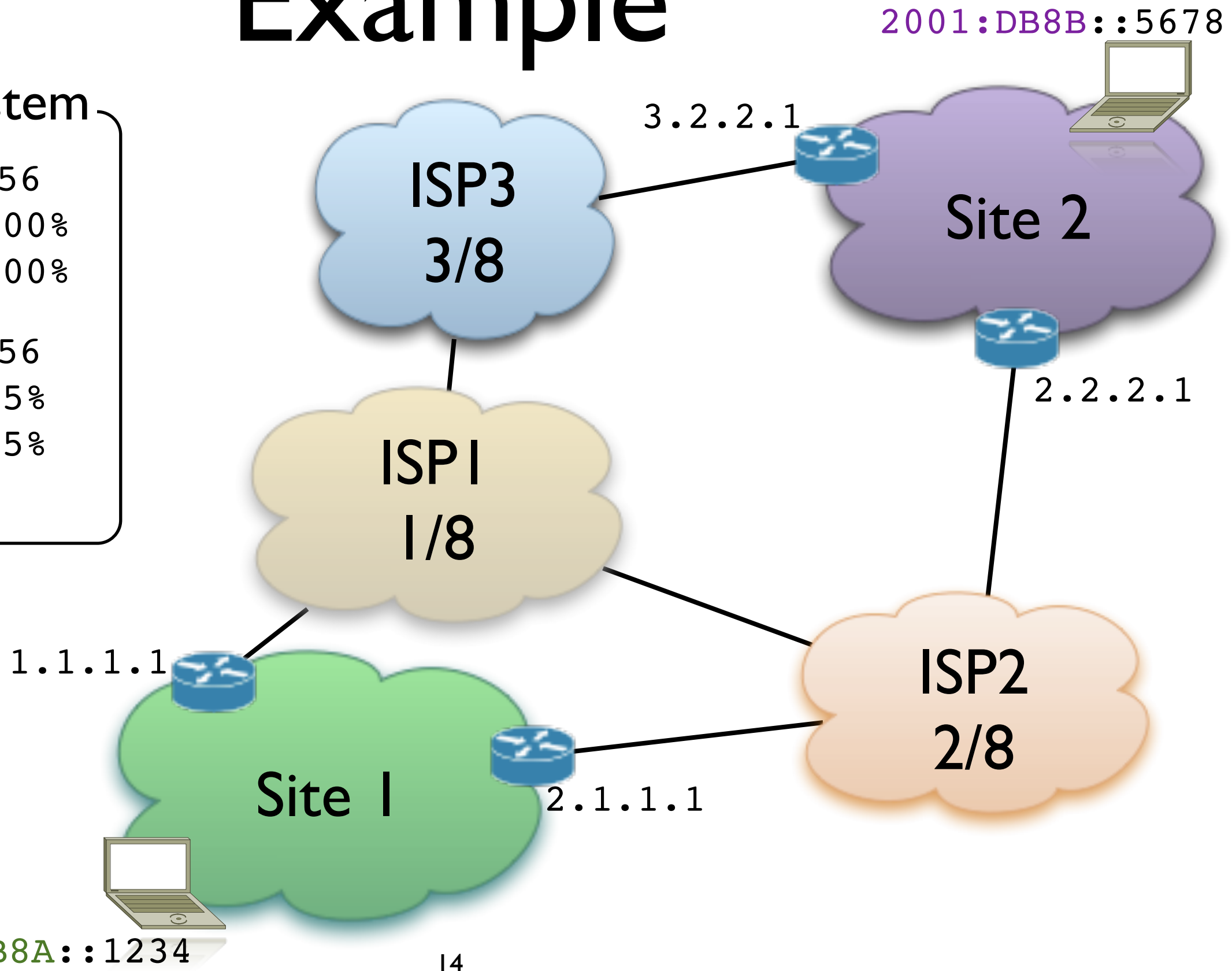
Mapping System

2001:DB8B::/56

- 3.2.2.1 1 100%
- 2.2.2.1 2 100%

2001:DB8A::/56

- 1.1.1.1 1 75%
- 2.1.1.1 1 25%



Example

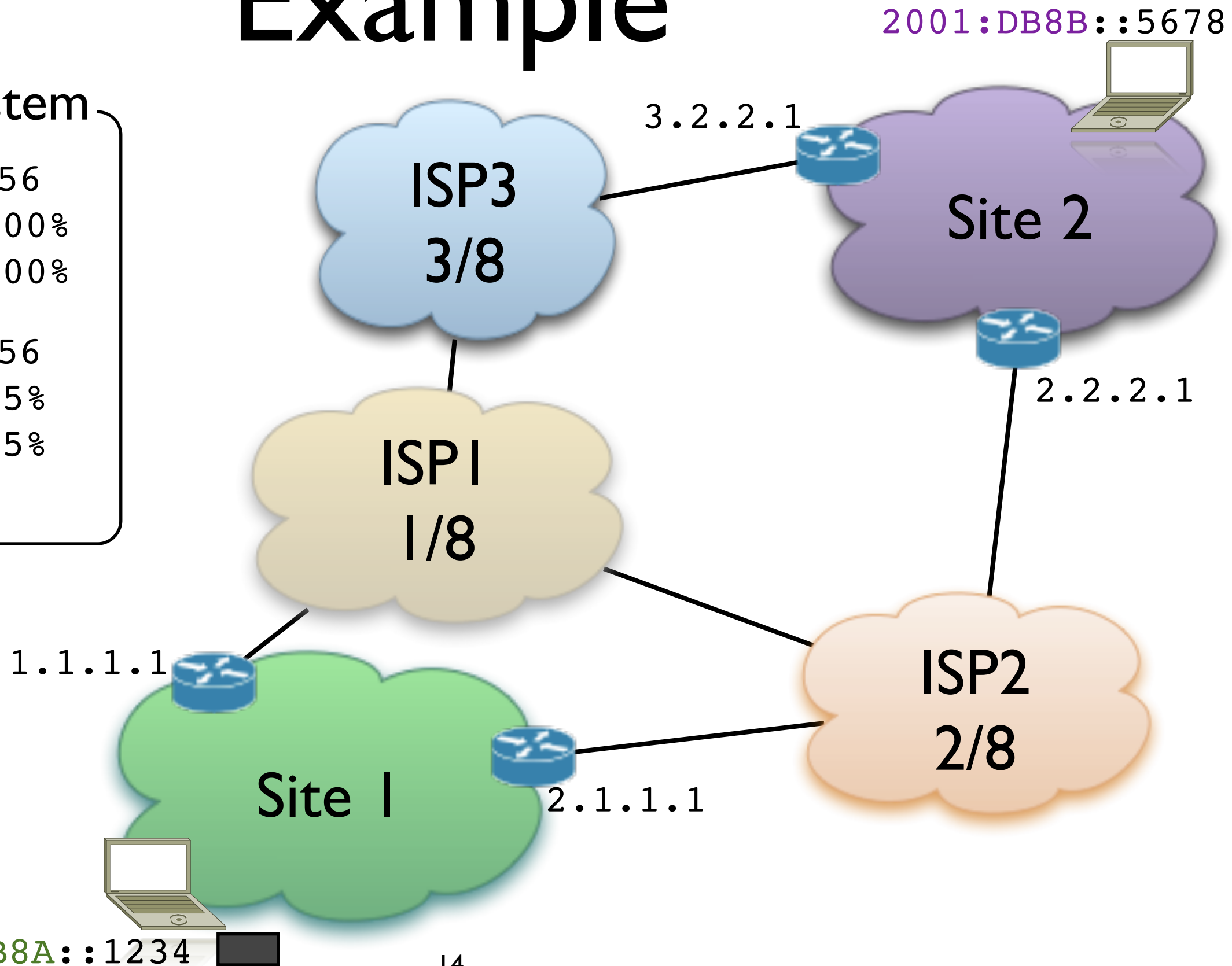
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2001:DB8A::1234

Example

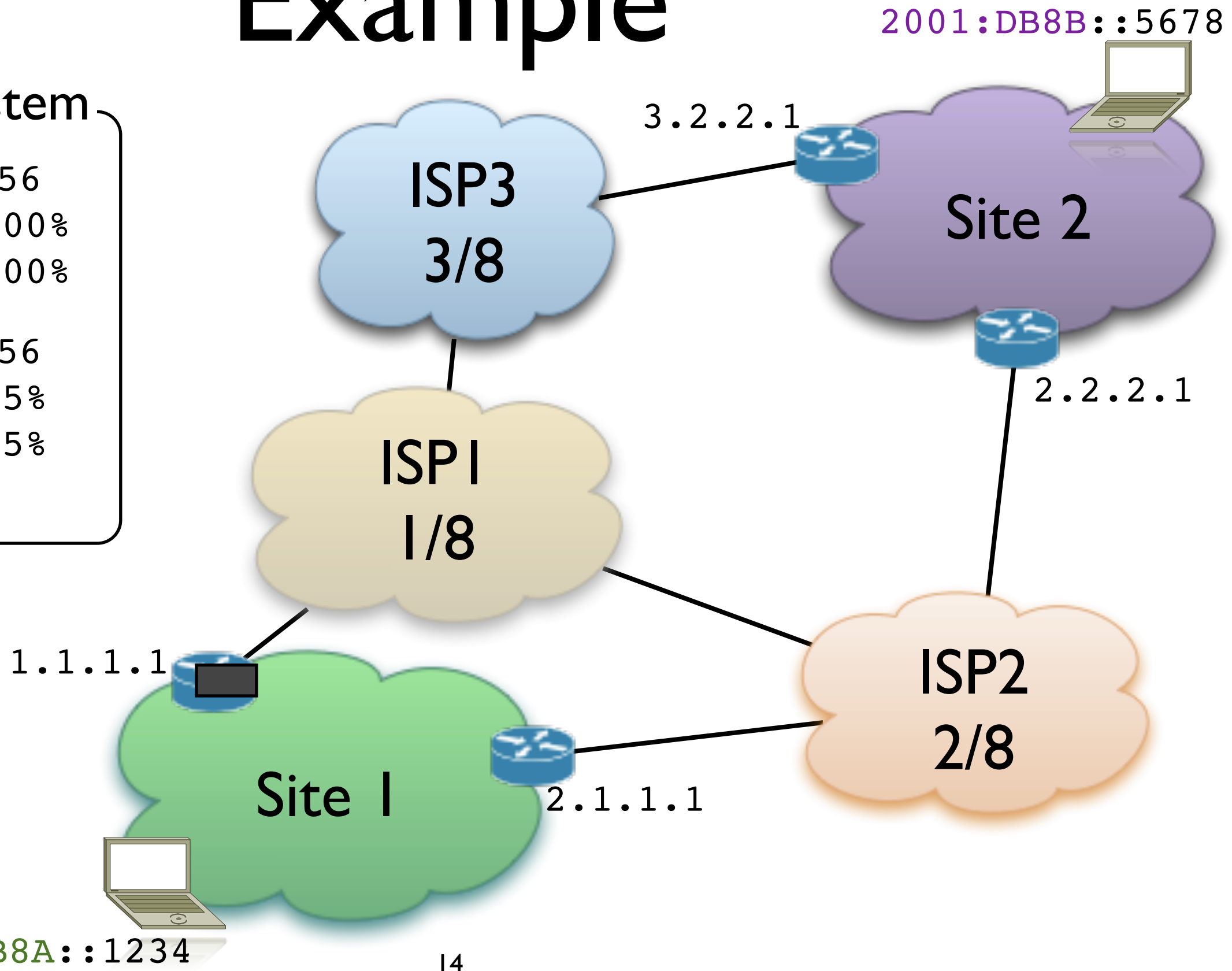
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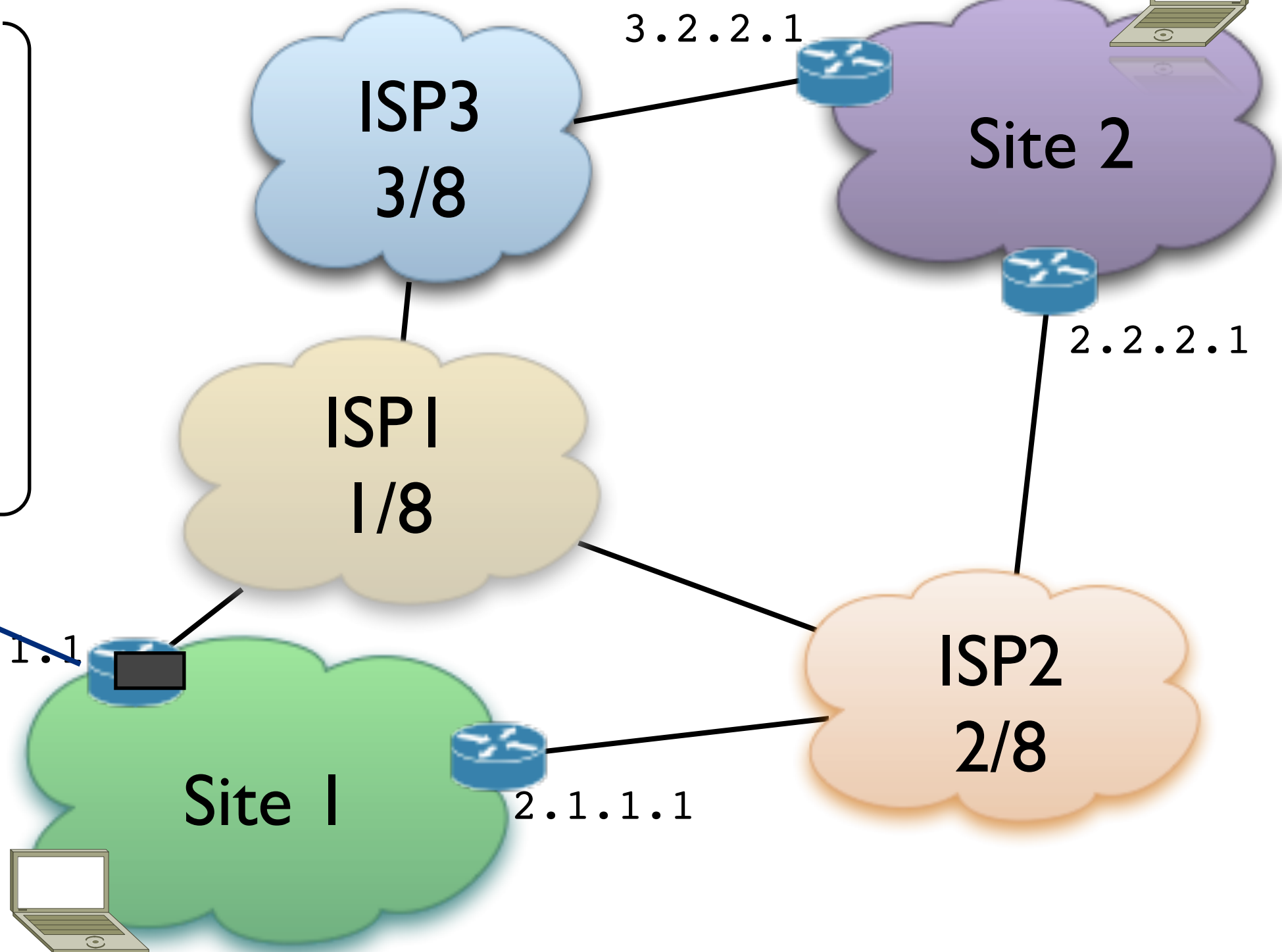
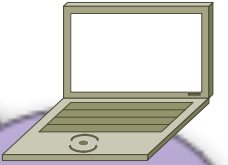
Map-Request:
2001:DB8B::5678?

1.1.1.1



2001:DB8A::1234

2001:DB8B::5678



Example

Mapping System

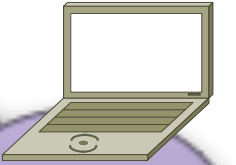
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2001:DB8A::/56

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- 2.1.1.1 1 25%

2001:DB8B::5678



3.2.2.1

ISP3
3/8

Site 2

2.2.2.1

ISP1
1/8

ISP2
2/8

Map-Reply:

2001:DB8B::/56

- 3.2.2.1 1 100%
- 2.2.2.1 2 100%

1.1.1.1

Site 1

2.1.1.1



2001:DB8A::1234

Example

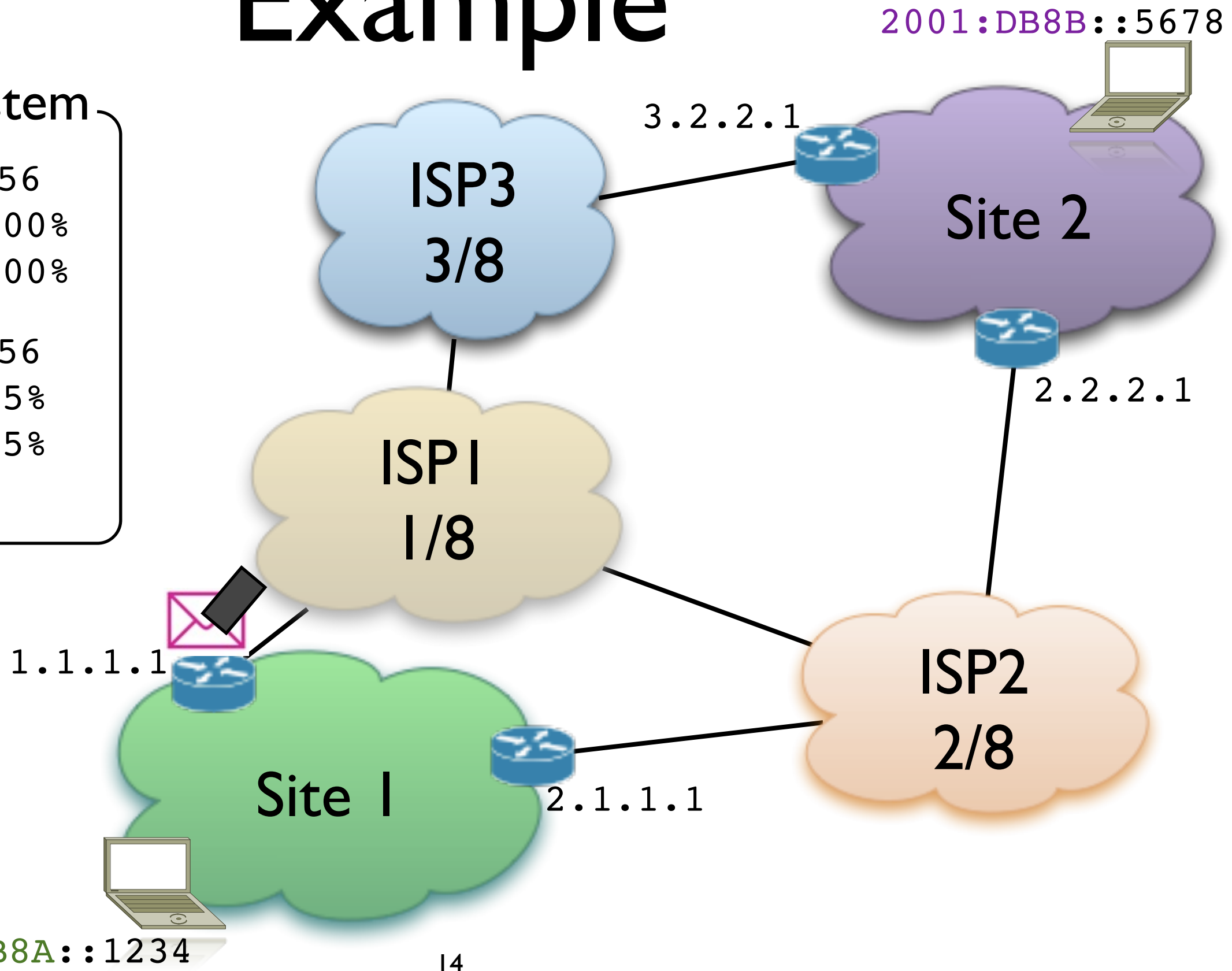
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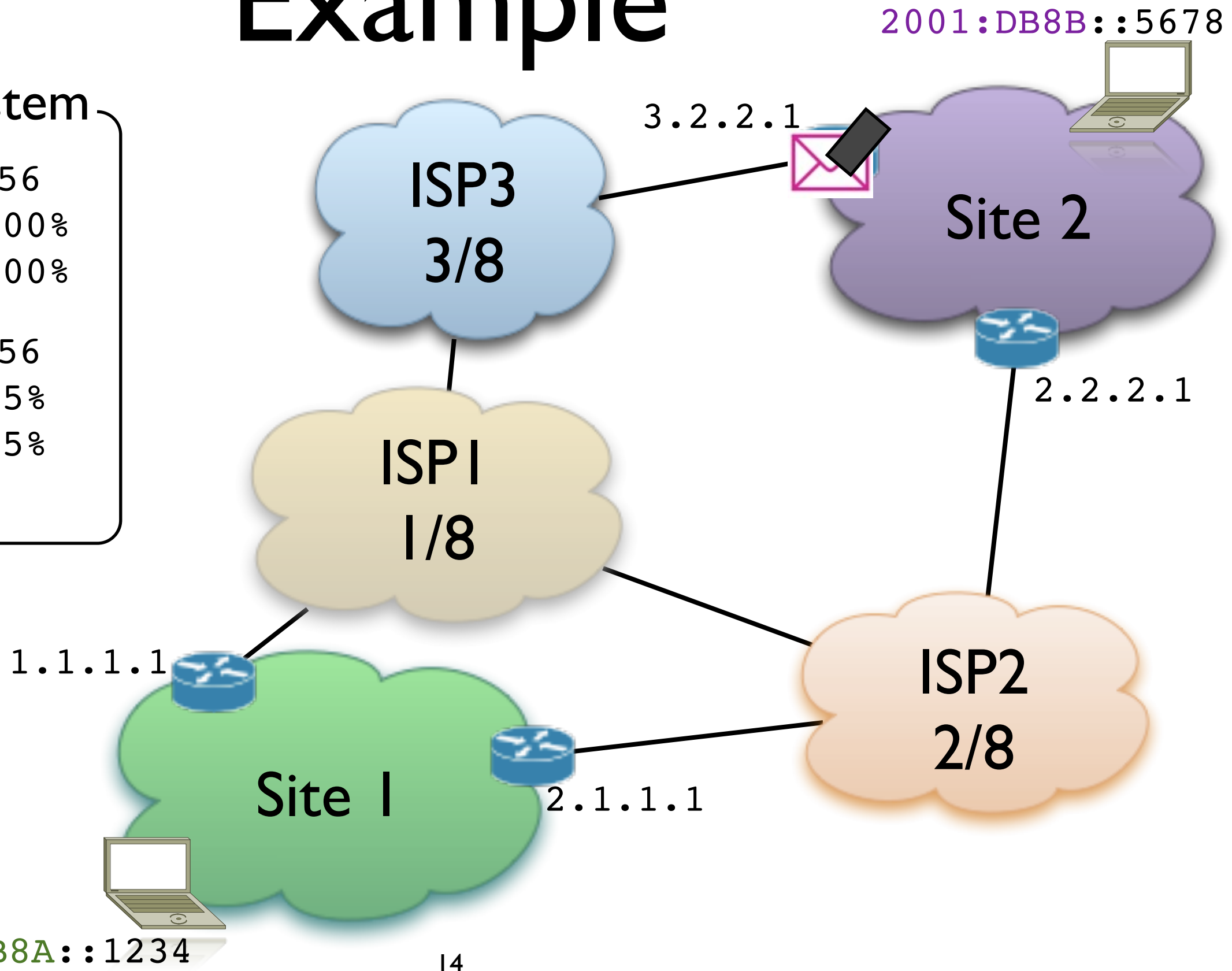
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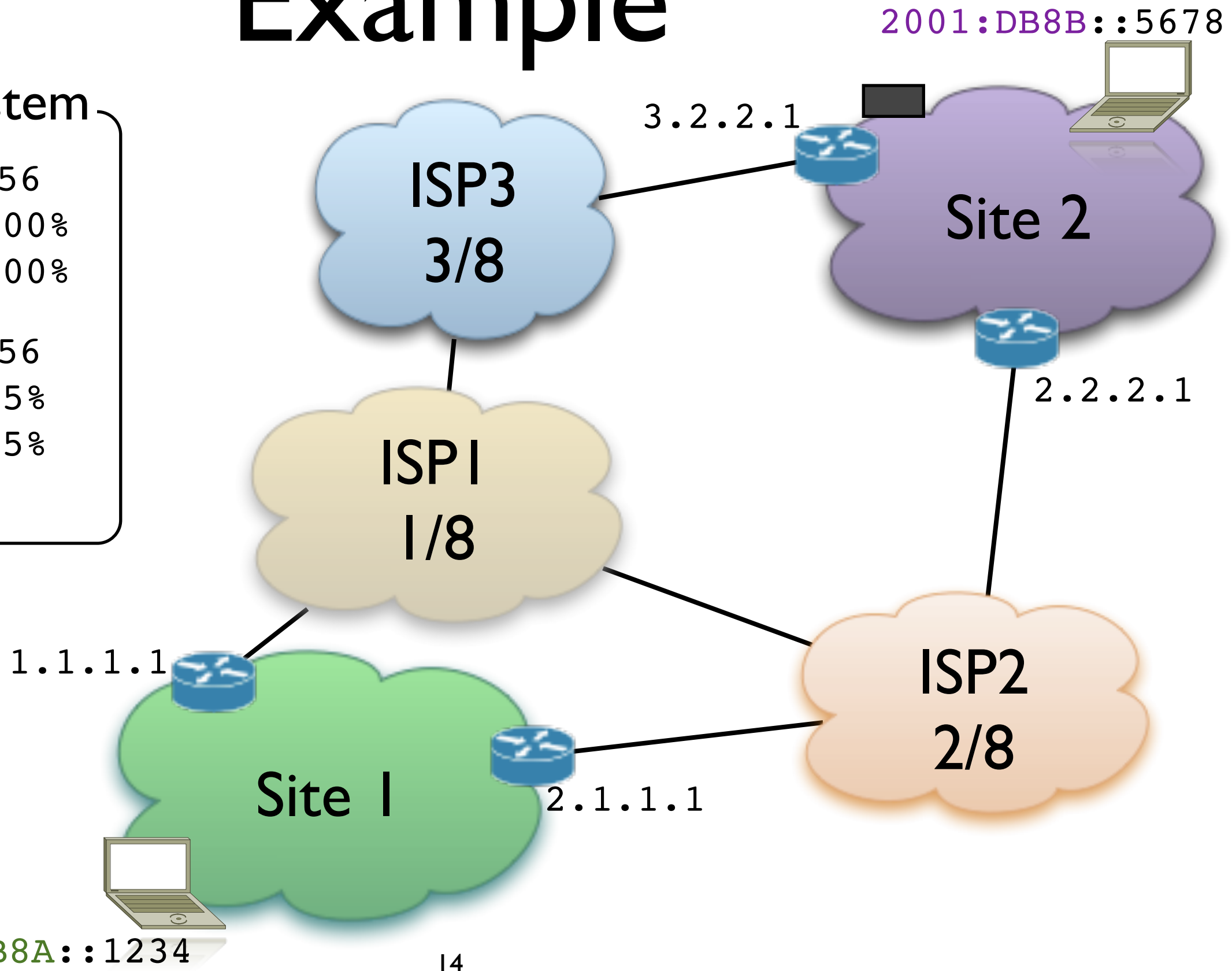
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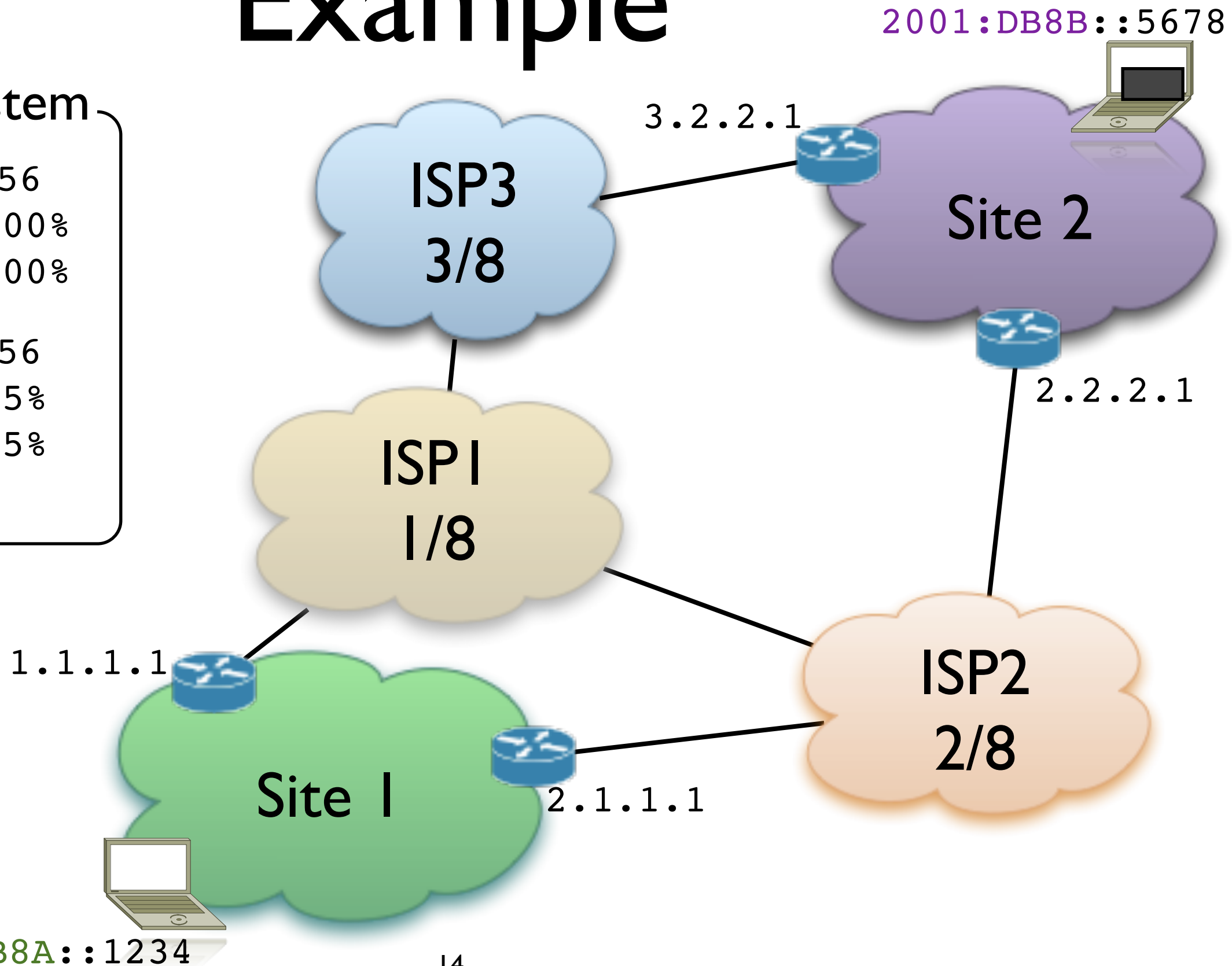
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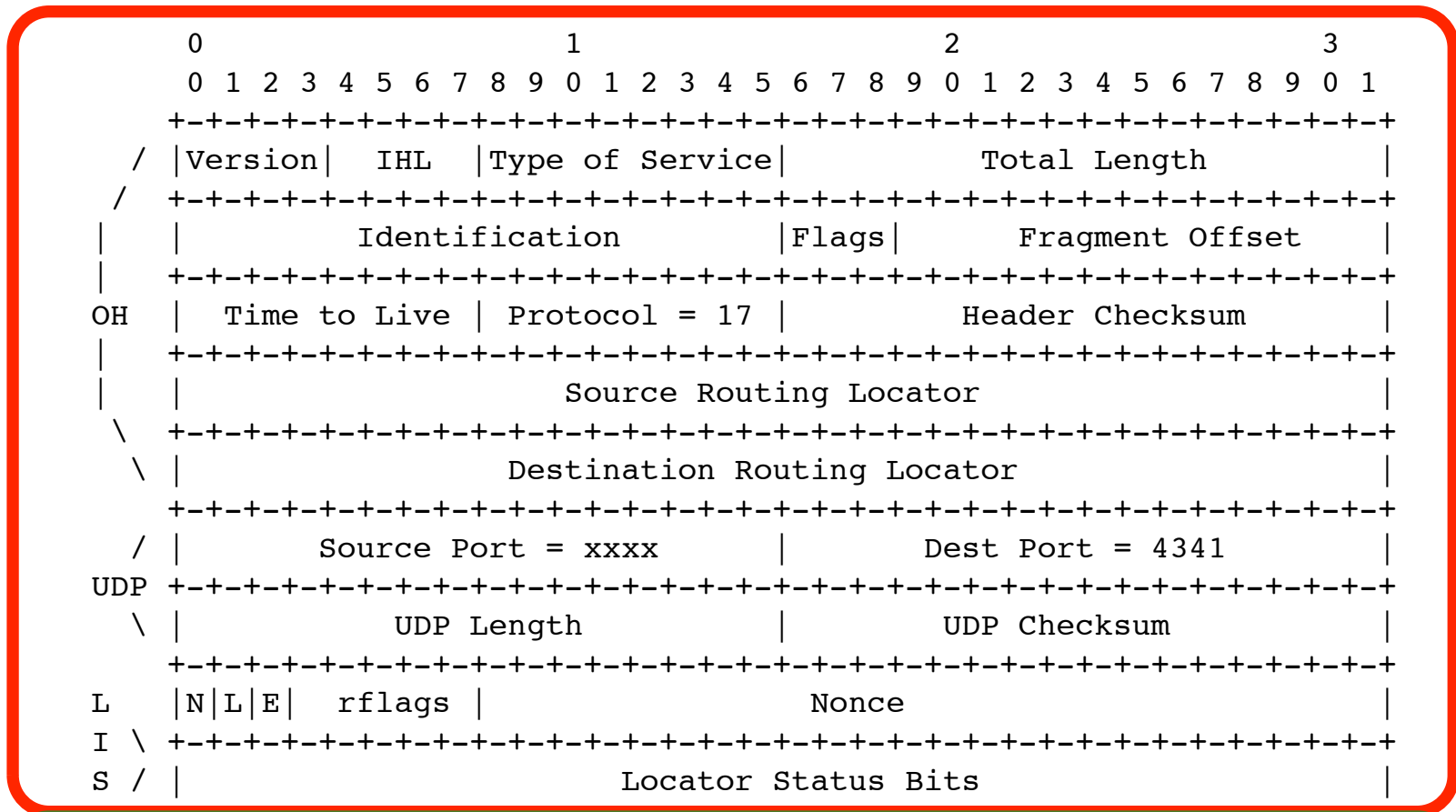
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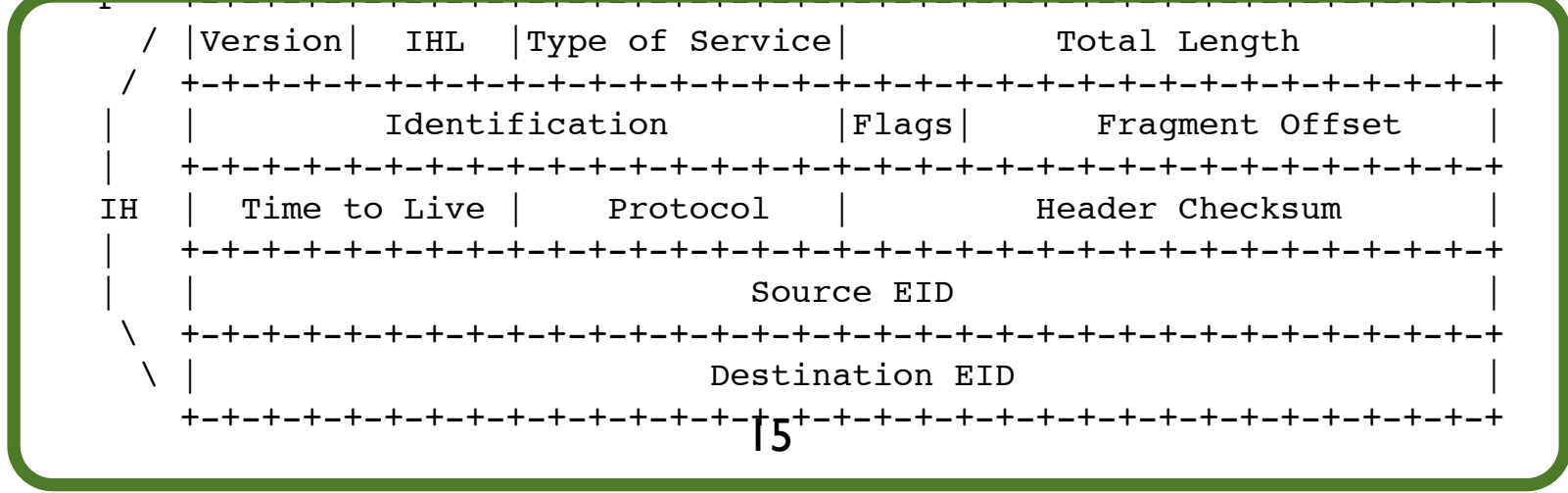


LISP Data Packet

(IP(UDP(LISP(IP))))



Encap



Payload

LISP Header

- `include/clicknet/lisp.h`

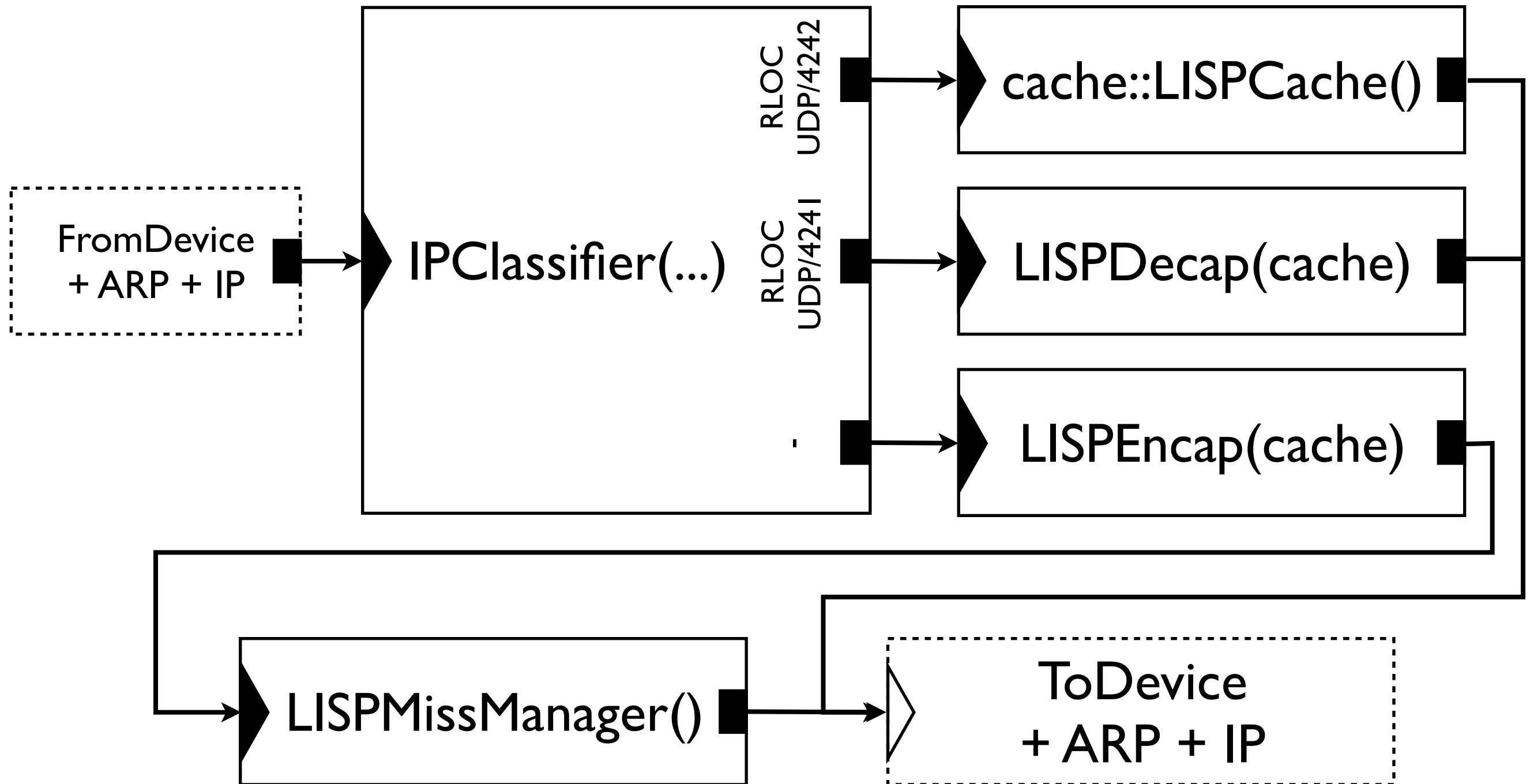
```
struct lisphdr {  
#if CLICK_BYTE_ORDER == CLICK_LITTLE_ENDIAN  
    unsigned rflags:5;  
    unsigned e_bit:1;  
    unsigned l_bit:1;  
    unsigned n_bit:1;  
#elif CLICK_BYTE_ORDER == CLICK_BIG_ENDIAN  
    unsigned n_bit:1;  
    unsigned l_bit:1;  
    unsigned e_bit:1;  
    unsigned rflags:5;  
#else  
#    error "unknown byte order"  
#endif  
    unsigned lisp_data_nonce:24;  
    uint32_t lisp_loc_status_bits;  
} __attribute__((__packed__));
```

Do you speak Click?

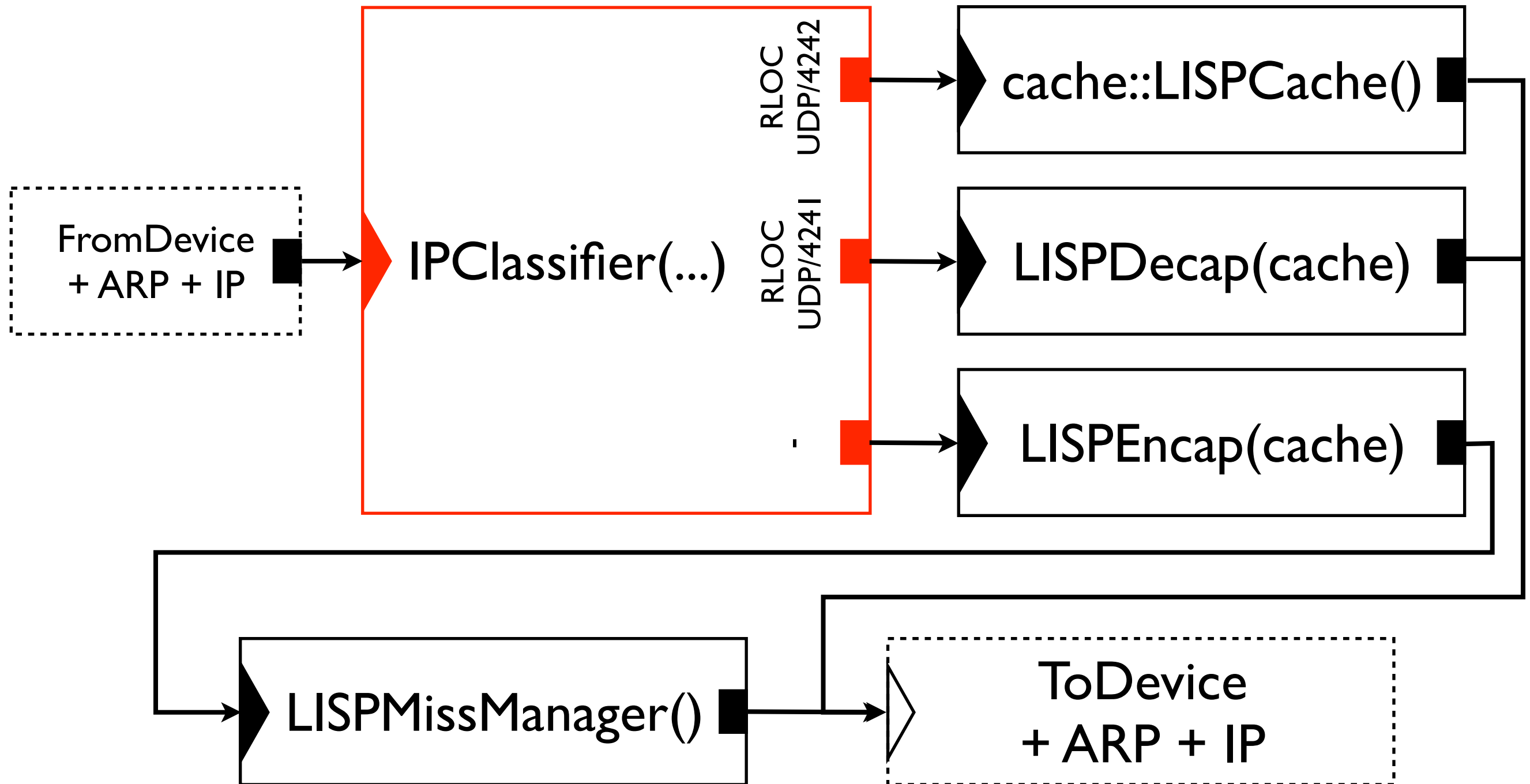
LISP-Click

- LISP Data Plane implementation
- user space
- interoperable with (x-test @ IETF76)
 - Cisco NX-OS implementation
 - OpenLISP
- very **experimental** code!

LISP-Click Graph



IPClassifier Element



IPClassifier Element

- Determine the action to take

1. Control plane packet

```
dst RLOC dst udp port 4342
```

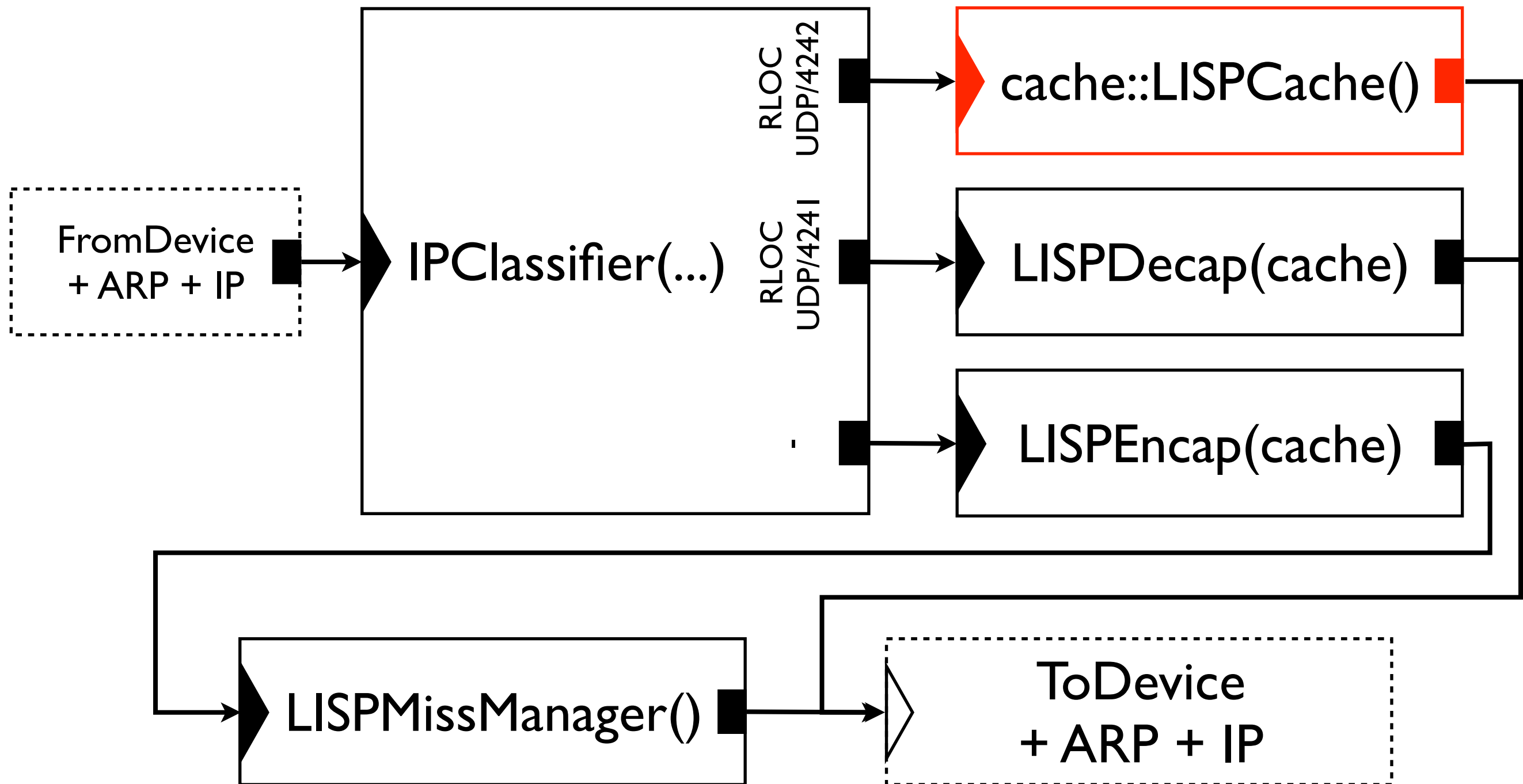
2. Decap

```
dst RLOC dst udp port 4341
```

3. Encap

-

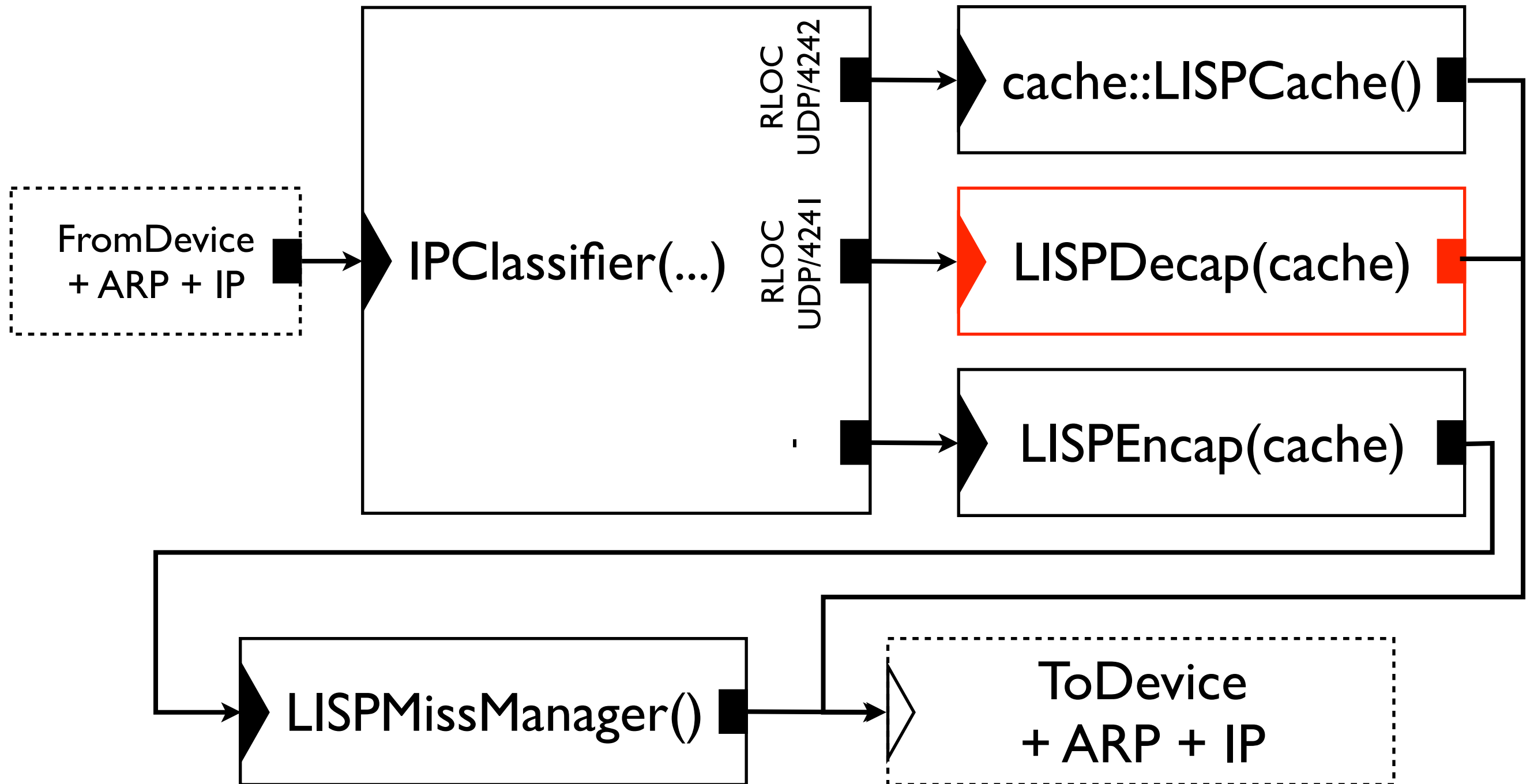
LISPCache Element



LISPCache Element

- Stores Mapping Database and Cache
 - 2 radix trees (Quagga implementation :-()
 - 2 priorities: primary OR backup
 - ECMP
- Stores RLOCs information
 - 1 Hashtable
- Control Socket (`write cache.insert/.update/.rloc`)
 - insert/update mappings
 - control RLOCs information
- Deal with control-plane packets (not implemented)

LISPDecap Element



LISPDecap Element

- Decap in Click
- `Packet::pull(len)`
 - *Remove a header from the front of the packet*

LISPDdecap Element

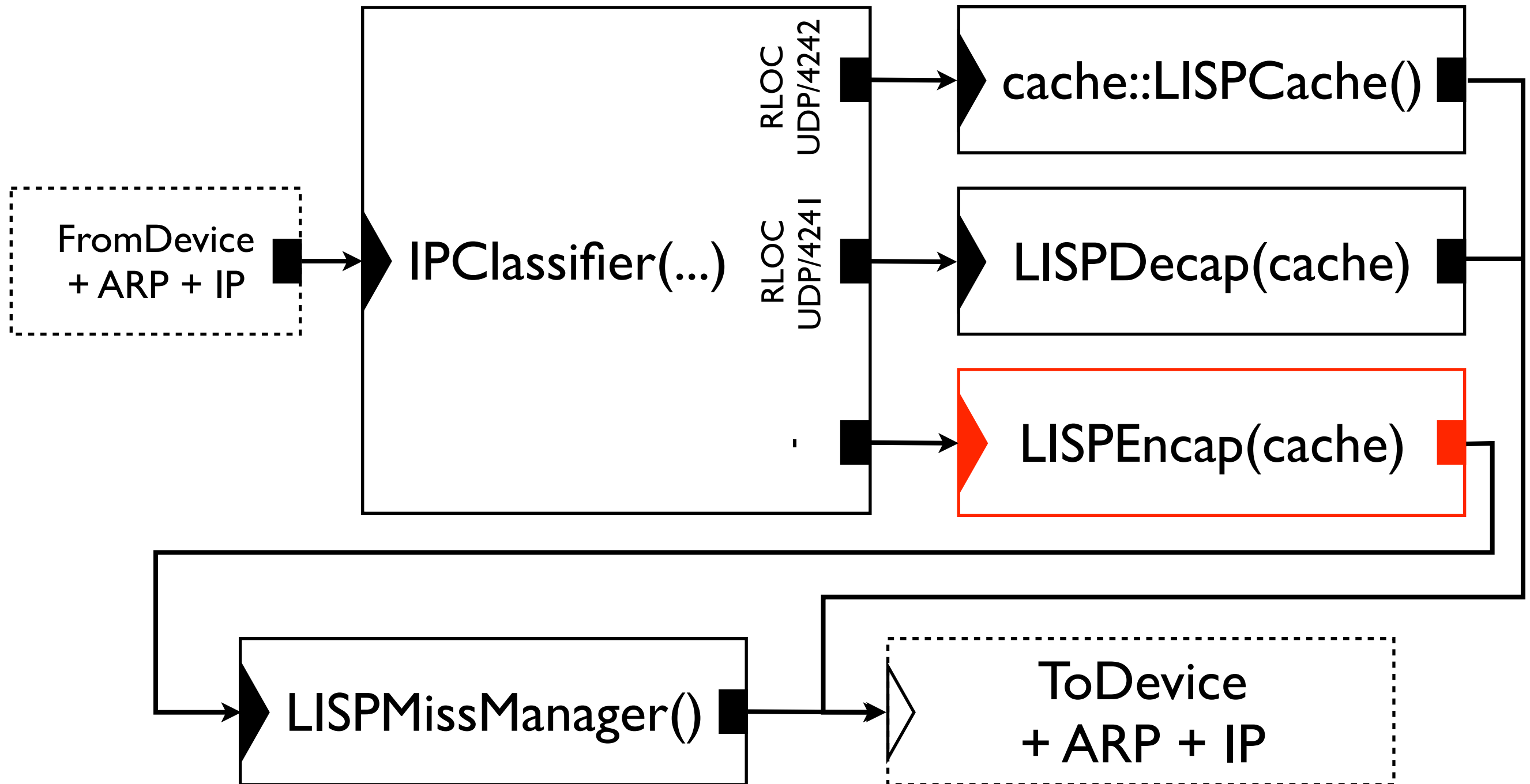
- elements/lisp/lispdecap.cc

```
Packet * LISPDdecap::simple_action(Packet *p_in)
{
    click_ip * ip = (click_ip *) (p_in->data());
    click_udp * udp = (click_udp *) (ip + 1);
    struct lisphdr * lisp = (struct lisphdr *) (udp + 1);
    click_ip * payload = (click_ip *) (lisp + 1);
    IPAddress deid = IPAddress(payload->ip_dst);

    ... // random stuff ;- )

    int hsize = sizeof(click_ip) + sizeof(click_udp)
                + sizeof(struct lisphdr);
    p_in->pull(hsize);
    p_in->set_dst_ip_anno(deid);
    return p_in;
}
```

LISPEncap Element



LISPEncap Element

- Encap in Click
- `Packet::push(len)`
 - *Add space for a header before the packet*

LISPEncap Element

- elements/lisp/lispencap.cc

```
void LISPEncap::push(int, Packet *p_in)
{
    IPAddress drloc = ... // chose a destination RLOC
    ... // random stuff ;- )

    int hsize = sizeof(click_ip) + sizeof(click_udp)
                + sizeof(struct lisphdr);
    WritablePacket *p_ = p_in->push(hsize);

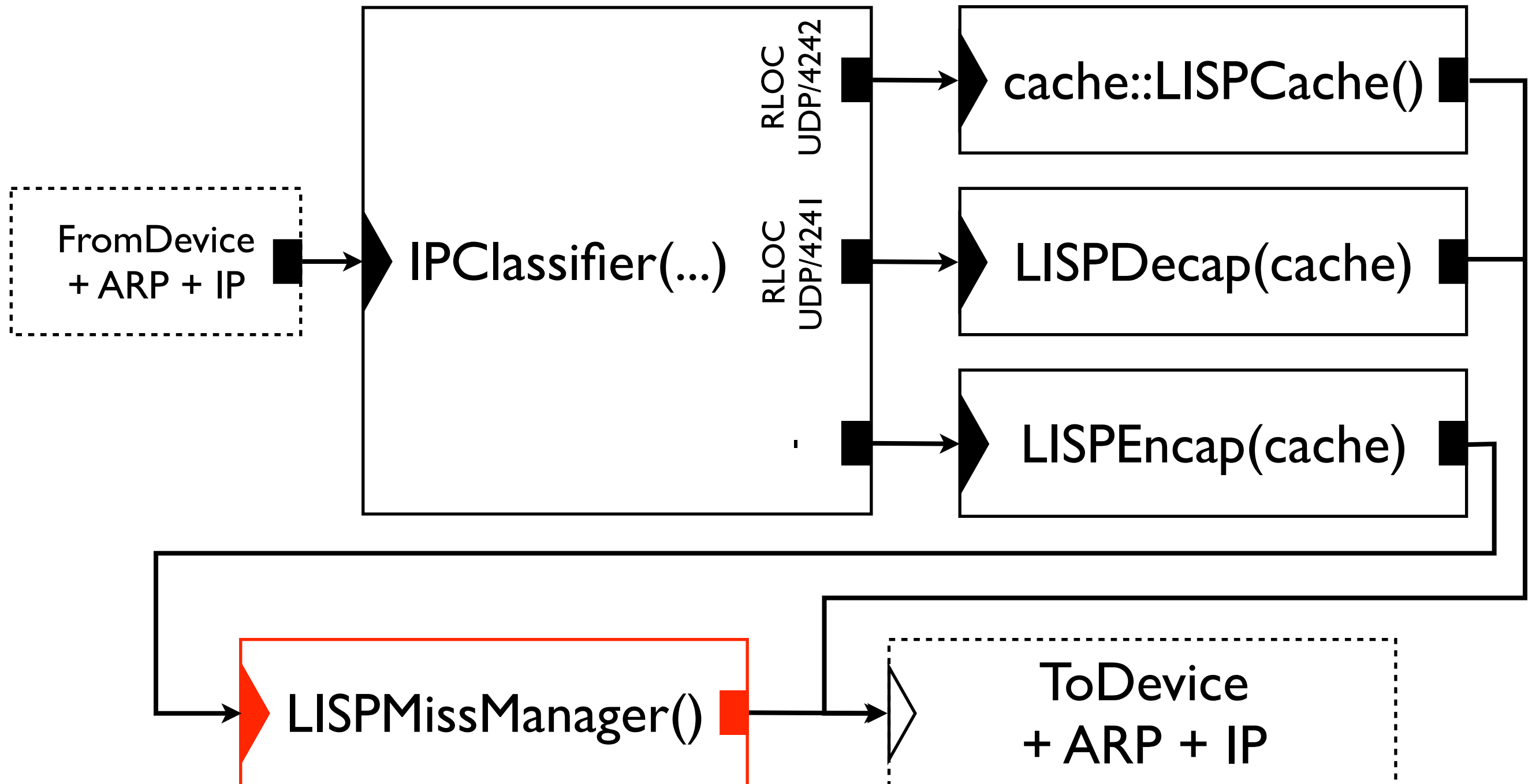
    ... // random stuff ;- )

    p_->set_dst_ip_anno(drloc);
    output(0).push(p_);
}
```

LISPEncap Element

- On cache miss, annotate the packet with `SET_MISS_ANNO`

LISPMissManager Element



LISPMissManager Element

- if packet is annotated
 - Send Map-Request for the EID (not implemented)
 - Drop the packet
- otherwise, send the packet

Further work

- IPv6 and cross AFI
- sanity checks
- smart source RLOC selection
- Full Control Plane
 - outside Click, via the Control Socket?

Want to see more?

- Meet @ 1:45pm for a live demo!

Need your help

- How to avoid ICMP port unreachable error?
- for the moment: run a dummy server listening on UDP port 4341

FromDevice(eth0, SNIFFER 0)

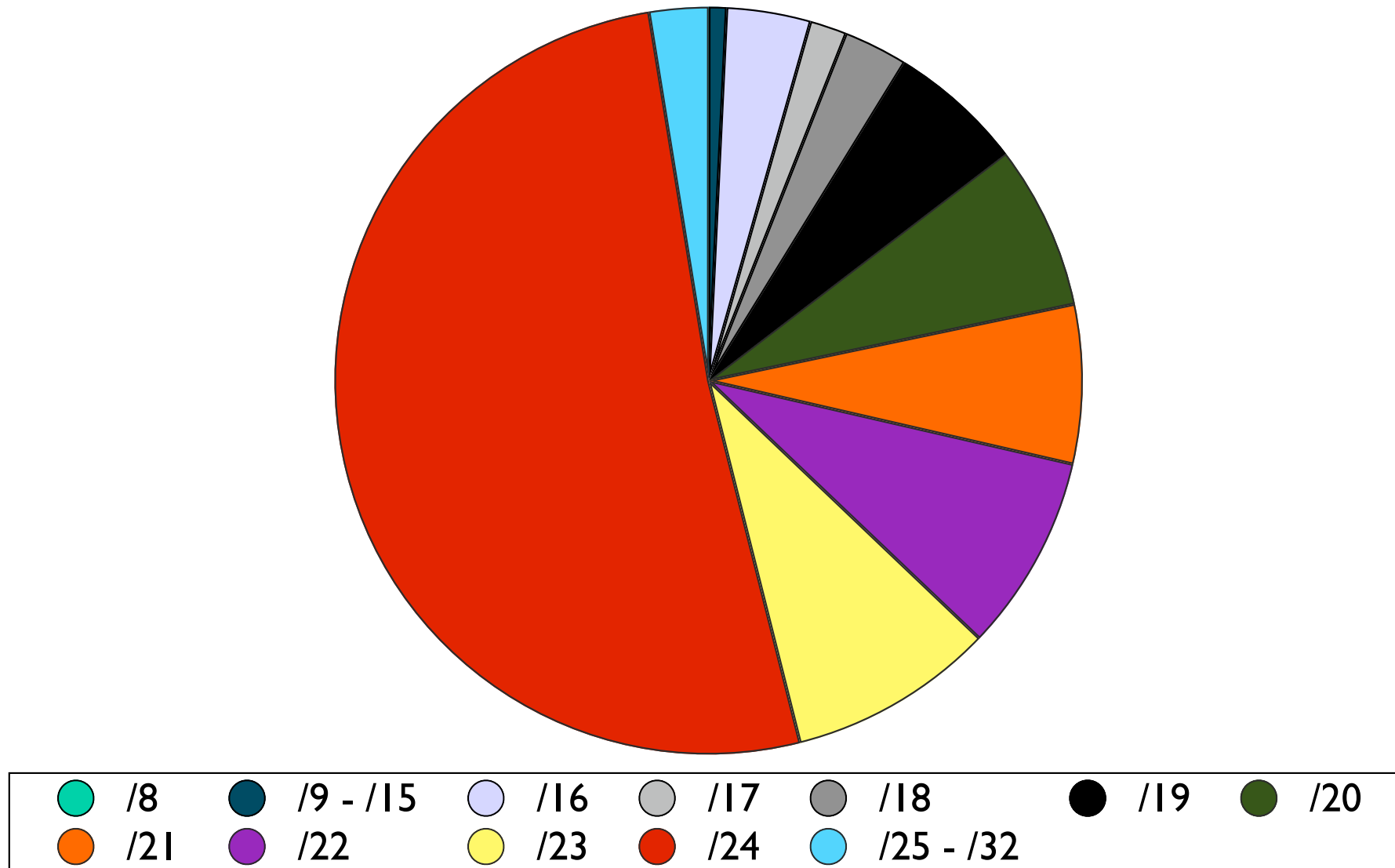
Thank you

<http://inl.info.ucl.ac.be>

Backup

Reasons of the BGP growth?

Prefix length distribution



Reasons of the BGP growth?

- Before CIDR: Classfull network
 - ask IANA for n addresses => receive a full class
 - Class C: up to 256 addresses
 - Classe B: up to 65,536 addresses
 - Classe A: up to 16,777,216
- First come, first served

18.0.0.0/8	MIT
19.0.0.0/8	Ford Motor Company
20.0.0.0./8	Computer Sciences Corporation

Reasons of the BGP growth?

- With CIDR: classless network
 - need n addresses \Rightarrow receive a $/\log(n)$ prefix
 - Provider Independent (PI)
 - Owned by sites and globally announced
 - Provider Aggregatable (PA)
 - Given by ISPs from their own address block to customers. Only announce the ISP prefix