

# Safe Routing Reconfiguration with Route Redistribution

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# Routing configuration matters

- It controls traffic paths
  - impacts QoE, business goals, SLAs, ...
  - enables resource optimization
- It is optimized wrt changing factors
  - traffic demands, working links, ...

# Live reconfigurations are important

- For traffic engineering
  - fine-tune how traffic flows in the network
- To keep devices up-to-date
  - e.g., new security patches, OS release, equipment renovation, ...
- For evolvability
  - new requirements or services
  - introduction of new protocols
    - e.g., OpenFlow?

# Reconfiguration techniques exists

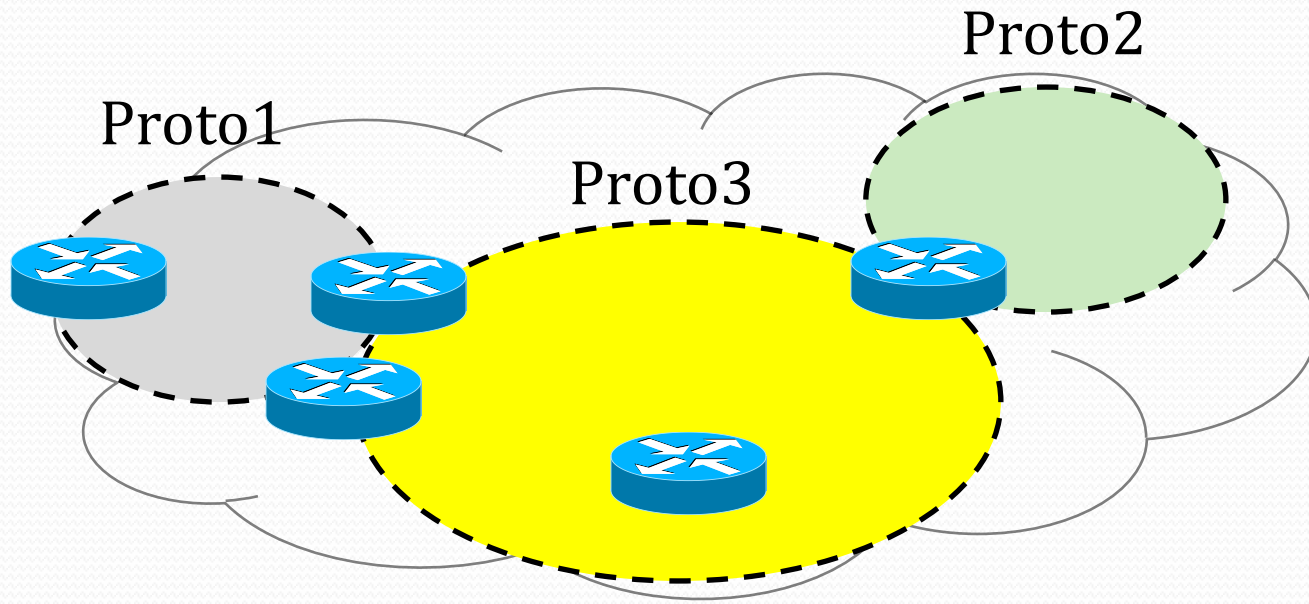
- Industrial guidelines
  - e.g., vendor-based [Herrero10]
- Research proposals
  - case specific, e.g., [Francois07]
  - more general, e.g., Ships-In-The-Night (SITN) [Vanbever11]

# Existing techniques oversimplifies

- Industrial guidelines provide **no guarantees on service continuity**
  - only rules of thumb [Herrero10]
- Research proposals for networks **with a single routing instance**
  - case specific, e.g., [Francois07]
  - more general, e.g., Ships-In-The-Night (SITN) [Vanbever11]

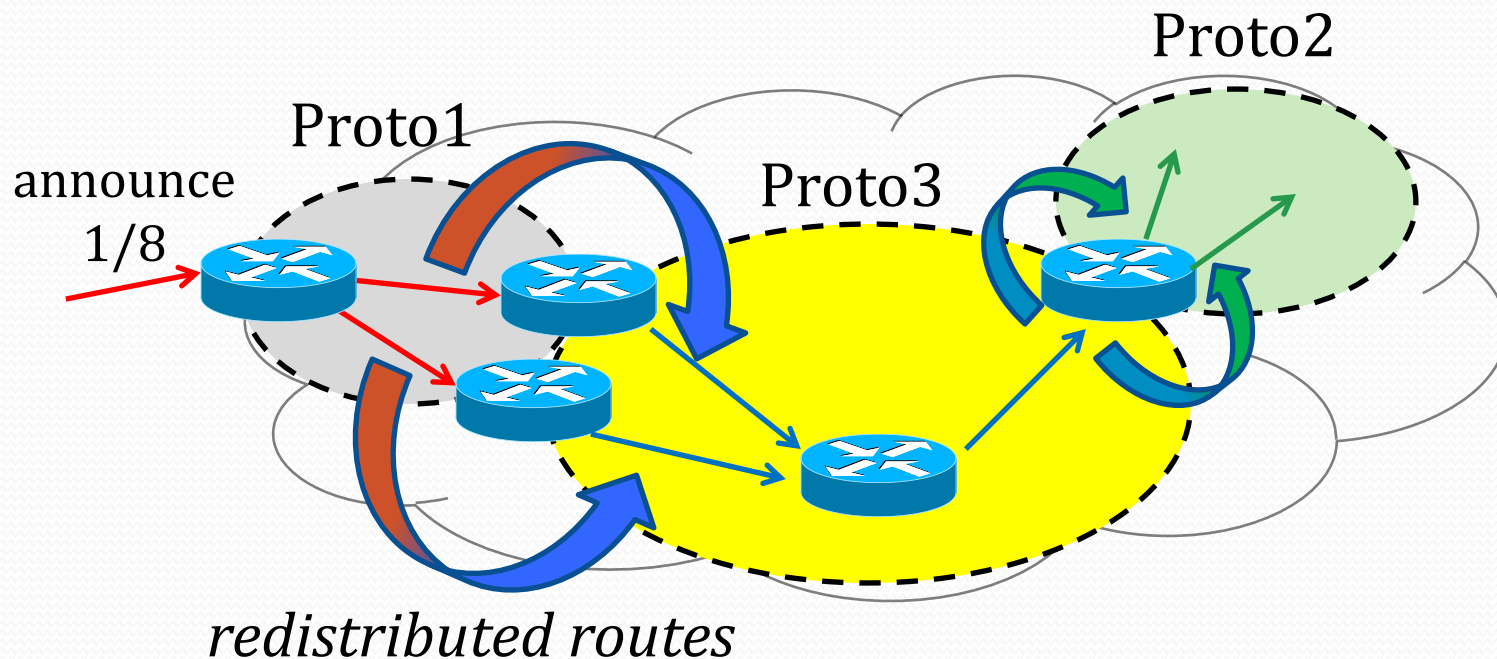
# Real networks are complex

- Multiple *Routing Domains (RD)*
  - each running a different *routing instance* ProtoX



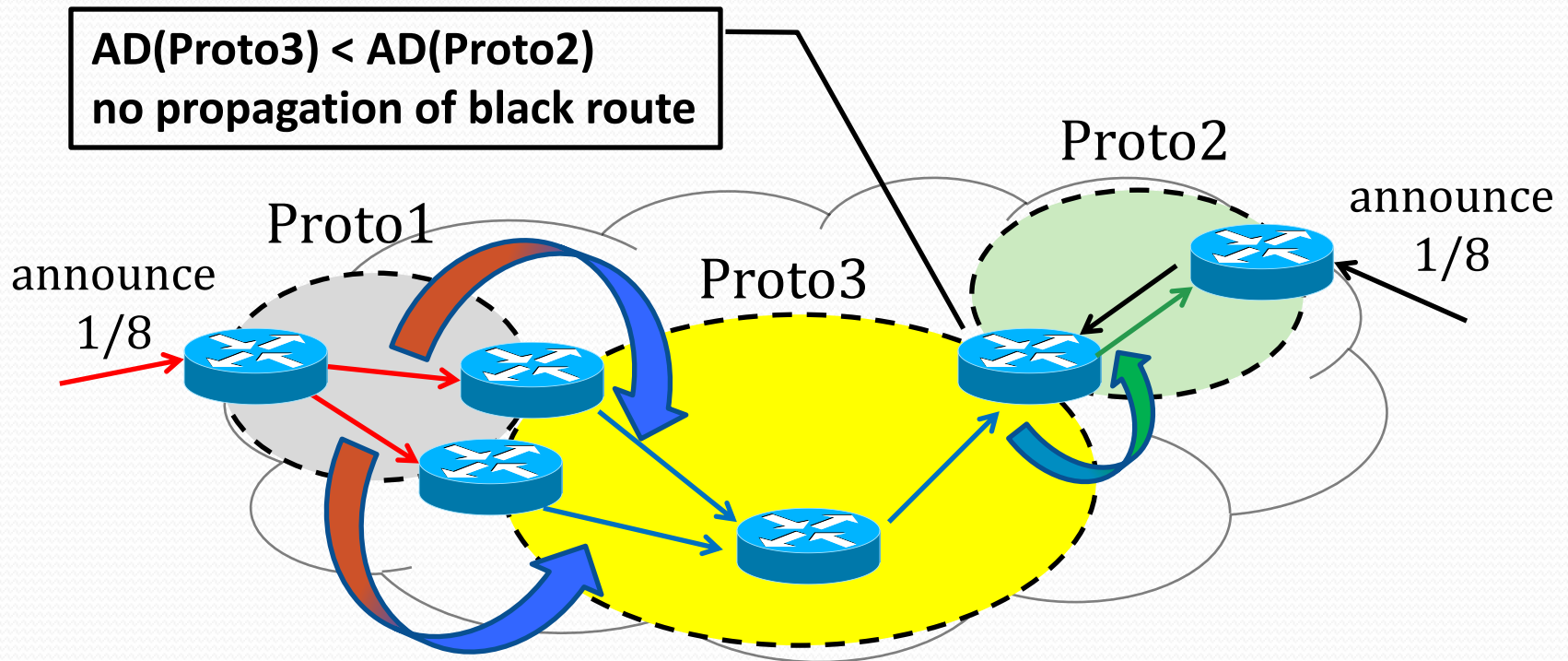
# Real networks are complex

- *Route redistribution (RR)* glues RDs together
  - propagating used routes across RDs



# Real networks are complex

- *Administrative Distance (AD)* encodes instance preferences





# We study practical reconfigurations

- On running networks
  - focus on large enterprises
- Incremental
  - for process control and debug/rollback
- No service disruptions
  - preventing possible routing/fwd anomalies
- No changes to router internals
  - working today

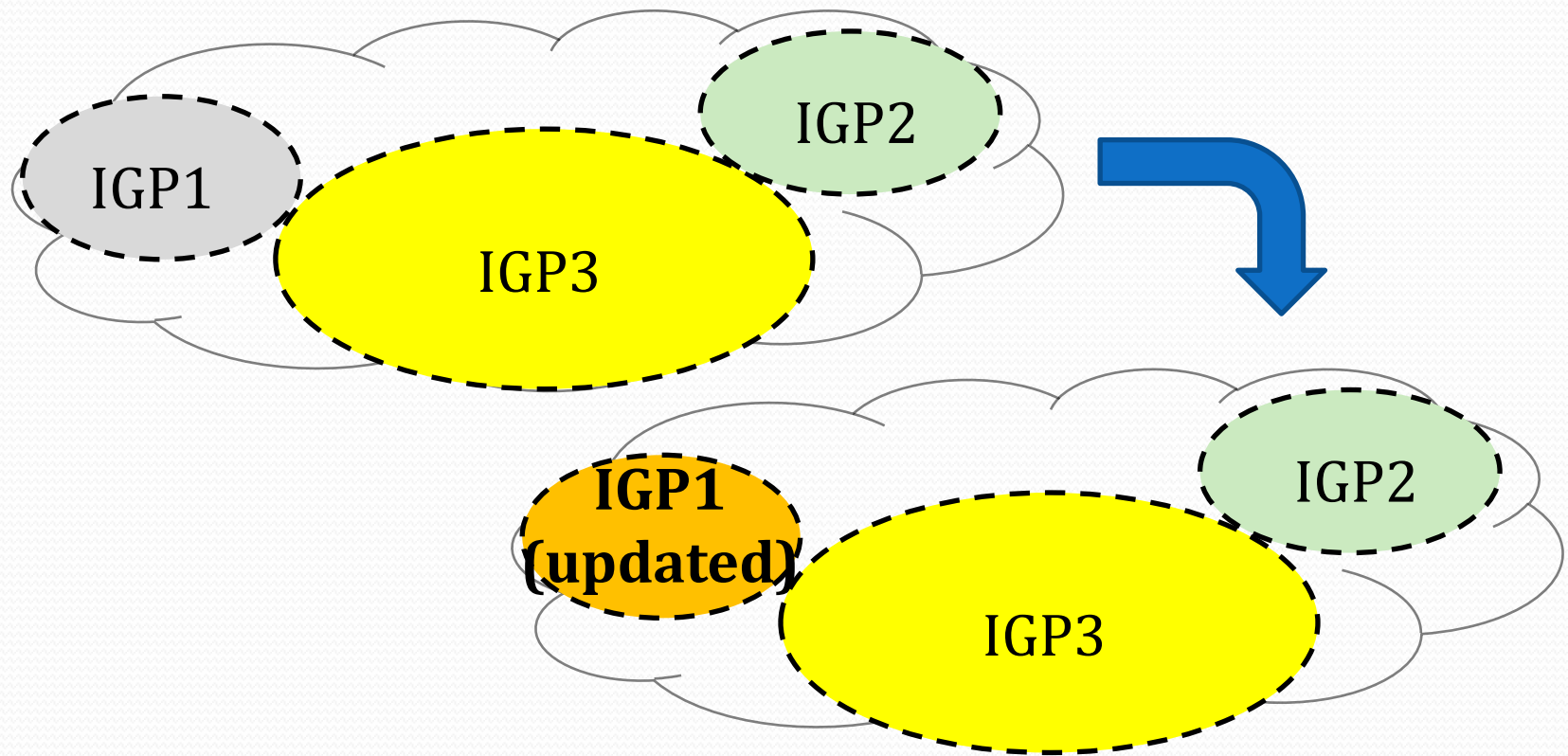
# Contribution overview

- Insight on the general problem
  - single-RD and multi-RD reconfigurations
  - anomalies can and do occur
  - RR affects prior work
- Practical solutions
  - new sufficient conditions for RR
  - provably safe procedures
  - prototype implementation and validation

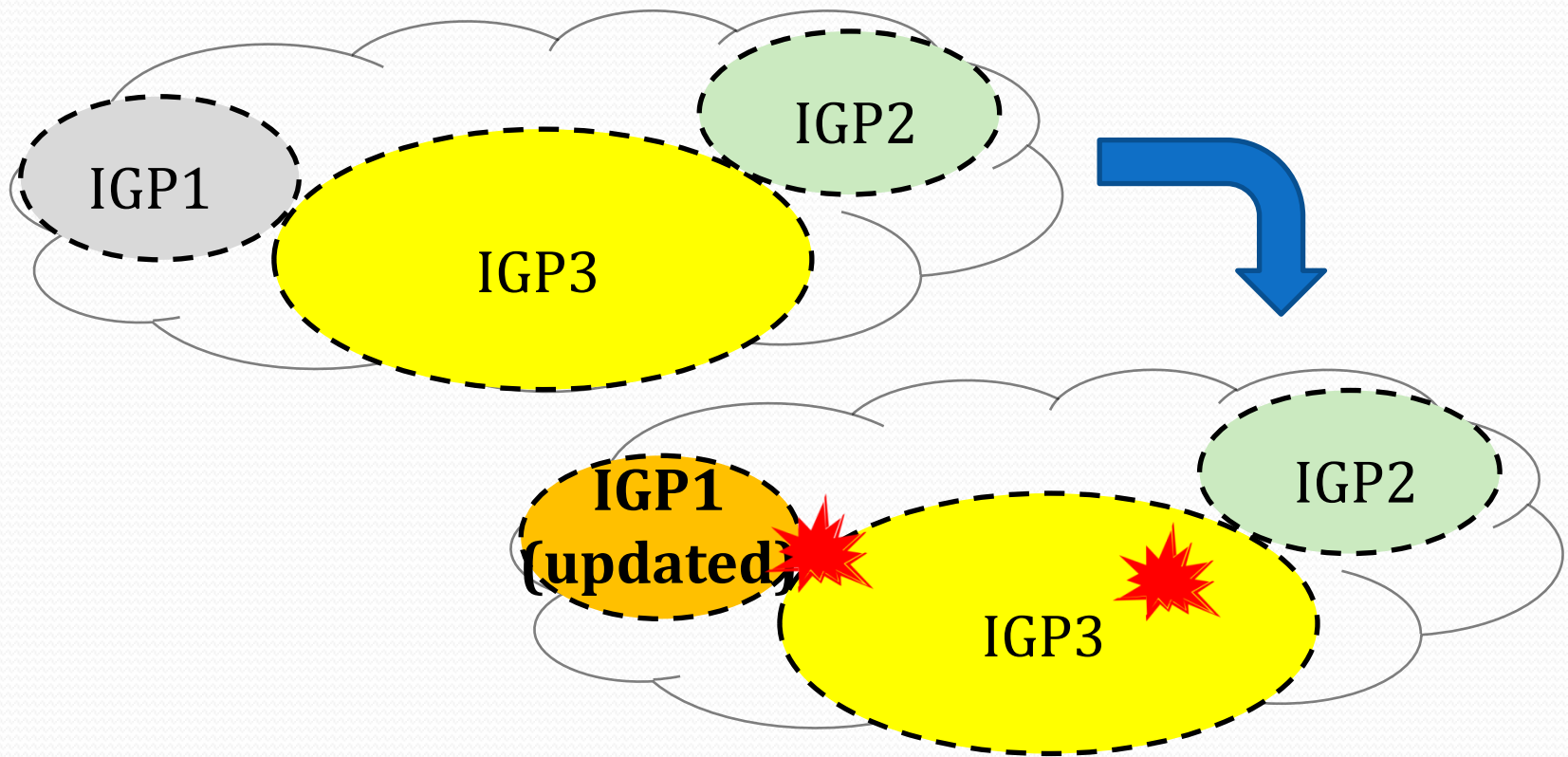
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# Single-RD reconfigurations



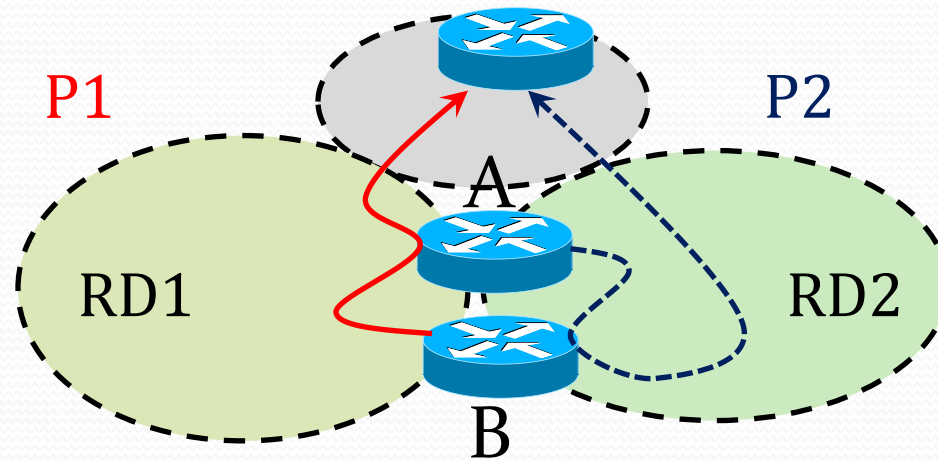
# Existing techniques may not work



# We now refer to SITN

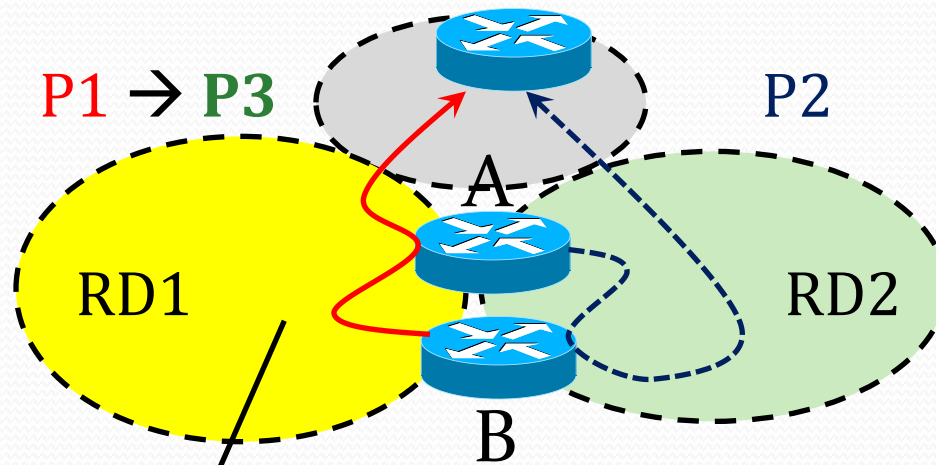
- introduce the final configuration
  - de-preferred
- swaps preference between initial and final configuration
  - on a per-router order
  - in a carefully-computed order

# SITN disruption example



- $AD(P_1) < AD(P_2)$  for both A and B

# SITN disruption example

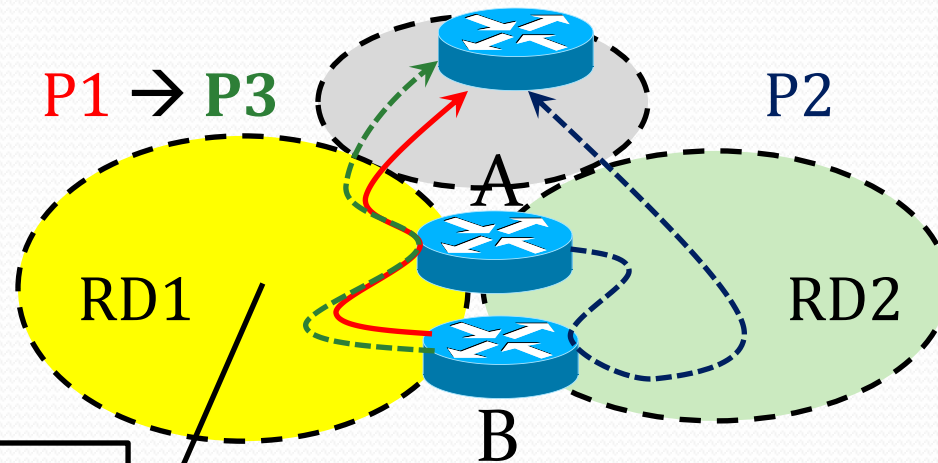


to be reconfigured

- $AD(P_1) < AD(P_2)$  for both A and B



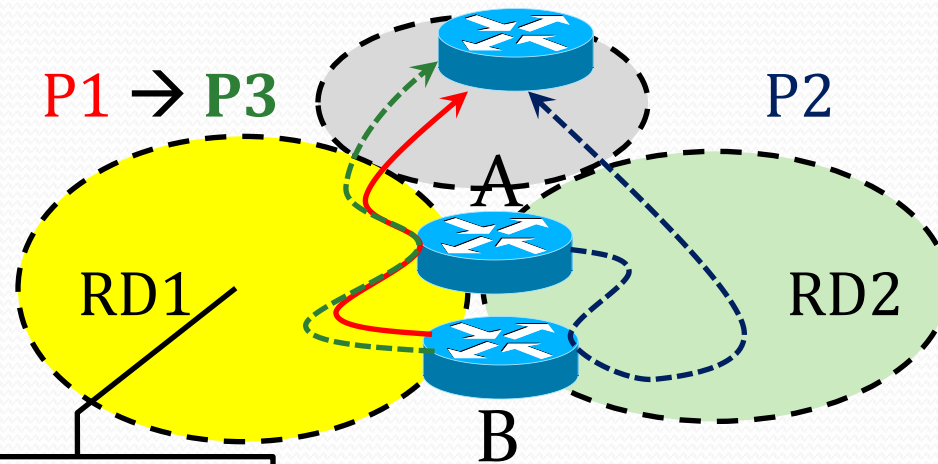
# SITN disruption example



SITN introduces P<sub>3</sub> with higher AD...

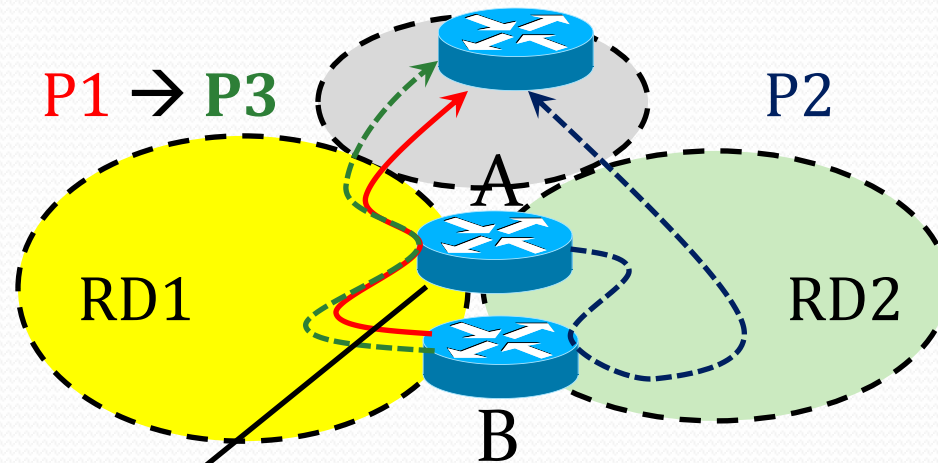
- $AD(P_1) < AD(P_2) < AD(P_3)$  for both A and B

# SITN disruption example



- $AD(P_1) < AD(P_2) < AD(P_3)$  for both A and B

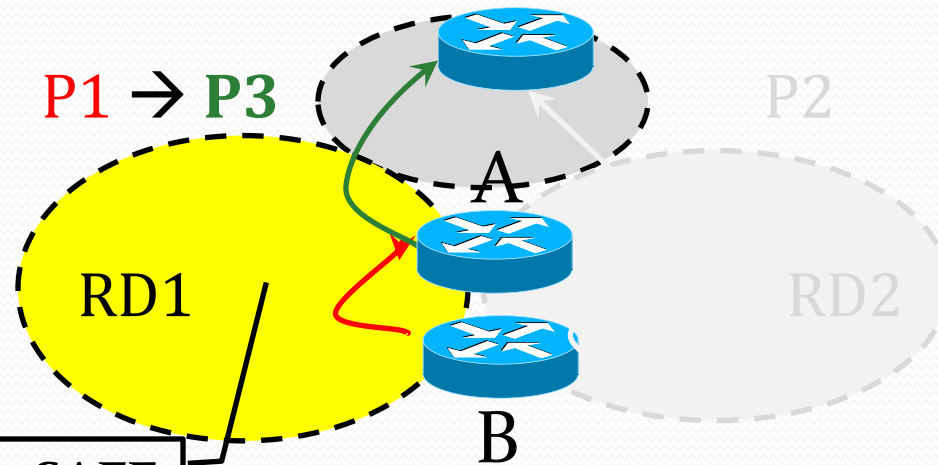
# SITN disruption example



e.g., starting from A

- $AD(P_2) < AD(P_3) < AD(P_1)$  for A
- $AD(P_1) < AD(P_2) < AD(P_3)$  for B

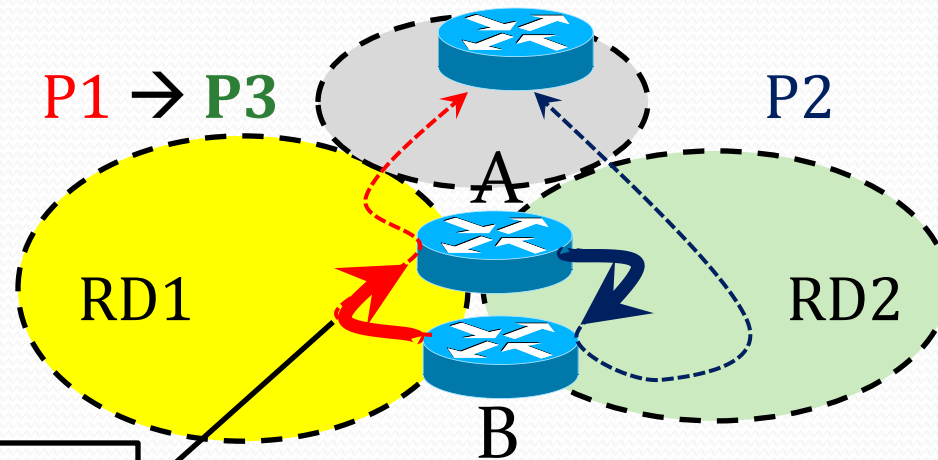
# SITN disruption example



starting from A is SAFE  
in a single-RD network

- $AD(P_2) < AD(P_3) < AD(P_1)$  for A
- $AD(P_1) < AD(P_2) < AD(P_3)$  for B

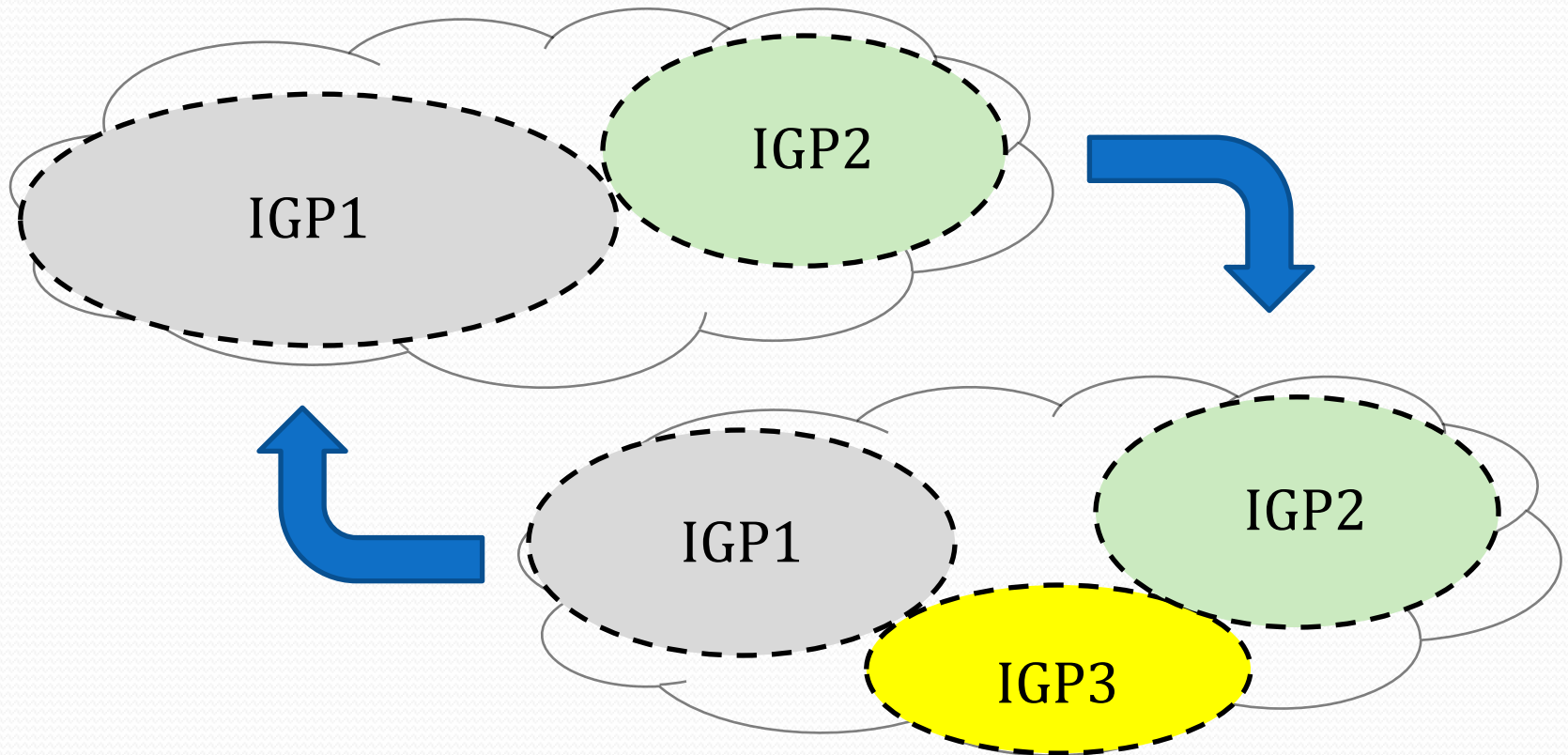
# SITN disruption example



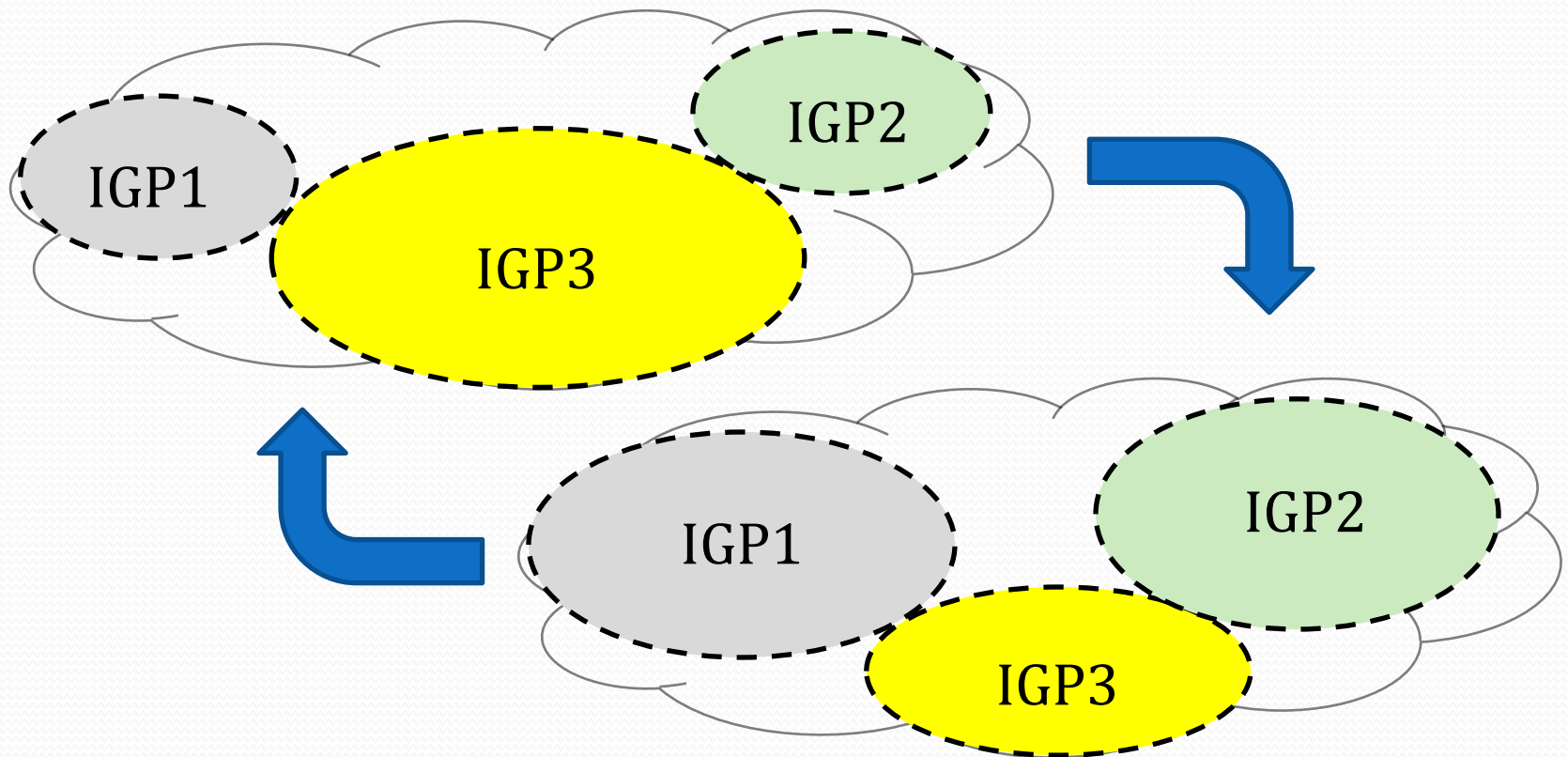
but causes a LOOP  
in the presence of P2

- $AD(P_2) < AD(P_3) < AD(P_1)$  for A
- $AD(P_1) < AD(P_2) < AD(P_3)$  for B

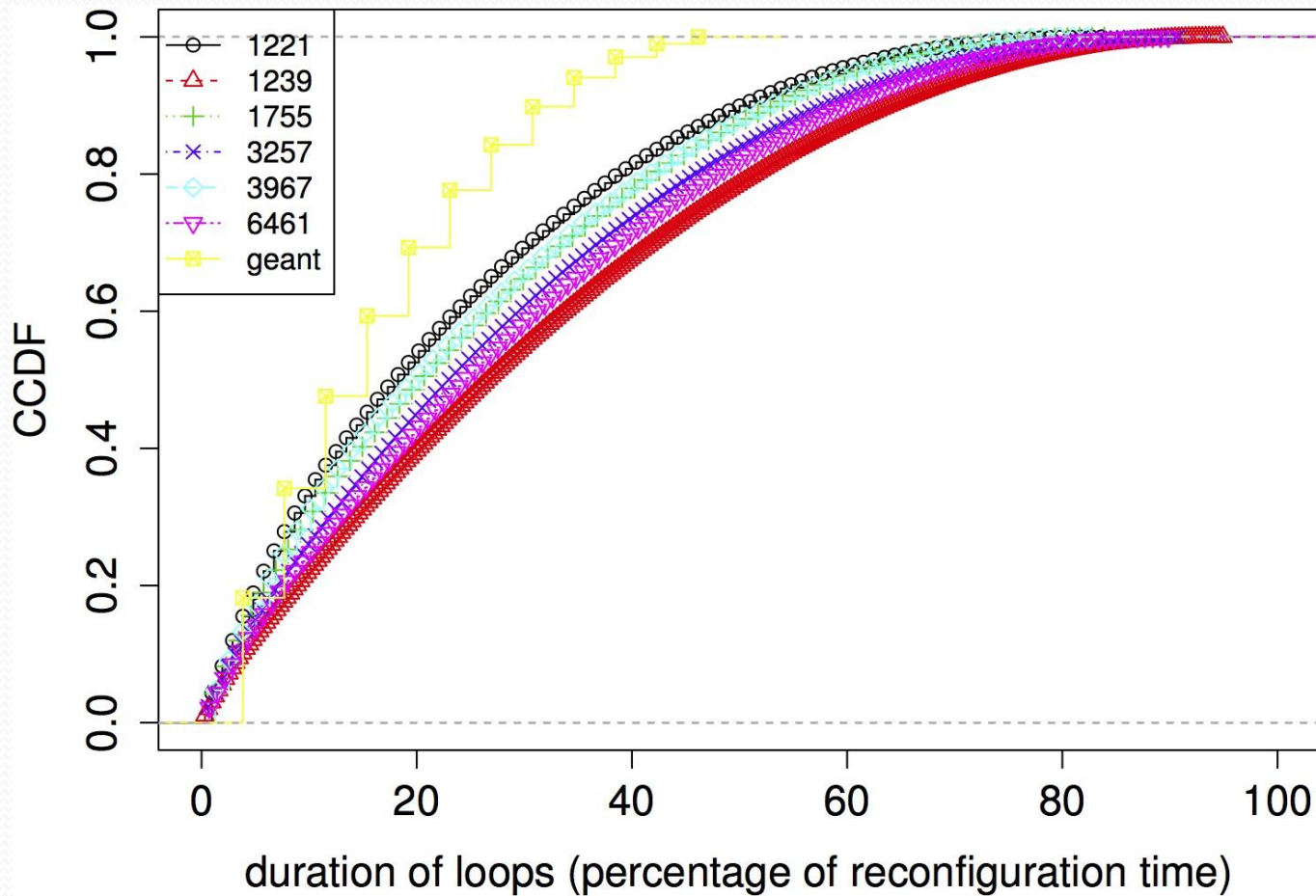
# Multi-RD reconfigurations



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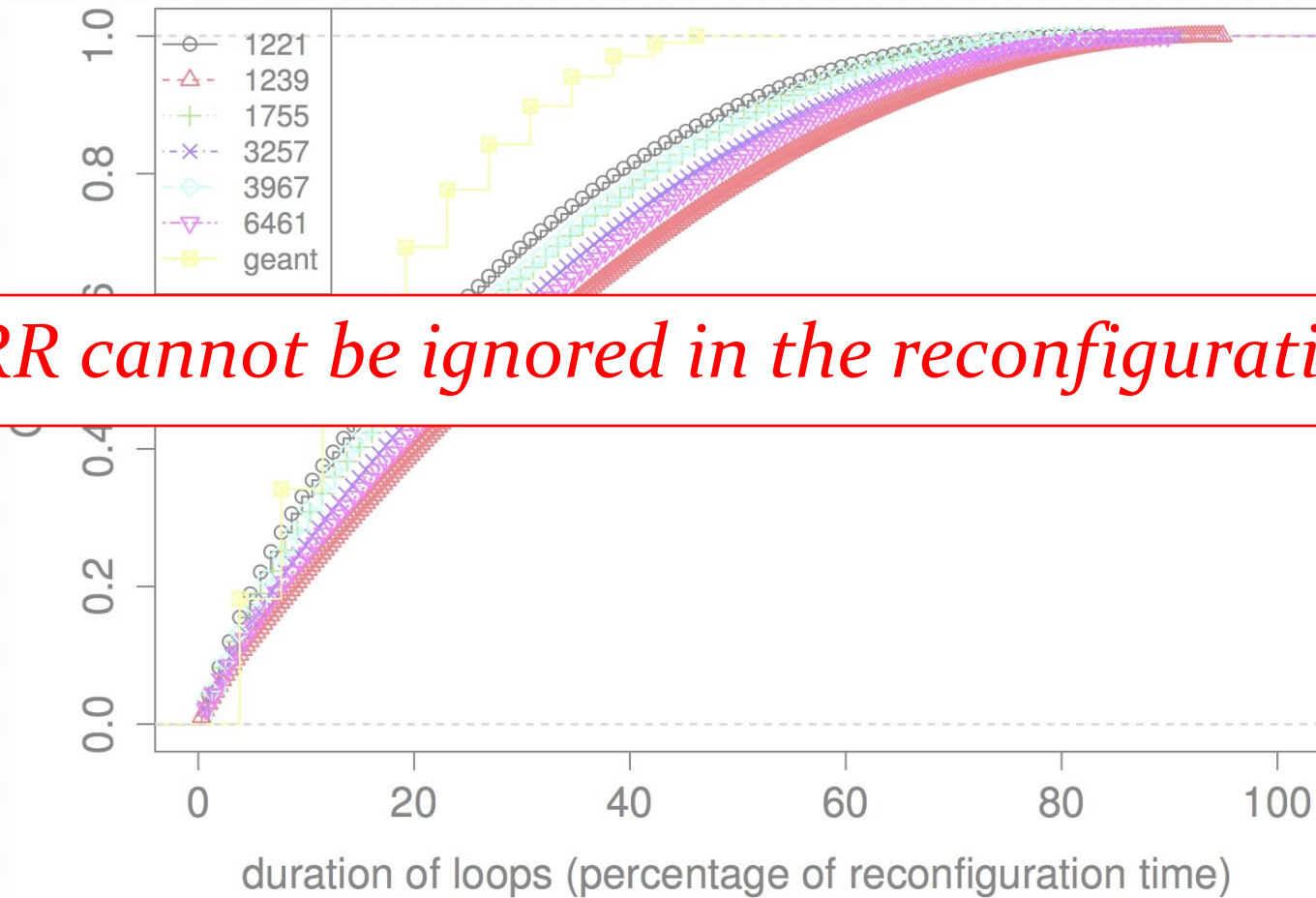


# Long-lasting loops can occur





# Long-lasting loops can occur

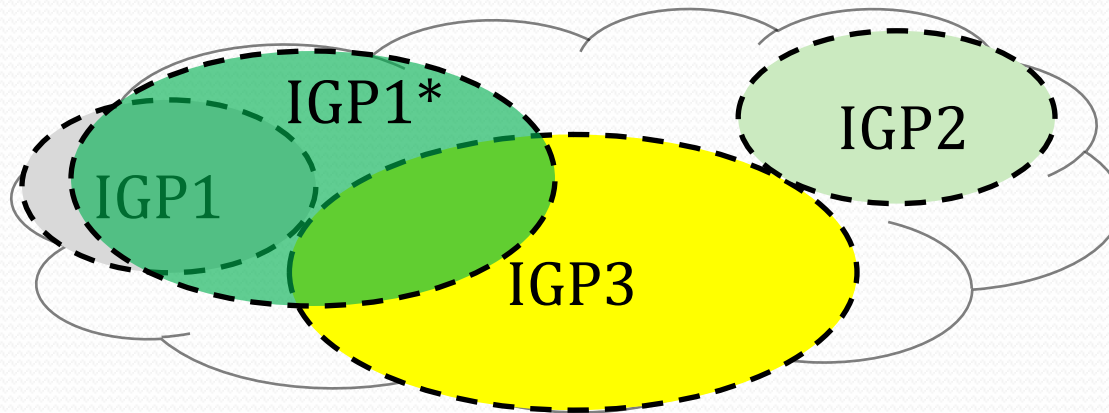


# Known techniques hard to extend

- Additional router-level operations
  - activating/de-activating RR
- Available routes change during the reconfiguration
  - route selection influences RR, and vice versa
  - new routes can be announced
  - existing routes can be withdrawn
- Previous RR theory does not apply
  - assumes one routing instance per RD

# Previous RR theory does not apply

- Multiple instances per RD
  - in single/multi-RD reconfigs
- Nested RDs
  - in multi-RD reconfigs



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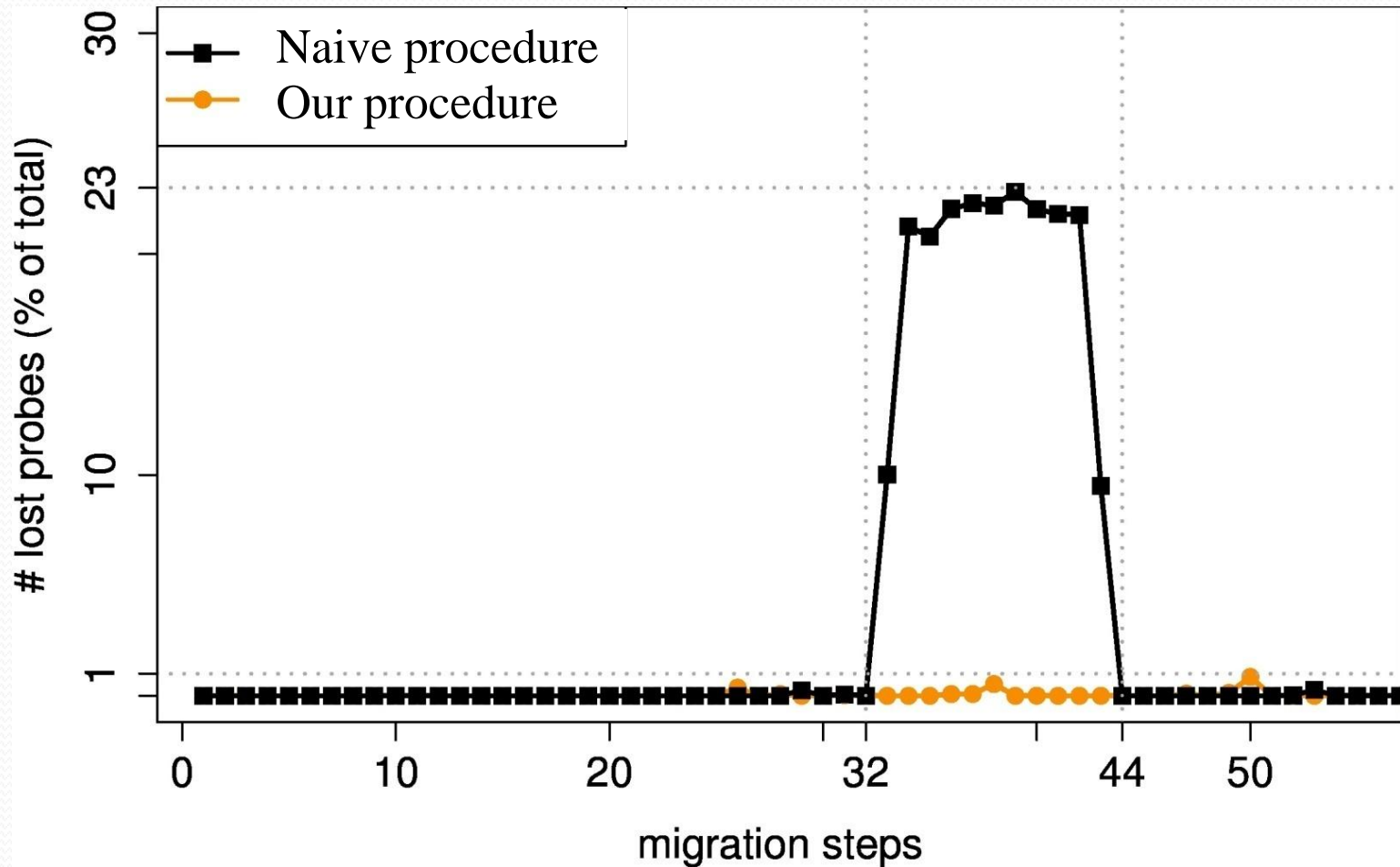
# Safe single-RD reconfigurations

- Looser RR correctness conditions for disjoint RDs
  - each RD is assigned a unique AD interval
  - absence of loops internal to any RD
- Safe procedures
  - extending previous techniques
  - always applicable via AD interval pre-adjustment

# Safe multi-RD reconfigurations

- New RR correctness conditions for nested RDs
  - all routers prefer the innermost RD
  - each shortest paths within any RD crosses at most one RR router
- Safe procedures
  - extending the SITN approach
  - supporting RD splitting/merging/reshaping

# Live safe RD split (on virtual Geant)



# Conclusions

- Study of practical reconfigurations
  - in enterprise networks with multiple RDs
  - overcoming limitations of prior work
- Extended RR theory
  - looser sufficient conditions for RR correctness
  - for both disjoint and nested RDs
- Safe reconfiguration procedures
  - based on our theory extension
  - validated via prototype implementation



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Thanks for attention!

Questions??