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Cisco Systems, Routing Symposium Monday, Oct. 5 2009

Introduction and motivation

Implementing CRS

Practical considerations and solutions

Conclusion

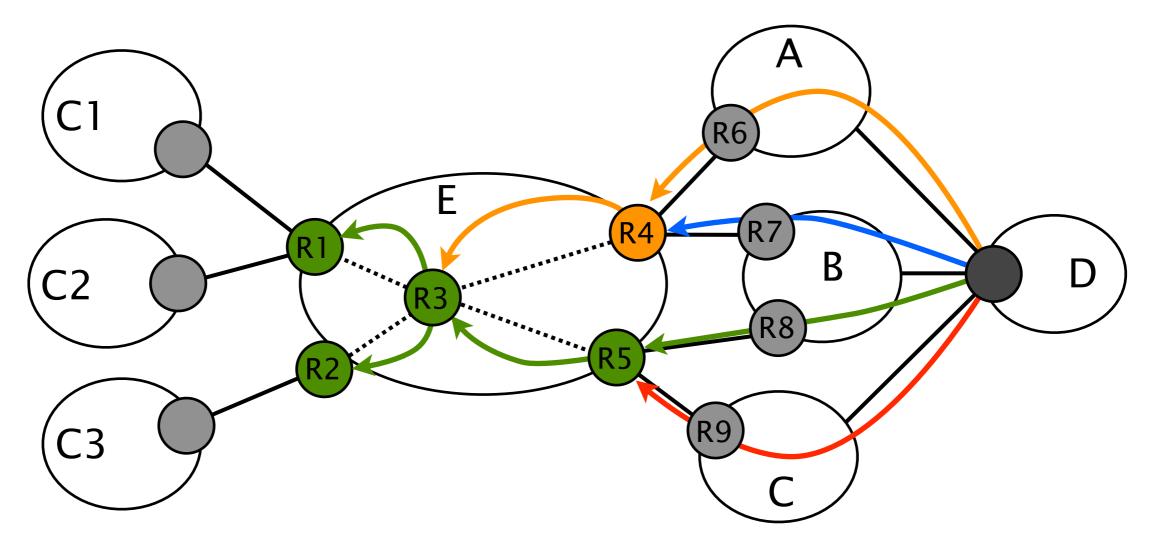
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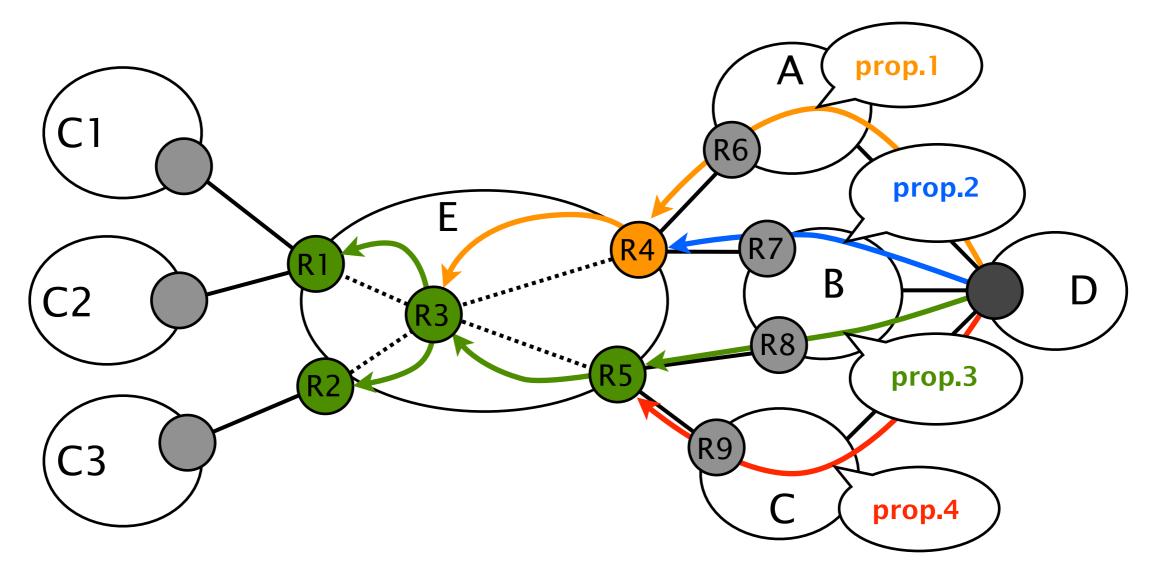
Practical considerations and solutions

Conclusion

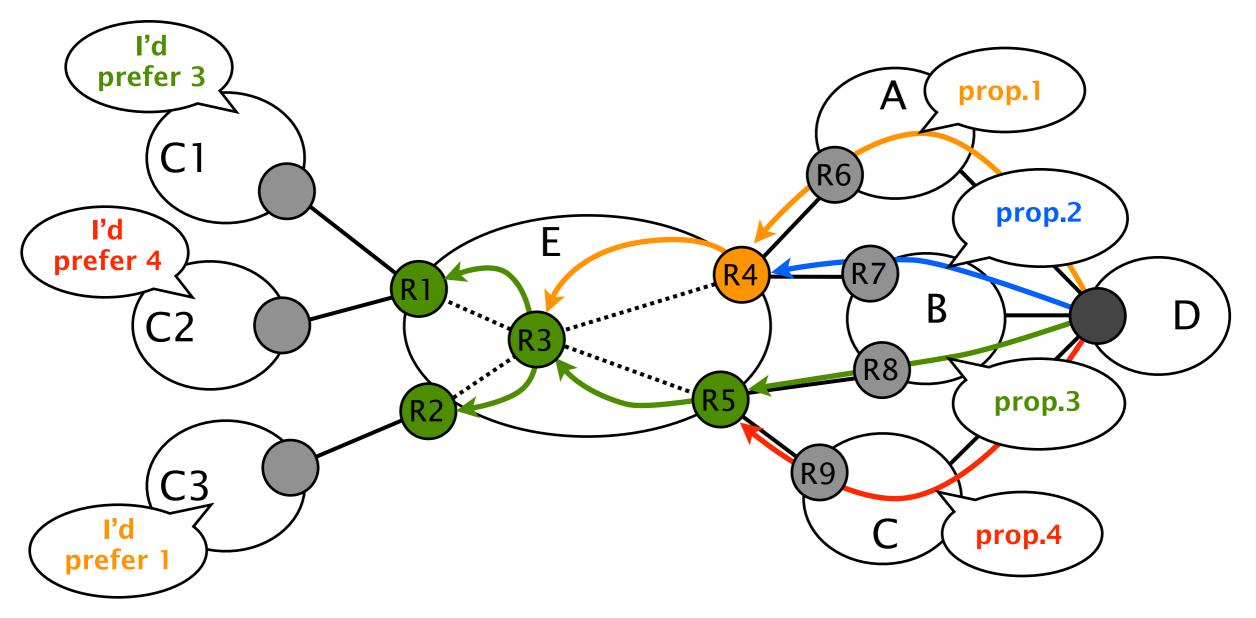
- A BGP router selects one best route for each destination
- Globally, AS E knows 4 paths towards D
 - Locally, some routers only know one path (*e.g.*, C1...C3)



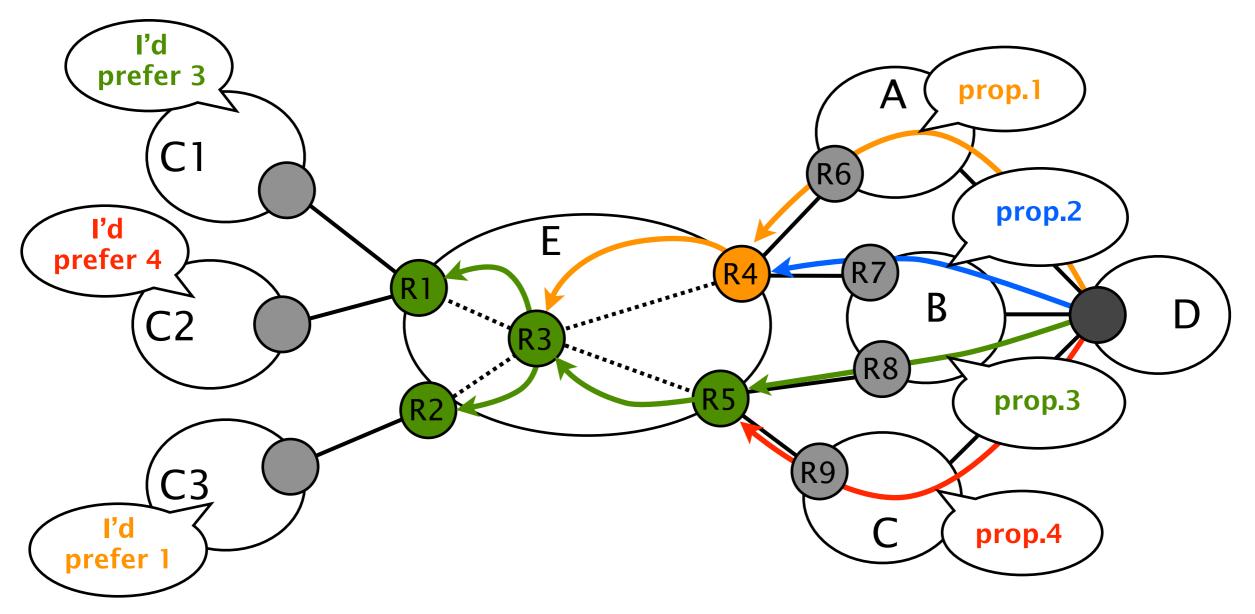
- Many ISPs have a rich path diversity
 - It is common to have 5-10 paths per prefix
- Different paths have different properties
 - It could be in terms of security, policies, etc.



- Clients may want different paths to the same prefix
 - If C1 is a competitor of C, he'd prefer to reach D via A or B
 - C1 may even want to pay an extra fee for that

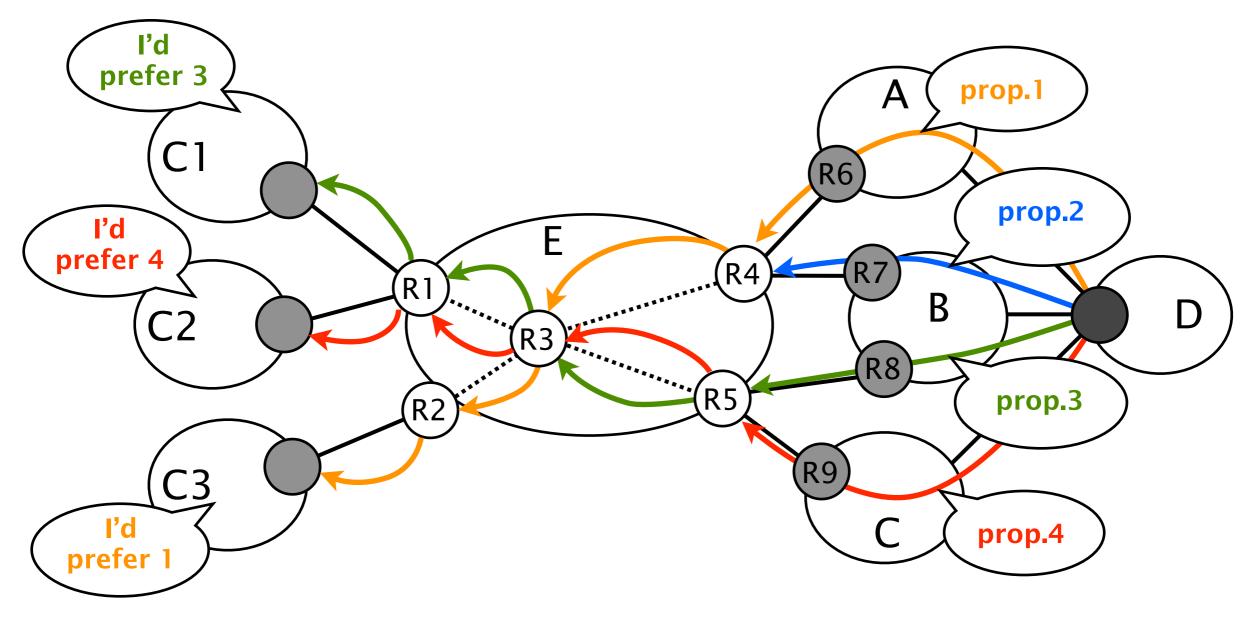


- With vanilla BGP, you can't match customers' preferences to available paths
 - Customers of a given PE receive the same path



CRS: Customized Route Selection

- Under CRS, one router can offer *different* interdomain routes to *different* neighbors
 - C1 reaches D via B, C2 reaches D via C



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Potential issues and solutions

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Two notions: class and service

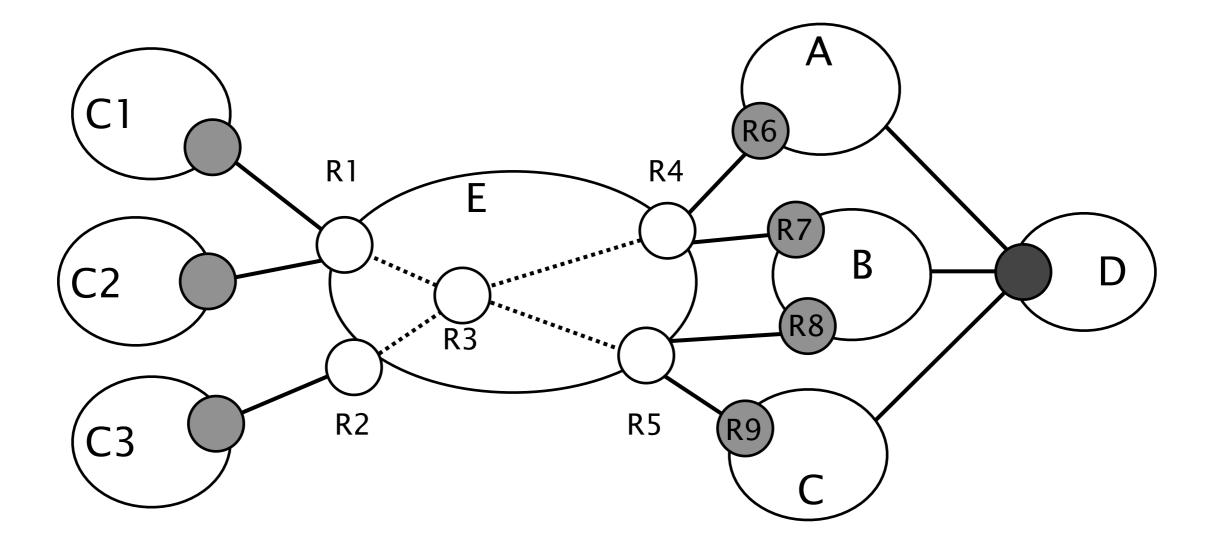
- A *class* is a set of routes sharing a property
 - *e.g.,* all the routes learned via provider *X*
 - One route can belong to more than one class
- A service is the union of one or more classes
 - Some classes can be preferred over others
 - e.g., service Y is the union of class 1 and class 2 where preference is given to class 1

What do we need to implement CRS with BGP MPLS VPNs ?

Mechanisms to disseminate and differentiate paths

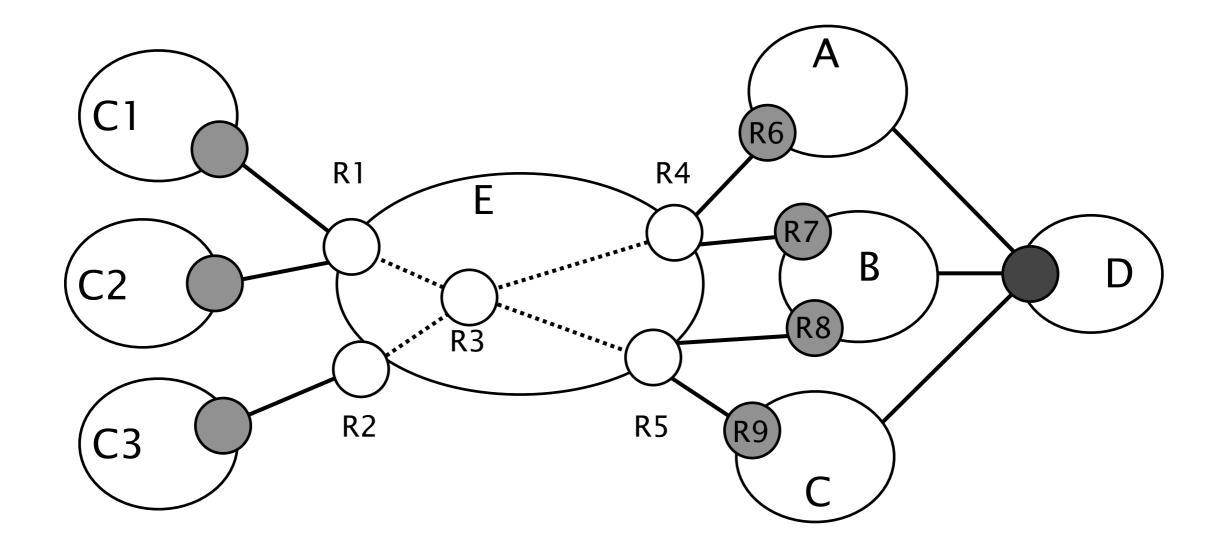
- Multiprotocol BGP is used as dissemination protocol
- Route Targets (RT) are used to identify classes
- Route Distinguishers (RD) are used to ensure diversity
- Customized route selection mechanisms at ASBR
 - Use Virtual Routing and Forwarding (VRF) instances to build services
- Traffic forwarding on the chosen paths
 - MPLS tunneling

How do we implement CRS with BGP MPLS VPNs ?

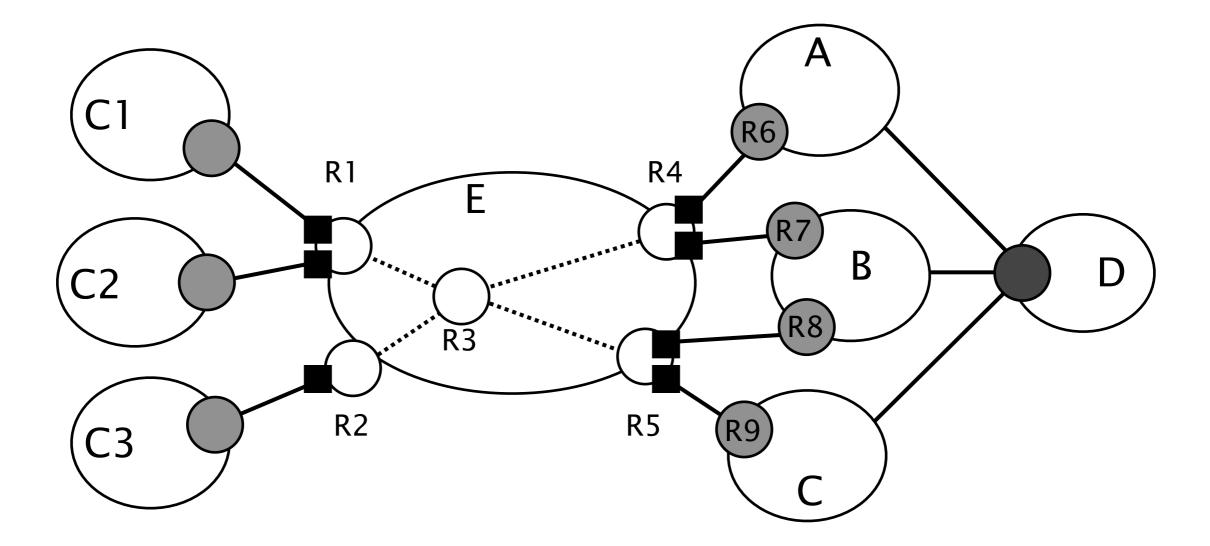


- C1 wants to reach D via B, C2 via C
- Define two services on R1: prefer B (resp. C) routes
- Define three classes: *learned* via A, B or C

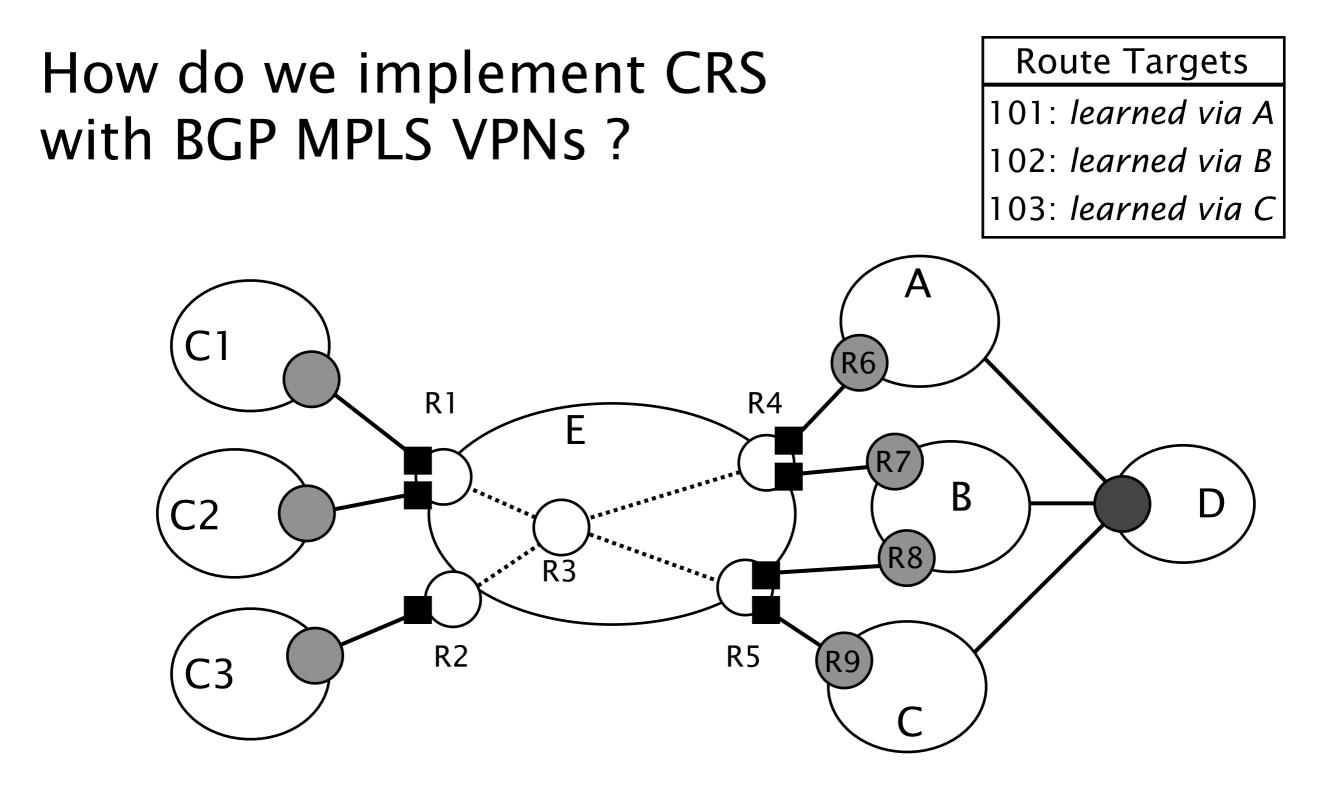
How do we implement CRS with BGP MPLS VPNs ?



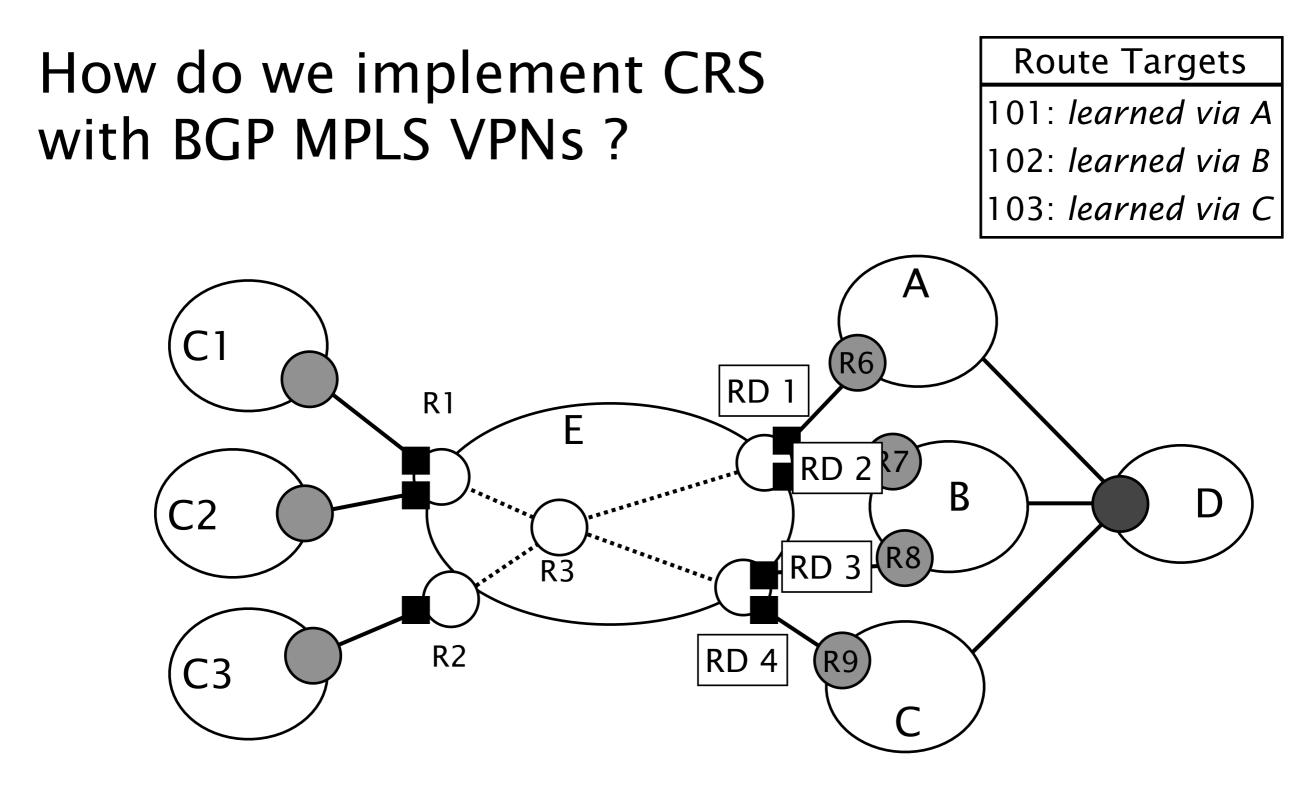
How do we implement CRS with BGP MPLS VPNs ?



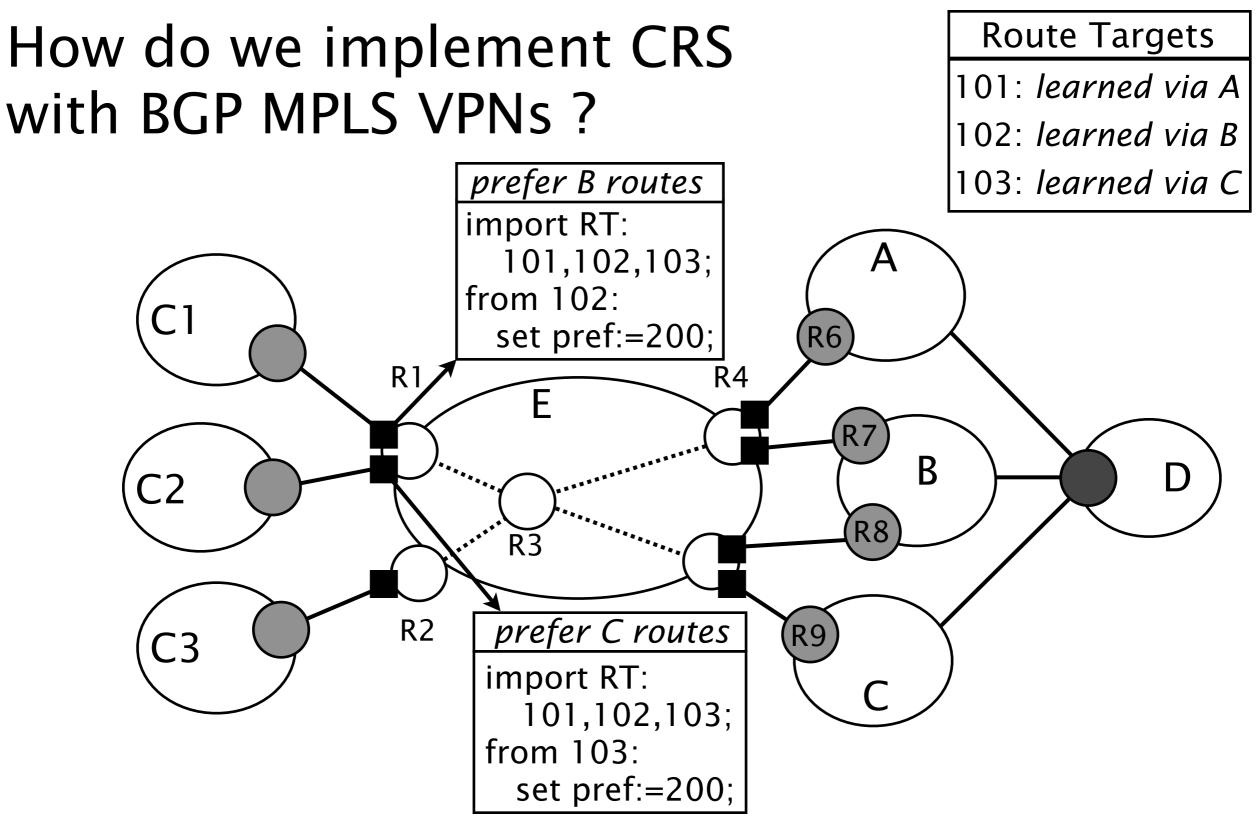
Consider peers as VPNs and put them in VRFs



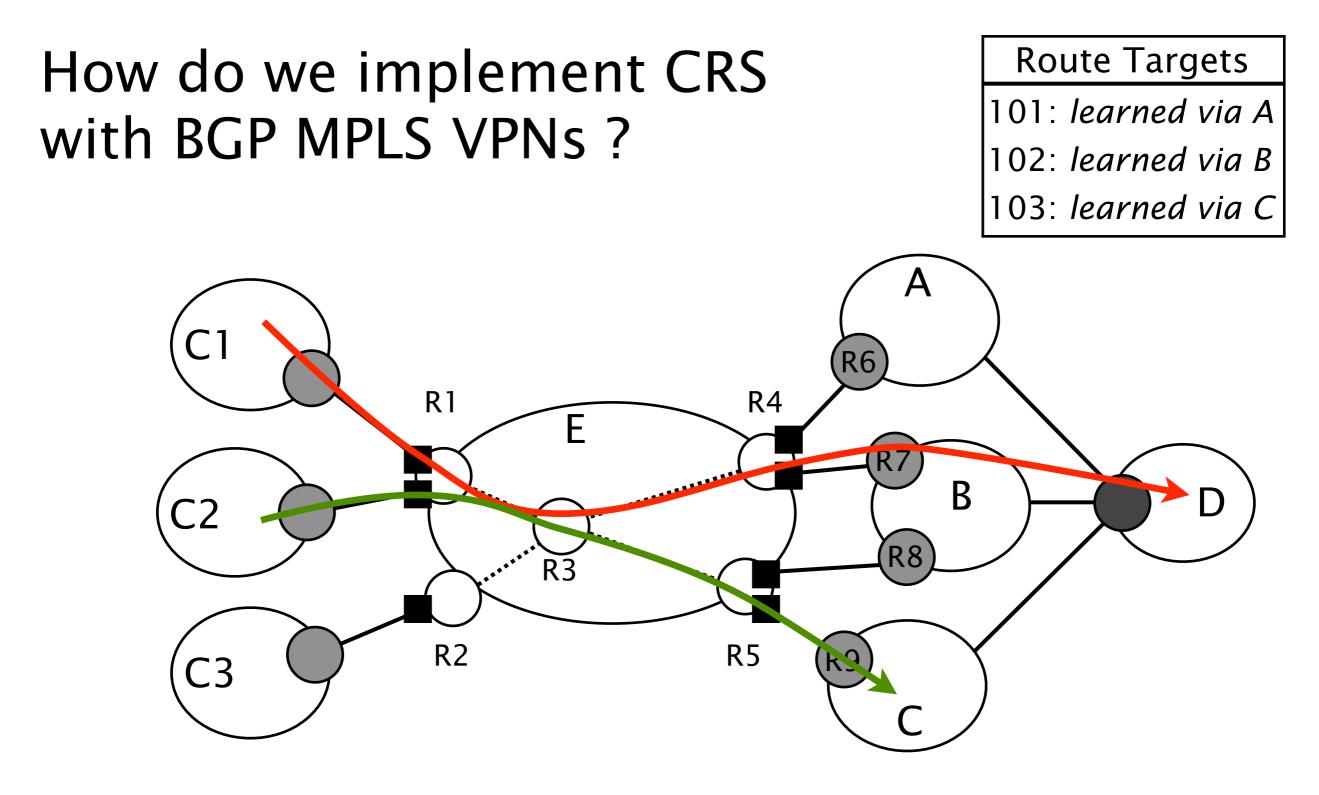
Consider peers as VPNs and put them in VRFs
Use RT to identify *classes*



- Consider peers as VPNs and put them in VRFs
- Use RT to identify classes
- Use different RD to differentiate routes



Define services by using VRFs' import filters



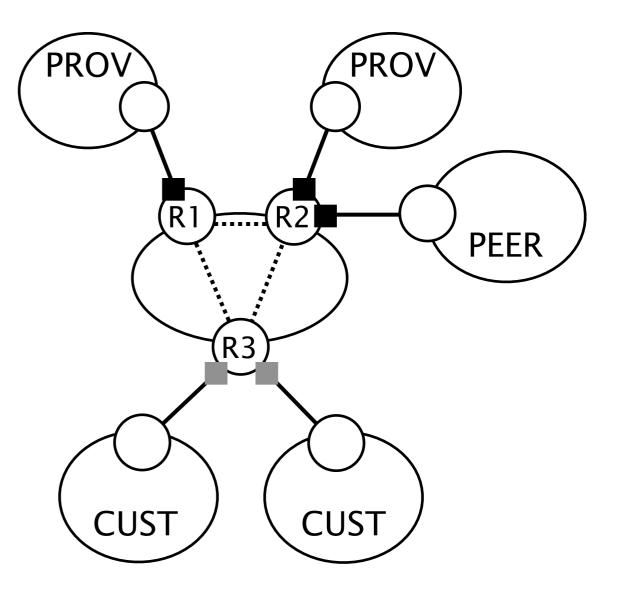
MPLS is used for forwarding

Two levels label stack

R3 only knows label to reach the PEs

CRS applied to *classical* policies

- Define three classes
 - Providers (RT 100)
 - Peers (RT 101)
 - Customers (RT 102)
- Define two services
 - VRF Provider/Peer (■)
 - import RT 102;
 - VRF Customers (■)
 - import RT 100,101,102;
- Thanks to VRF isolation, policies violations vanish



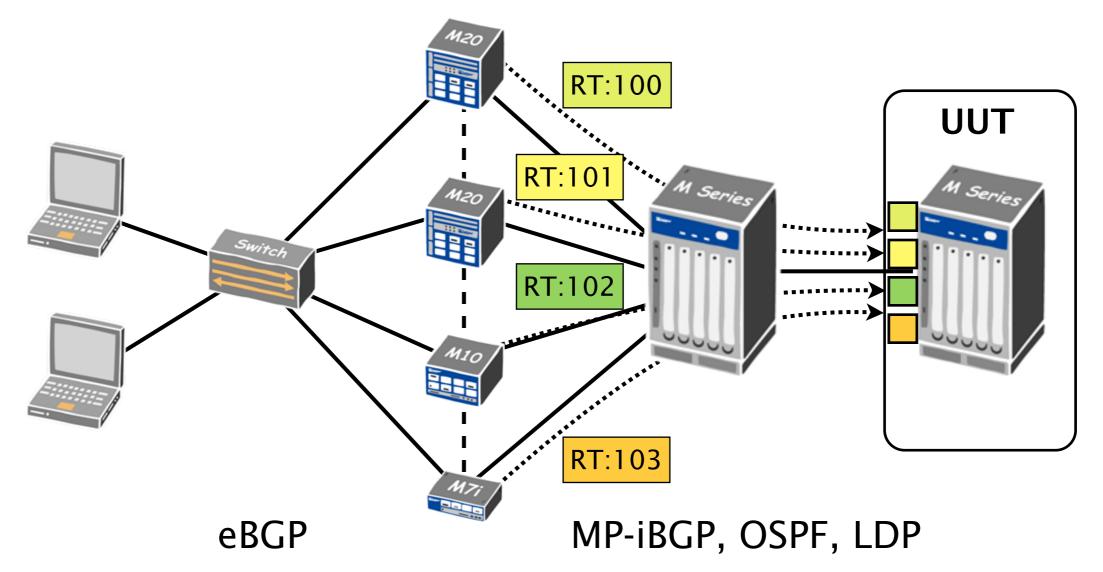
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Is CRS pushing a M120 to the limit ?



Four services are defined on the Unit Under Test (UUT)

- Each service is fed with one class (one RT)
- In each class, ~300k routes (1 path per route)
- In the end, 1.200.000 routes in **RIB** & **FIB**

Is CRS pushing a M120 to the limit ?

- UUT was a Juniper M120 [JunOS 9.3R2.8]
 - Routing Engine (RE) has 4 GB DRAM
 - Forwarding Engine Boards (FEB) have 512 MB DRAM

	RE	FEB
empty	17%	9%
<i>fully-loaded</i> (1.200.000 routes)	38%	39%

- FIB could handle more than 2.000.000 routes
 - Enough to support a few services without modifications

More services ? scalability and...scalability

- Routes dissemination overhead
 - All PEs receive all VPN routes
- Routes storage overhead
 - RIB
 - Modest performance demand
 - Add more DRAM to support CRS ?
 - FIB
 - CRS's biggest challenge
 - Sharing between the VRFs in the FIB?

How could we improve CRS FIB's scaling: *Selective VRF Download*

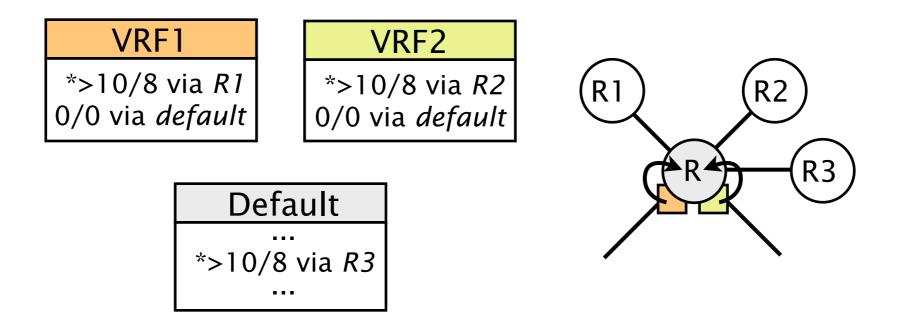
By default, *all* VRFs are installed on *all* line cards

		Temp	CPU Utilization (%)		Memory	Utilization (%)	
Slot	State	(C)	Total	Interrupt	DRAM (MB)	Неар	Buffer
2	Online	24	1	0	512	39	59
3	Online	28	1	0	512	39	59

- Customers ask for the same services ?
 - Connect them on the same line card
 - Download VRFs only to line cards that need them
- It could be a management nightmare...

How could we improve CRS FIB's scaling: *Cross-VRF Lookup*

- Specific routing for a small set of prefixes ?
 - Create one small VRF per service
 - Add default entry towards a default VRF
- The price to pay is 2 IP lookups

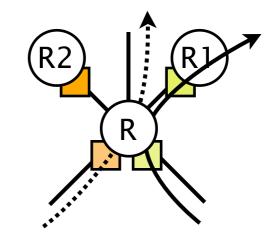


How could we improve CRS FIB's scaling: *Distributed VRF*

- Distribute VRFs among routers which can afford extra load
 - PEs do not maintain complete VRFs anymore
 - PEs default route traffic towards these routers
- Increase in latency and load
- Distributed version of Cross-VRF Lookup

R maintain small VRFs and default rest to R1 or R2 → detour path

→ direct path



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CRS is feasible

- Implementable
 - It can be realized on today's routers
 - It uses well known BGP MPLS/VPNs techniques
- Scalable (for a few services)
 - "Modest" message and storage overhead
 - Lab experiments tend to confirm that
- Guaranteed interdomain convergence
 - Extra flexibility does not compromise global routing stability¹

¹ Proof in SIGMETRICS'09 paper by Y. Wang, M. Schapira, and J. Rexford

Questions ?

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