



"oFIQUIC: Leveraging QUIC in OSPF for seamless network topology changes"

Rybowski, Nicolas ; Pelsser, Cristel ; Bonaventure, Olivier

ABSTRACT

Link state-routing protocols such as OSPF and ISIS are used in most if not all Internet Service Provider and enterprise networks. They both rely on flooding to distribute the network topology to all routers. Upon topology changes, all routers update their forwarding tables asynchronously which leads to transient events such as micro-loops and packet losses. We propose two improvements to OSPF in an extension called oFIQUIC. First, we use QUIC to exchange routing information between neighboring routers. Second, we revisit the OSPF flooding process. Instead of relying entirely on flooding to distribute topology changes, we establish secure remote QUIC sessions with distant OSPF routers to inform them of topology changes. This enables oFIQUIC to prevent transient loops by ordering the updates of the forwarding tables of all routers after a topology change. We add oFIQUIC to the BIRD implementation of OSPF. Our evaluation demonstrates that oFIQUIC prevents loops and converges quickly in different topologies.

CITE THIS VERSION

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oFIQUIC: Leveraging QUIC in OSPF for seamless network topology changes

Nicolas Rybowski, Cristel Pelsser, Olivier Bonaventure
firstname.lastname@uclouvain.be



Agenda

- **Topology changes in IGP**s
- Ordered FIB updates (oFIB)
- oFIQUIC design
- oFIQUIC performances

Seamless topology changes is necessary

Topology changes in IGPs

Planned modification

Unplanned modification

Seamless topology changes is necessary

Topology changes in IGP

Planned modification

Unplanned modification

- Adding / removing links / nodes
 - Physically upgrading backbone network
 - Logically reconfiguring optical underlay
 - e.g. add virtual link, disable link, modify available bandwidth

Seamless topology changes is necessary

Topology changes in IGP

Planned modification

- Adding / removing links / nodes
 - Physically upgrading backbone network
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 - e.g. add virtual link, disable link, modify available bandwidth
- Modifying links metric
 - Reflect link bandwidth
 - Traffic engineering

Unplanned modification

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Unplanned modification

- Link failures
 - e.g. Fiber cut, failed NIC
- Node failure

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Topology changes in IGPs

Planned modification

- Adding / removing links / nodes
 - Physically upgrading backbone network
 - Logical changes

- Modifying link weights

- Reflecting changes
- Traffic engineering

Unplanned modification

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 - e.g. Fiber cut, failed NIC
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Table 2

About topological changes, a spatial aggregation attempt of listener events (per day statistics).

Types	Average	Median	Max	Total
Link down/up	37.35	6	2074	16,848
Router down/up	0.3	0	8	136
Weight change	0.12	0	4	57

Seamless topology changes is necessary

Topology changes in IGP

Planned modification

- Adding / removing links / nodes
 - Physically upgrading backbone network
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- Modifying

- Reflection
- Traffic

Unplanned modification

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Table 2

About topological changes, a spatial aggregation attempt of listener events (per day statistics).

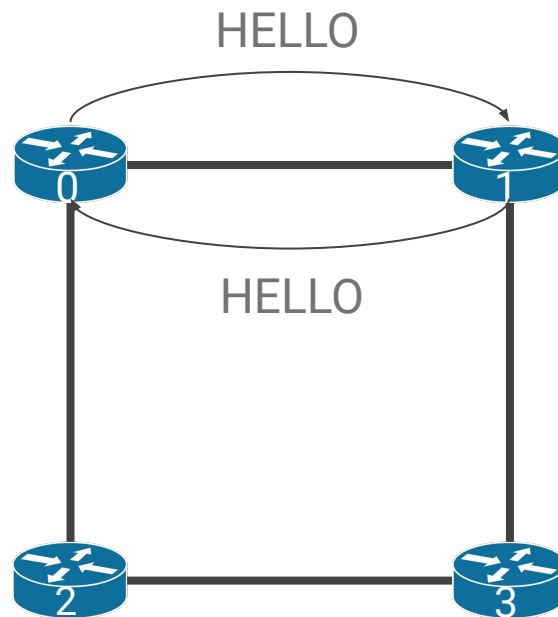
Types	Average	Median	Max	Total
Link down/up	37.35	6	2074	16,848
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Issue: IGP not designed for frequent changes

OSPF discovers neighbors

IGP remainder

1. Neighbor discovery
(2-way connectivity check)

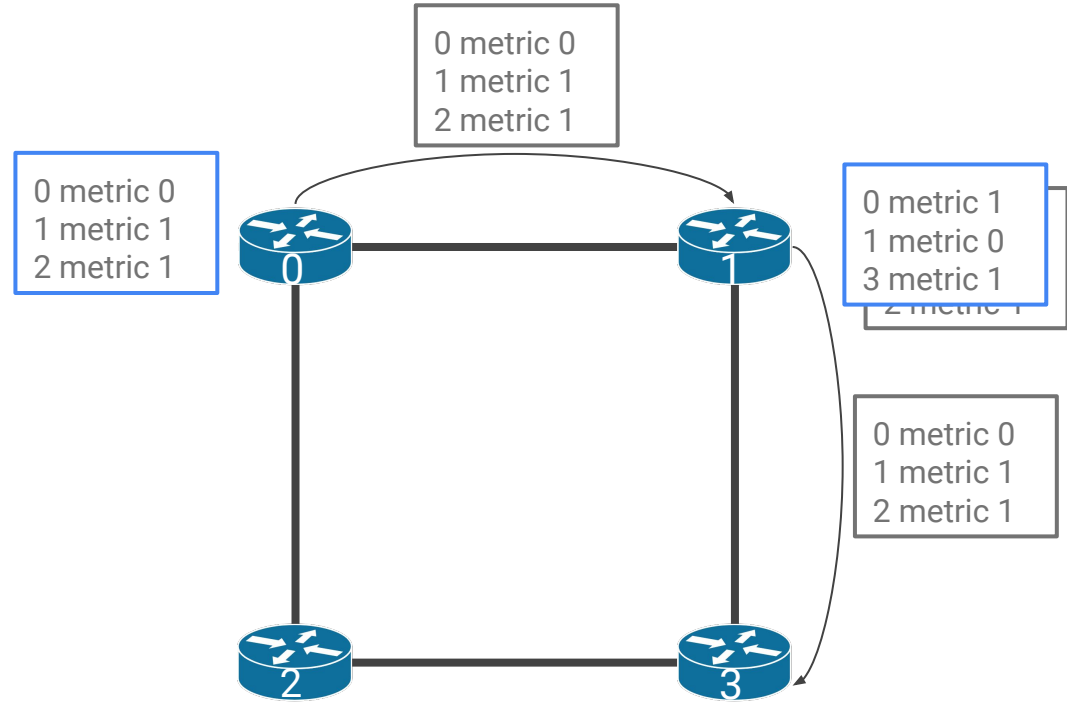


All metrics are unitary

OSPF floods Link State Advertisements (LSAs)

IGP remainder

1. Neighbor discovery
(2-way connectivity check)
2. Flood Link State
Advertisement (LSAs)

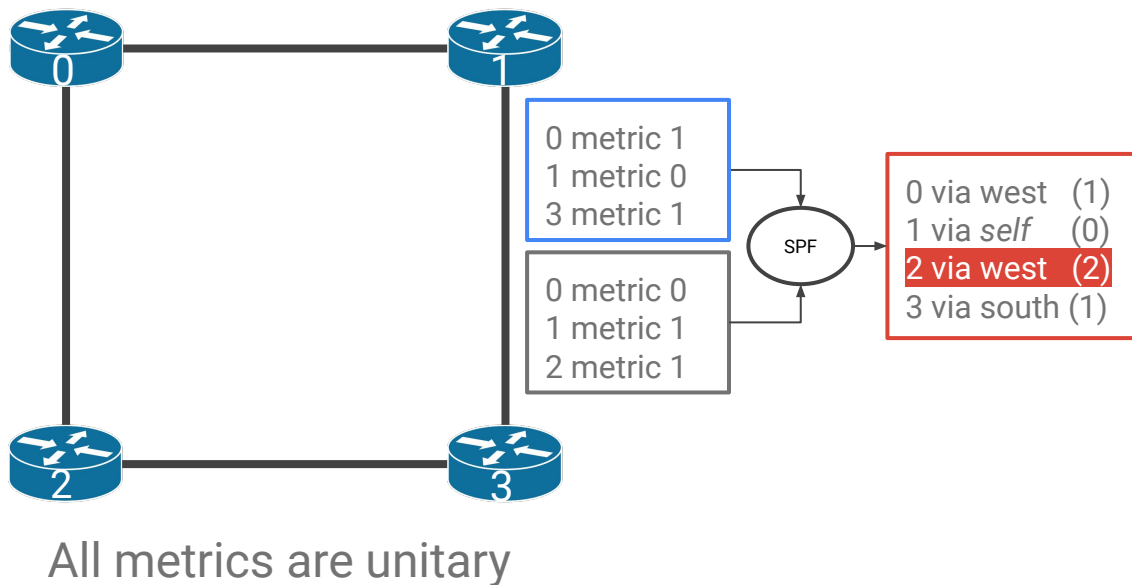


All metrics are unitary

OSPF computes the shortest paths

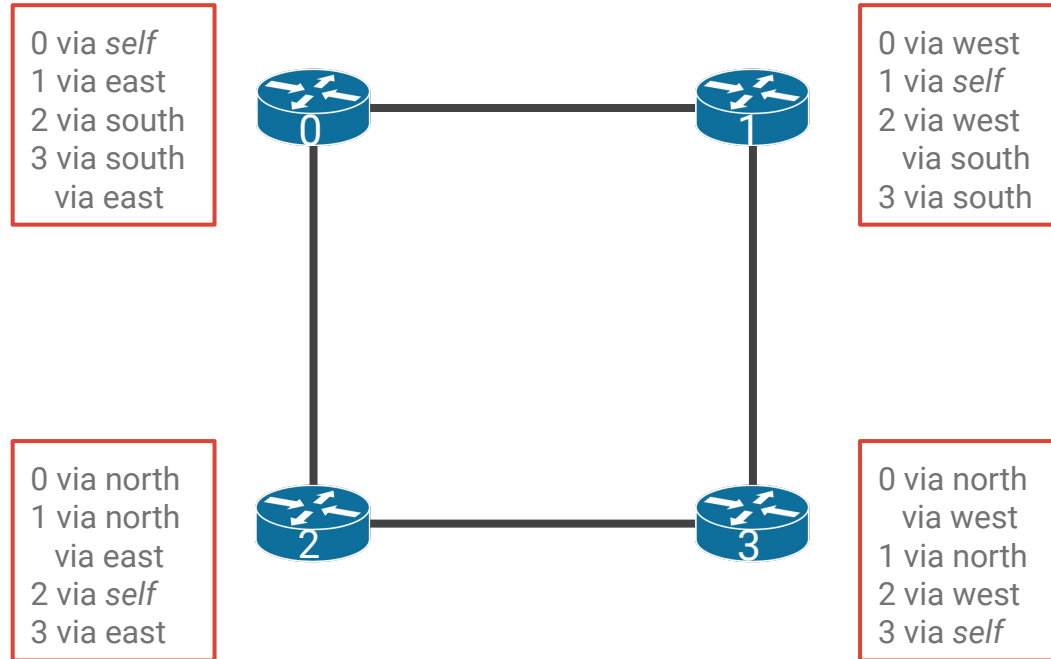
IGP remainder

1. Neighbor discovery
(2-way connectivity check)
2. Flood Link State
Advertisement (LSAs)
3. Shortest paths computation



OSPF converges

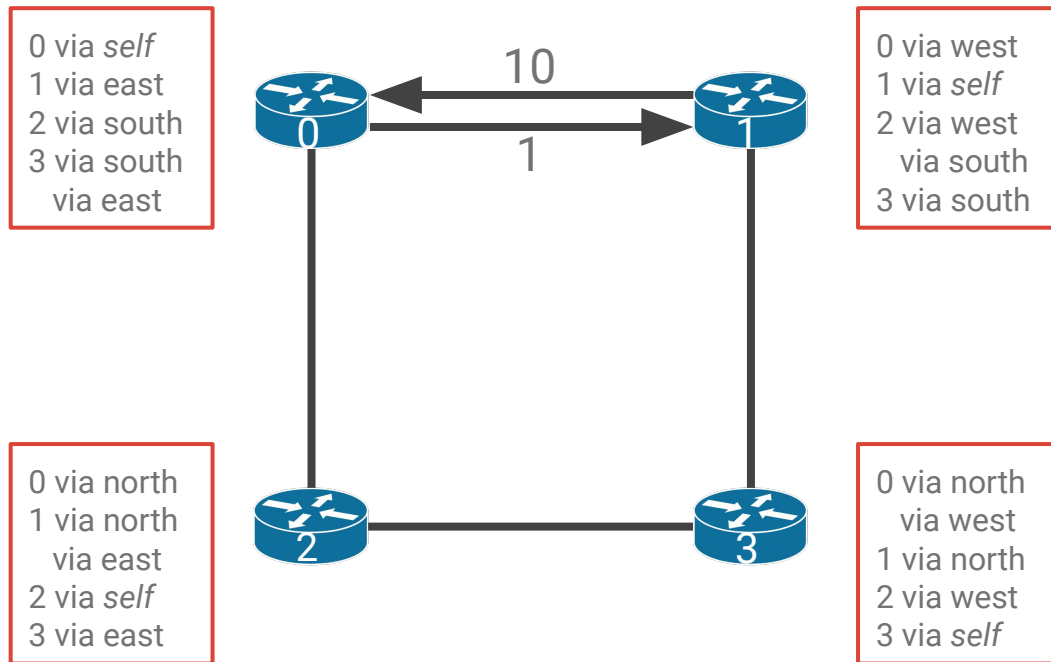
IGP remainder



Increasing link metric

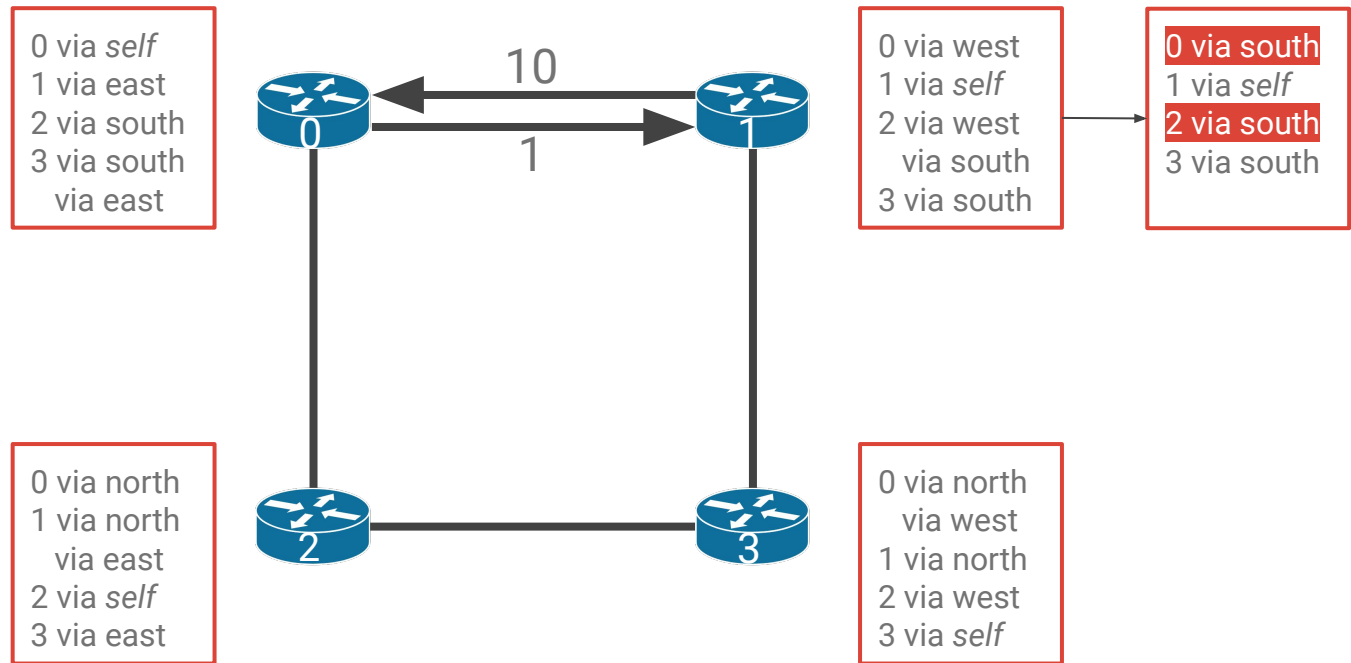
Topology changes in IGP

Virtually remove a link with max metric - 1



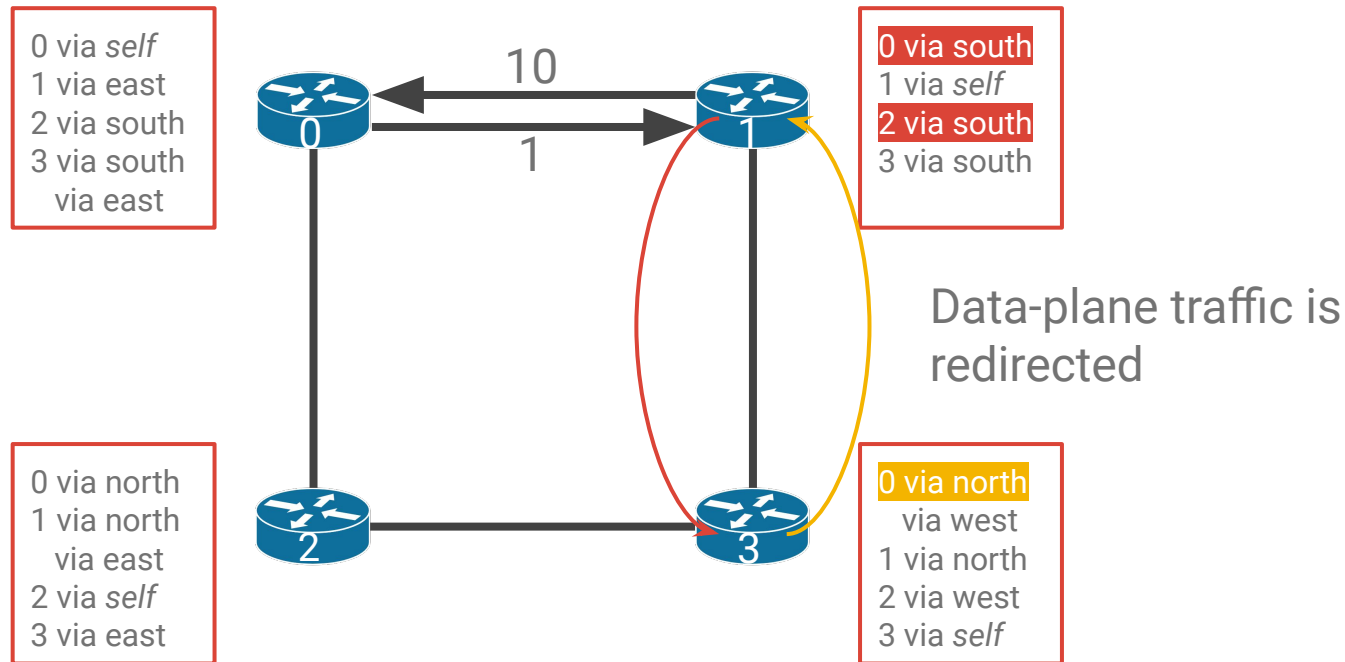
Increasing link metric

Topology changes in IGP



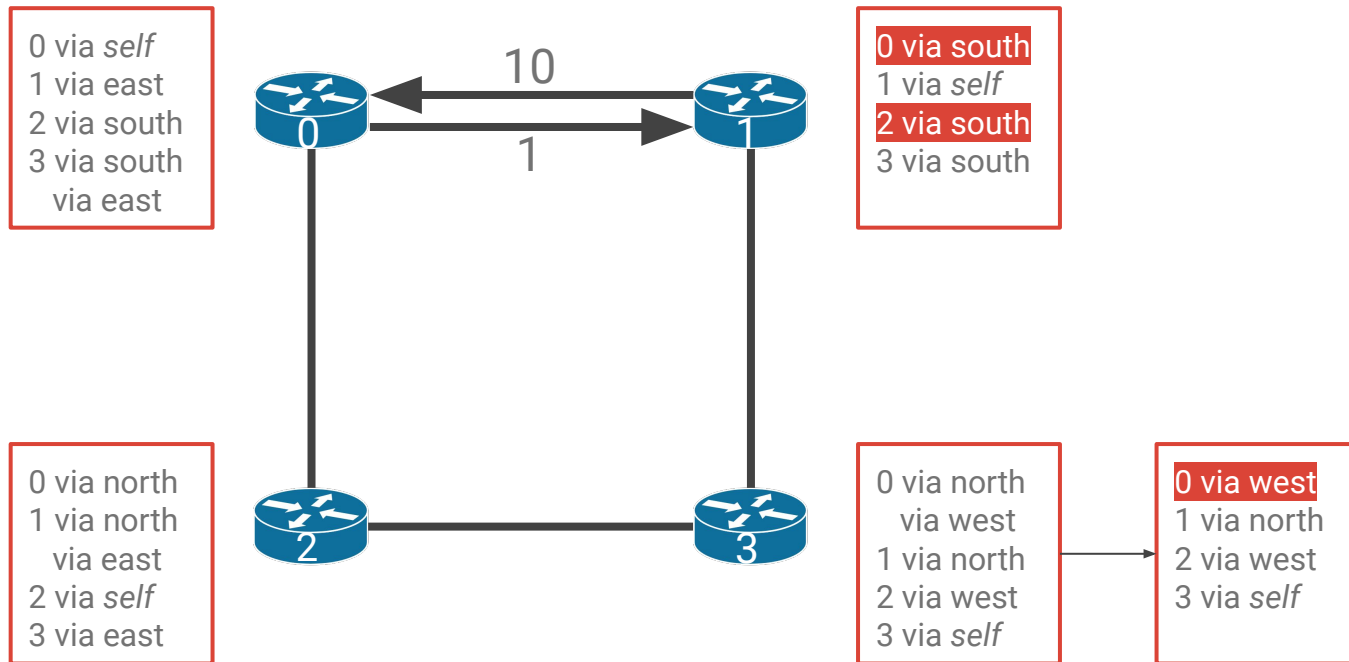
Increasing link metric triggers micro-loops

Topology changes in IGP



Increasing link metric triggers micro-loops

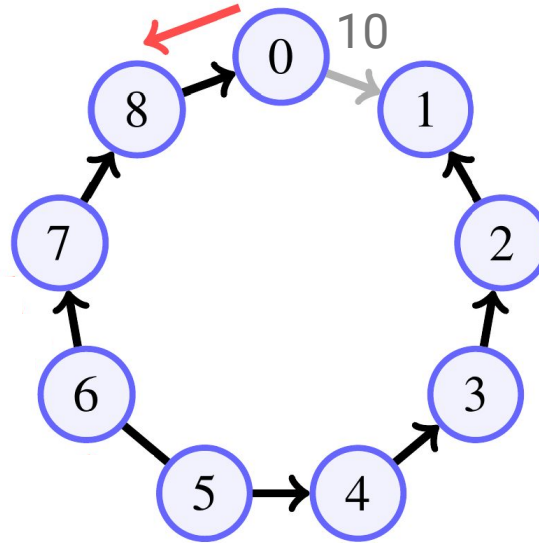
Topology changes in IGP



Micro-loops propagate

Topology changes in IGP

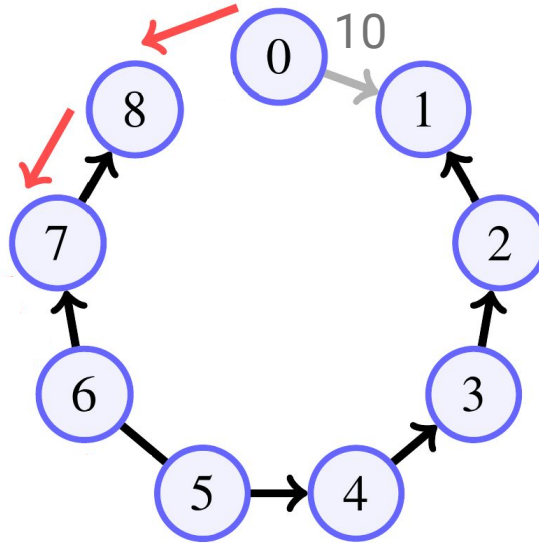
- Unitary link metric



Micro-loops propagate

Topology changes in IGP

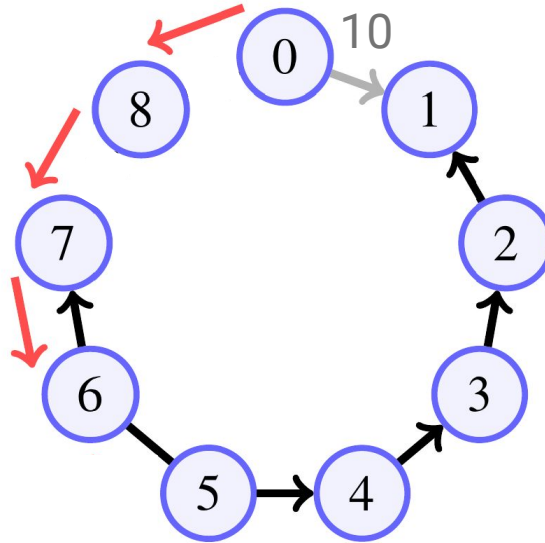
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Micro-loops propagate

Topology changes in IGP

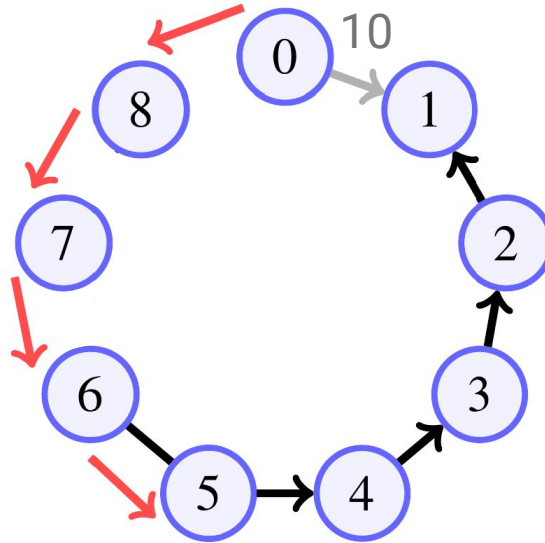
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Micro-loops propagate

Topology changes in IGP

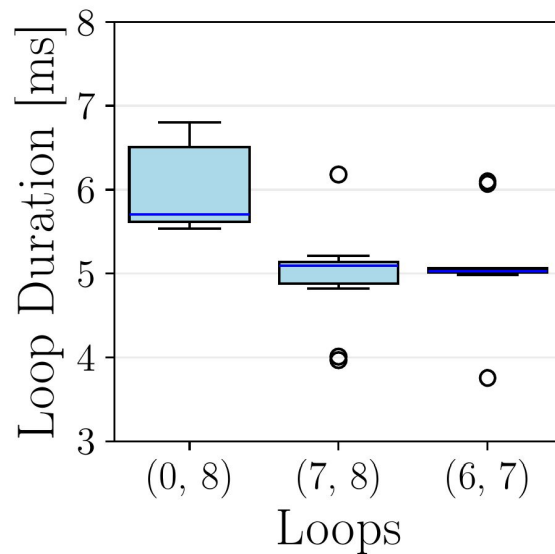
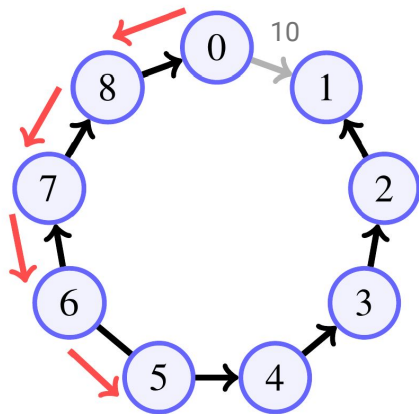
- Unitary link metric



Measuring micro-loops duration

Topology changes in IGP

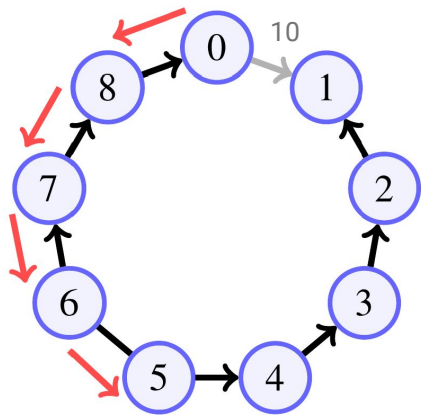
- Unitary link metric
- 5ms link delay



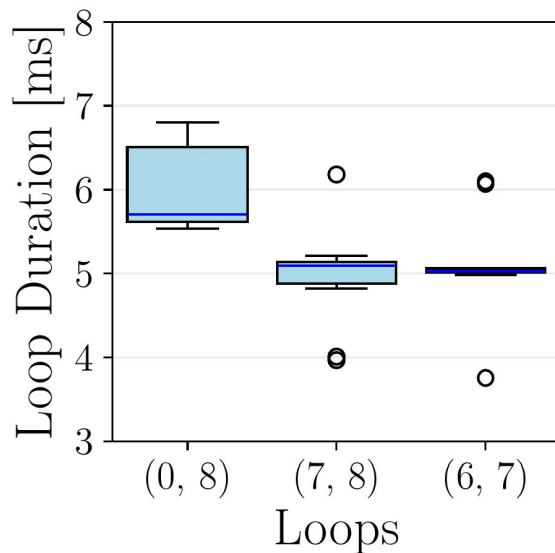
Measuring micro-loops duration

Topology changes in IGP

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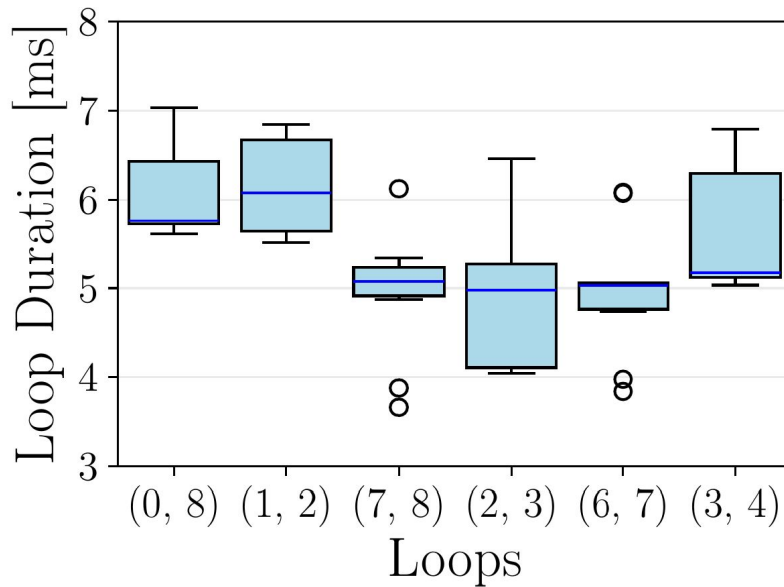
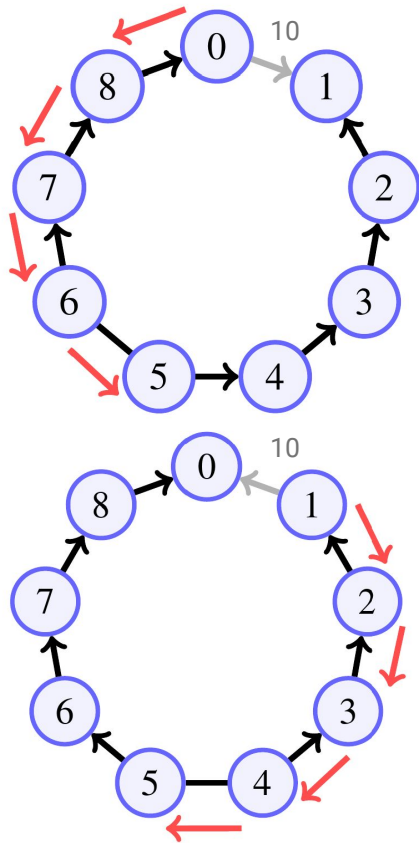
- Micro-loop duration factors
 - Link delay
 - # prefixes
 - SPF computation
 - FIB injection duration (injection rate)
 - ...



Symmetric update increases the number of loops

Topology changes in IGP

- Unitary link metric
- 5ms link delay



Impact of micro-loops

Topology changes in IGP

- Packet reordering

Impact of micro-loops

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- TTL reaches 0 if micro-loop lasts long enough => transient black-hole

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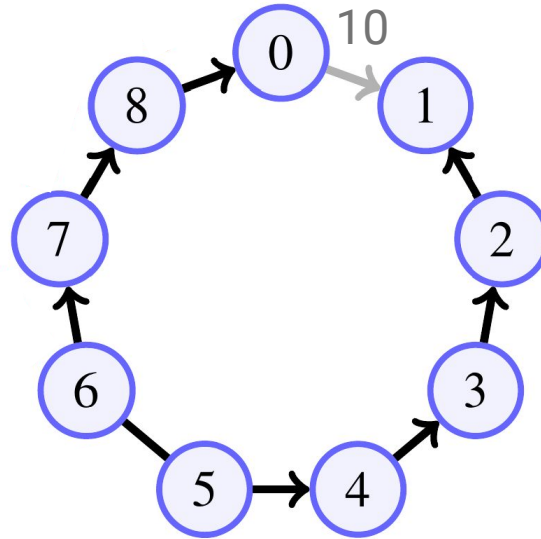
IGPs trigger traffic disruption by design

Agenda

- Topology changes in IGPs
- **Ordered FIB updates (oFIB)**
- oFIQUIC design
- oFIQUIC performances

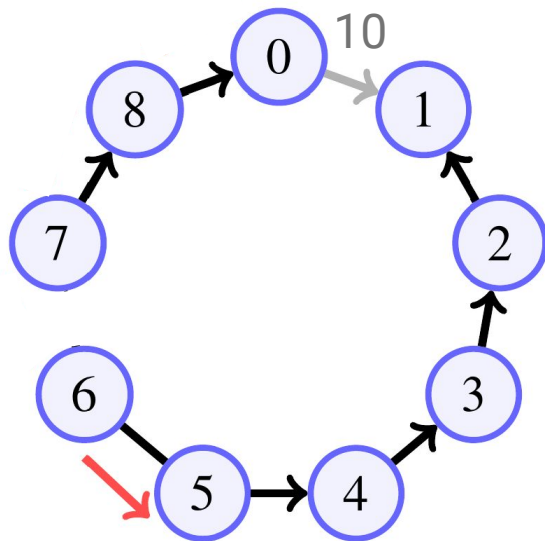
Ordering FIB updates suppresses micro-loops

Ordered FIB updates (oFIB)



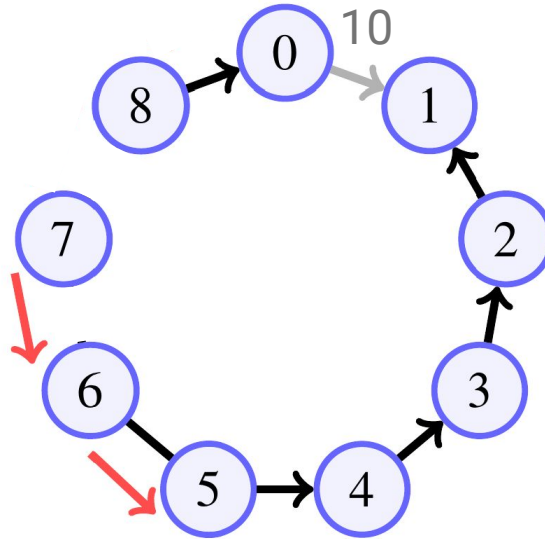
Ordering FIB updates suppresses micro-loops

Ordered FIB updates (oFIB)



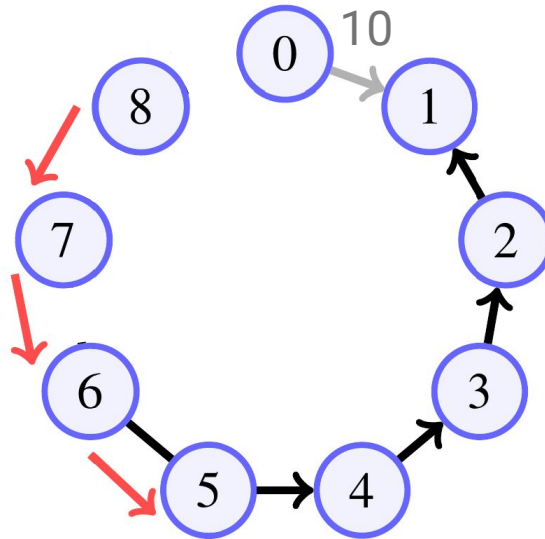
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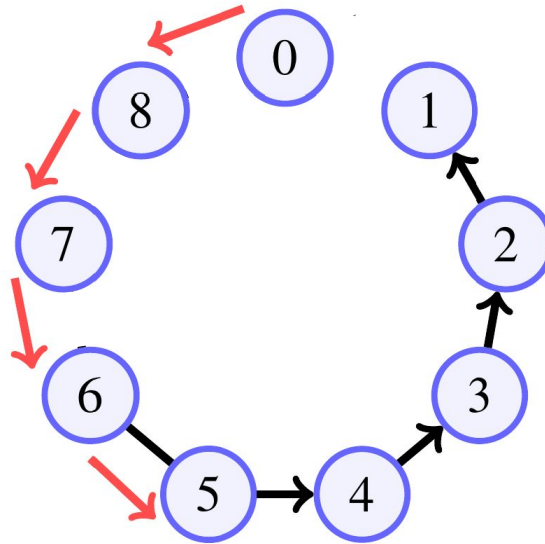
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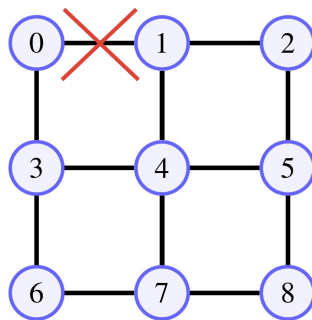
Ordering FIB updates suppresses micro-loops

Ordered FIB updates (oFIB)



Defining FIB update ordering

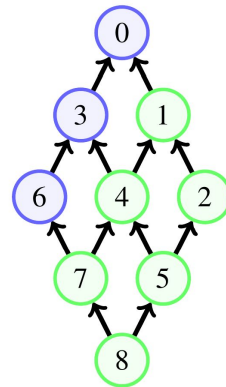
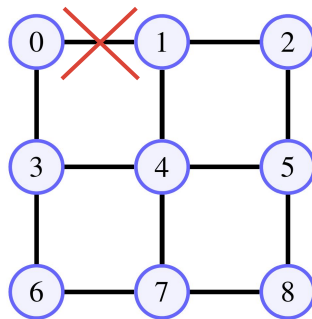
Ordered FIB updates (oFIB)



Defining FIB update ordering

Ordered FIB updates (oFIB)

i. Compute Reverse Shortest Path Directed Acyclic Graph (rspDAG) rooted at R0 [1]

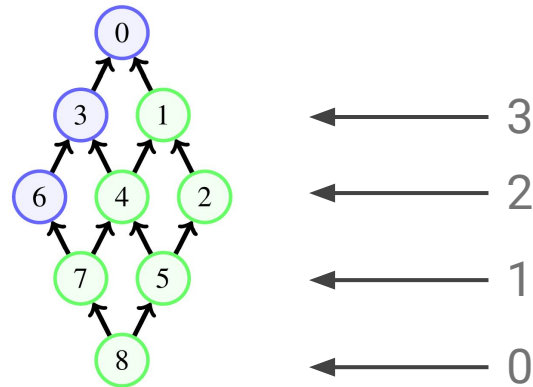
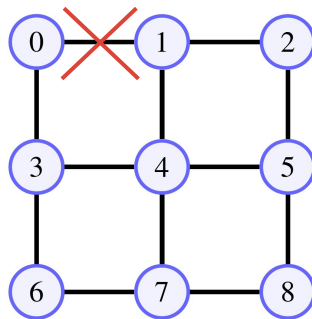


$rspDAG_{1 \rightarrow 0}(0)$

Defining FIB update ordering

Ordered FIB updates (oFIB)

ii. Attribute a rank to each node in the sub-tree [1]



$rspDAG_{1 \rightarrow 0}(0)$

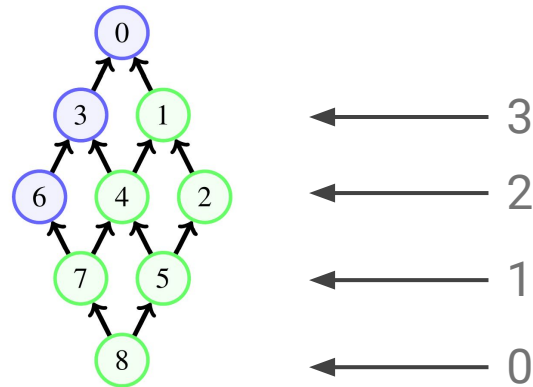
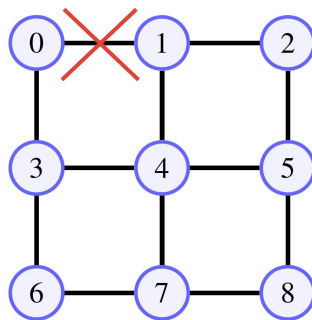
Defining FIB update ordering

Ordered FIB updates (oFIB)

iii. Associate a timer to each node in the sub-tree [1]

$$T(R) = \text{Rank}(R) \times \text{MAX_FIB_TIME}$$

=> Update FIB at timer expiration



$rspDAG_{1 \rightarrow 0}(0)$

Drawbacks of time-based ordering

Ordered FIB updates (oFIB)

- Worst-case timer delays whole convergence
 - Possible shortcut with completion messages [1]
- Distributed computation
- Asynchronous update

=> not suitable for frequent changes

Agenda

- Topology changes in IGPs
- Ordered FIB updates (oFIB)
- **oFIQUIC design**
- oFIQUIC performances

Connecting remote IGP peers

oFIQUIC design

- Scheduling node
 - Ordering computed once
 - Timer-less synchronous updates

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 - Multi-hop transport layer
 - Authentication (Authorization)

Connecting remote IGP peers

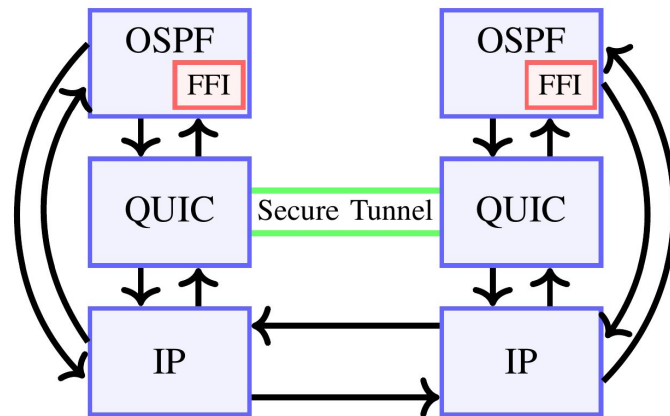
oFIQUIC design

- Scheduling node
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 - Multi-hop transport layer
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- Leverage QUIC transport
 - Complete transport protocol based on UDP
 - Native support for TLS (mutual TLS)

Connecting remote IGP peers

oFIQUIC design

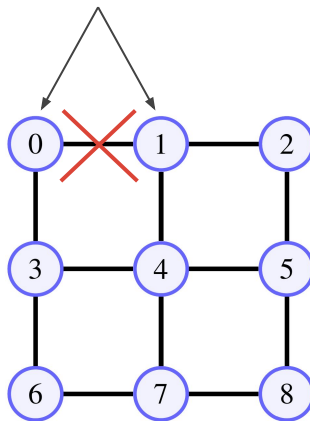
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Ordering FIB updates with oFIQUIC

oFIQUIC design

scheduling nodes

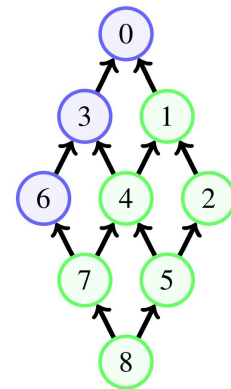
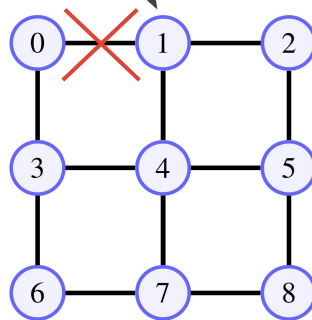


Ordering FIB updates with oFIQUIC

oFIQUIC design

1. Compute rspDAG

scheduling node



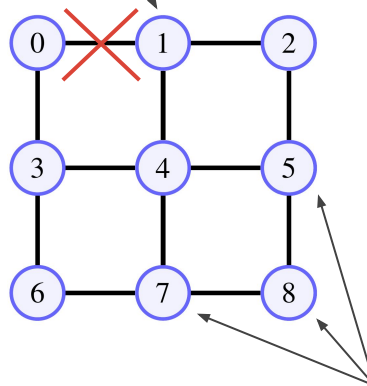
$rspDAG_{1 \rightarrow 0}(0)$

Ordering FIB updates with oFIQUIC

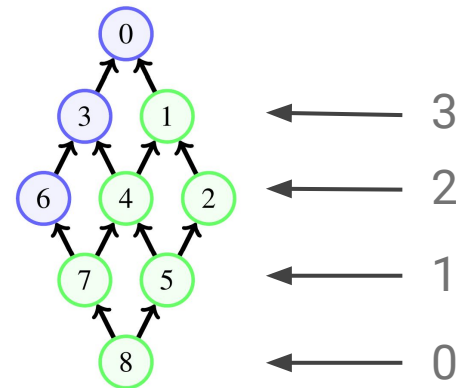
oFIQUIC design

1. Compute $rspDAG$
2. Derive an ordered list of (distant) nodes

scheduling node



distant nodes



$rspDAG_{1 \rightarrow 0}(0)$

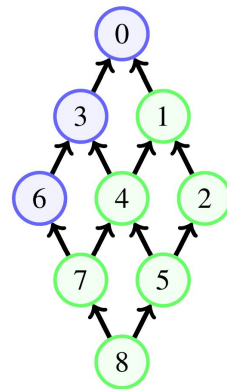
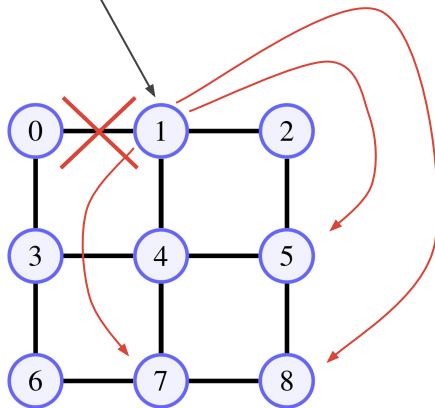
Ordered list: [R8, R5, R7, R2, R4]

Ordering FIB updates with oFIQUIC

oFIQUIC design

1. Compute rspDAG
2. Derive an ordered list of (distant) nodes
3. Initiate QUIC tunnels

scheduling node



$rspDAG_{1 \rightarrow 0}(0)$

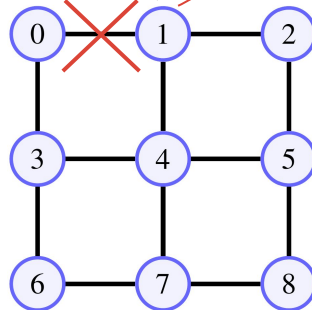
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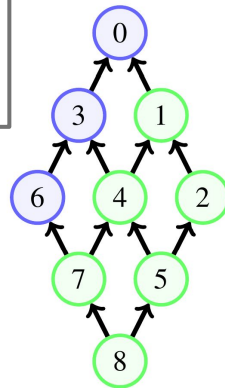
oFIQUIC design

1. Compute rspDAG
2. Derive an ordered list of (distant) nodes
3. Initiate QUIC tunnels
4. Send LSA to first node

scheduling node



0 metric ∞
2 metric 1
4 metric 1



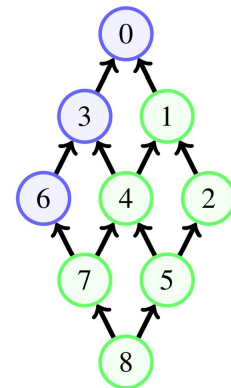
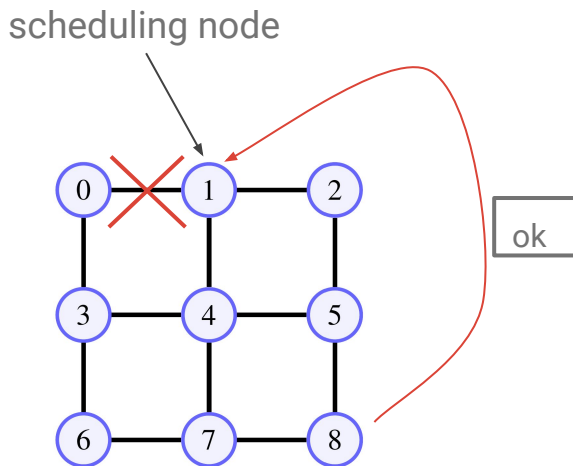
$rspDAG_{1 \rightarrow 0}(0)$

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Ordering FIB updates with oFIQUIC

oFIQUIC design

1. Compute $rspDAG$
2. Derive an ordered list of (distant) nodes
3. Initiate QUIC tunnels
4. Send LSA to first node
5. Wait for distant node confirmation



$rspDAG_{1 \rightarrow 0}(0)$

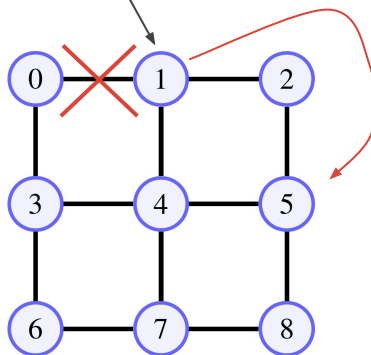
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Ordering FIB updates with oFIQUIC

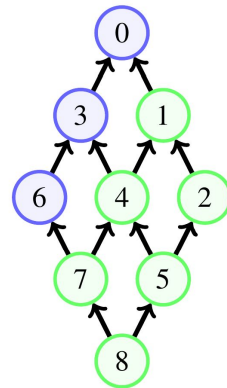
oFIQUIC design

1. Compute rspDAG
2. Derive an ordered list of (distant) nodes
3. Initiate QUIC tunnels
4. Send LSA to first node
5. Wait for distant node confirmation
6. Repeat steps 4 - 5 until list is empty

scheduling node



0 metric ∞
2 metric 1
4 metric 1



$rspDAG_{1 \rightarrow 0}(0)$

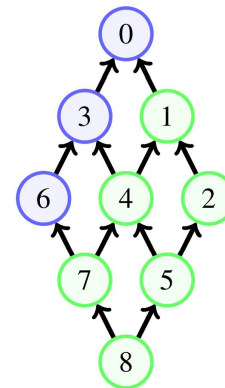
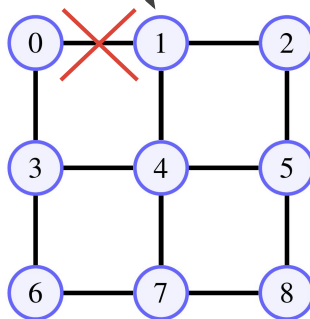
Ordered list: [R8, R5, R7, R2, R4]

Ordering FIB updates with oFIQUIC

oFIQUIC design

1. Compute rspDAG
2. Derive an ordered list of (distant) nodes
3. Initiate QUIC tunnels
4. Send LSA to first node
5. Wait for distant node confirmation
6. Repeat steps 4 - 5 until list is empty
7. Final flooding

scheduling node

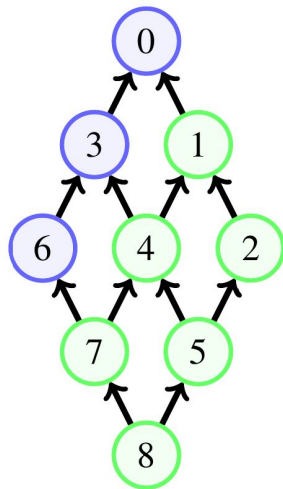
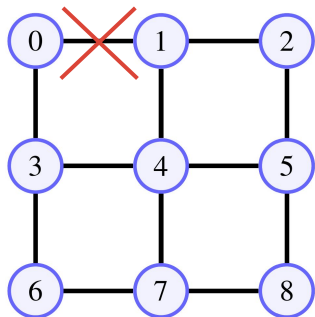


$rspDAG_{1 \rightarrow 0}(0)$

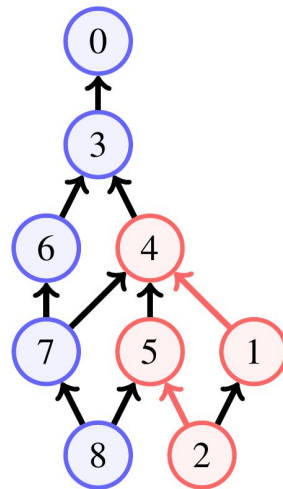
Ordered list: [R8, R5, R7, R2, R4]

Partial order vs Total order

oFIQUIC design



pre-convergence



post-convergence

Total order list: [R8, R5, R7, R2, R4]

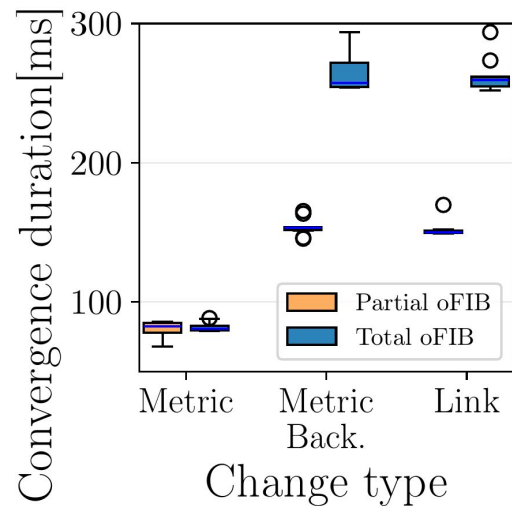
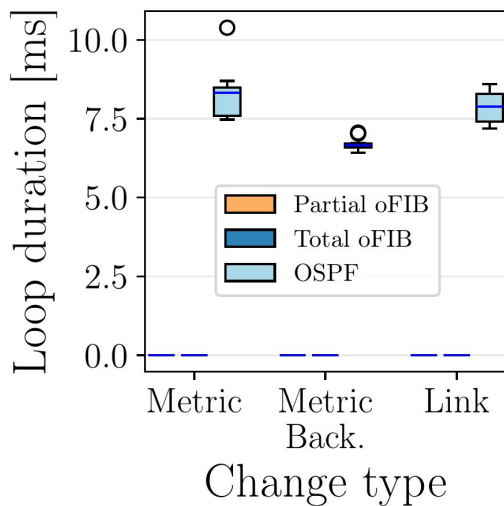
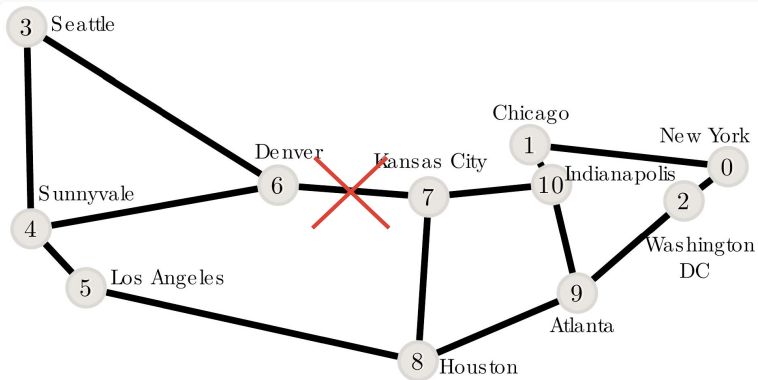
Partial order list: [R5, R2, R4]

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- oFIB: Ordered FIB updates
- oFIQUIC design
- **oFIQUIC performances**

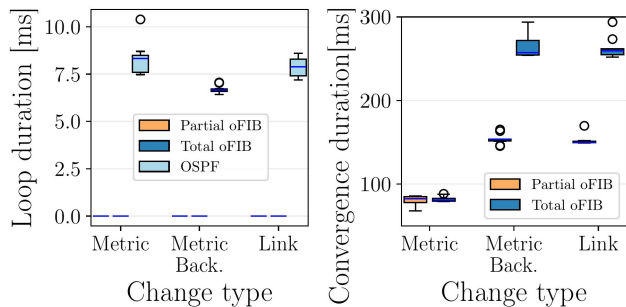
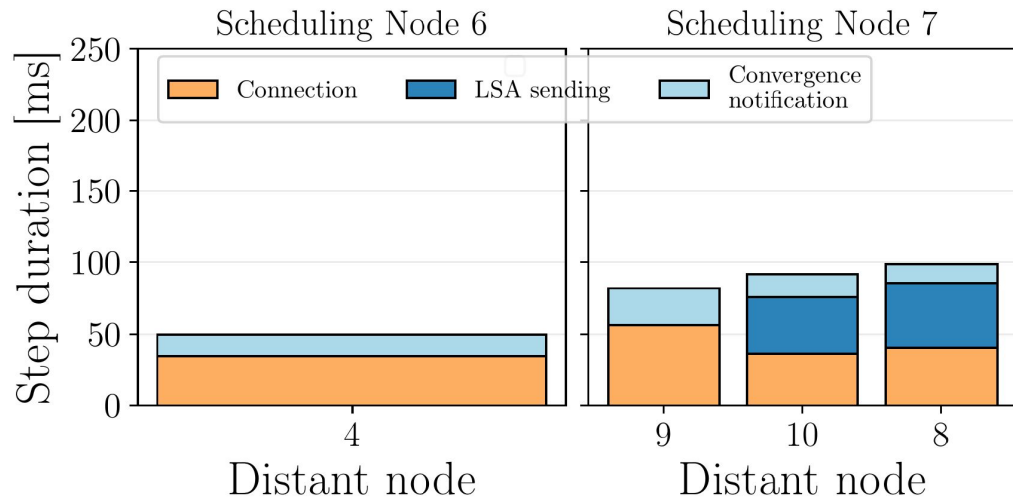
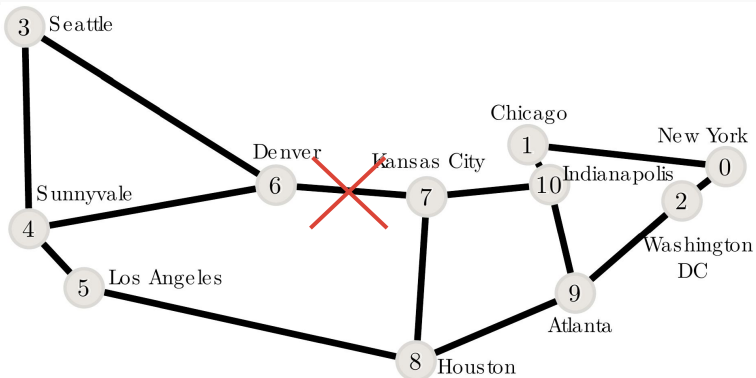
oFIQUIC convergence duration in Abilene

oFIQUIC performances



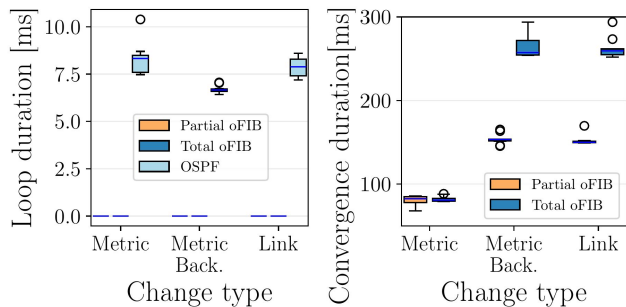
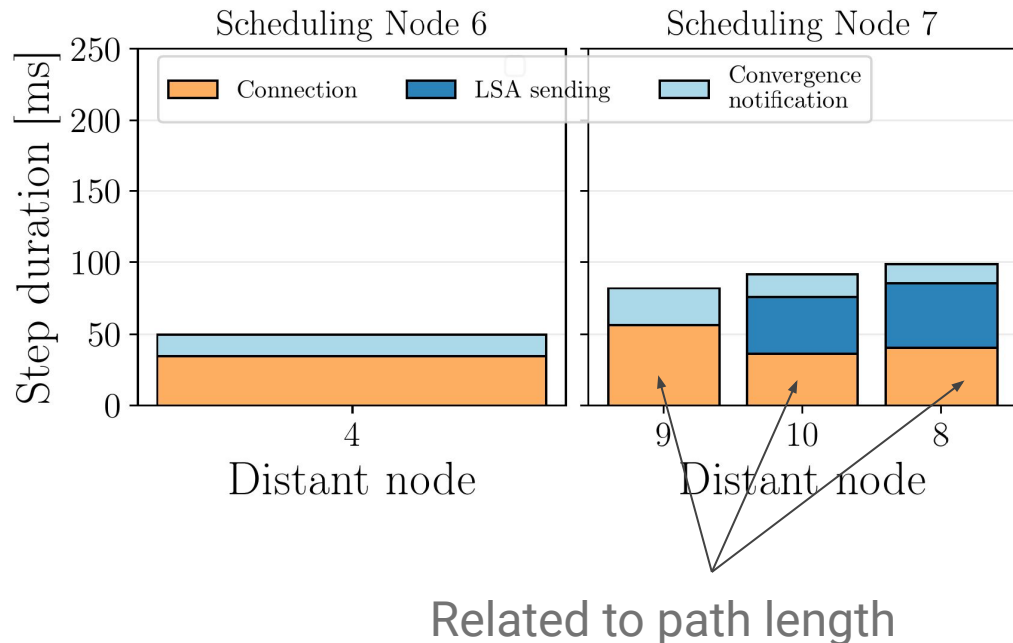
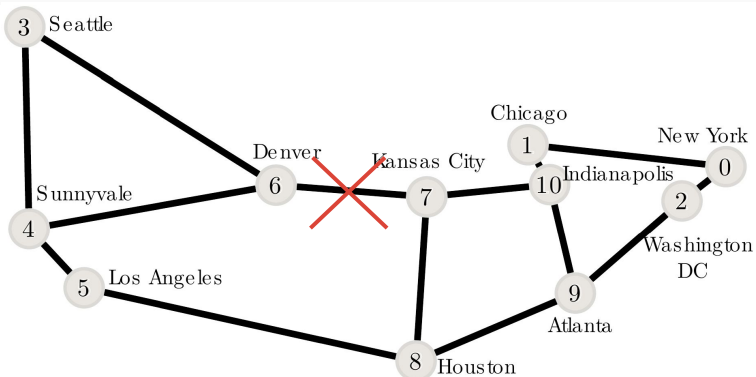
oFIQUIC contributions in Abilene

oFIQUIC performances



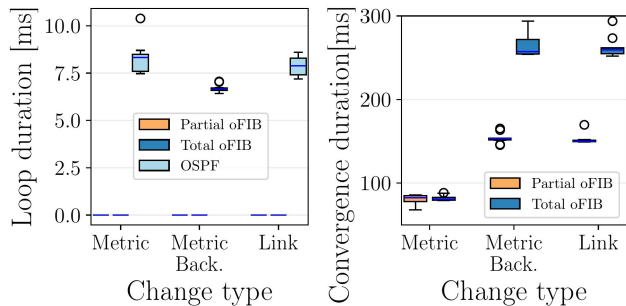
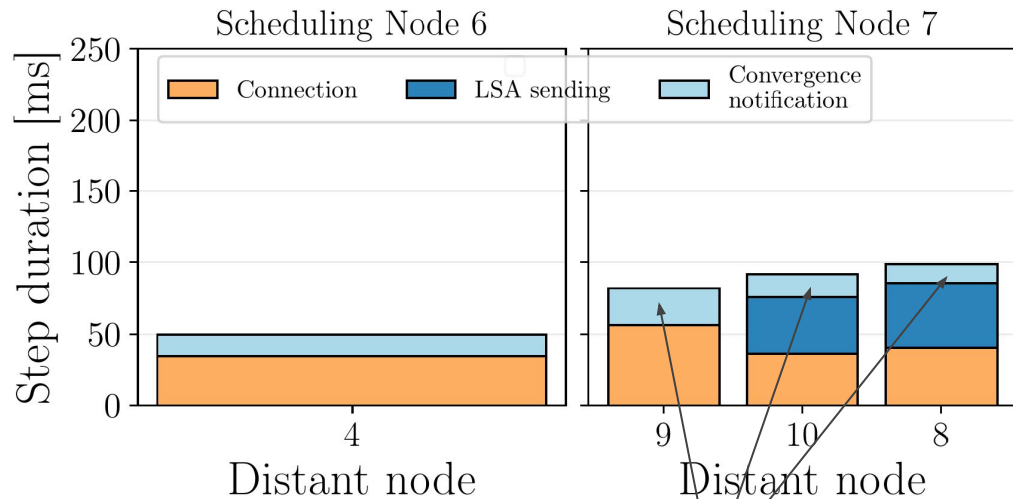
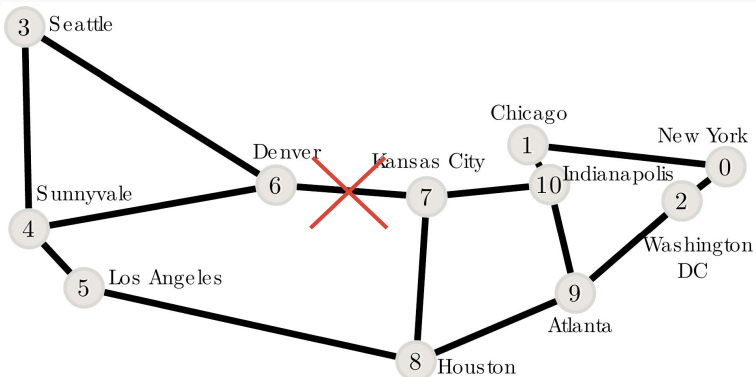
oFIQUIC contributions in Abilene

oFIQUIC performances



oFIQUIC contributions in Abilene

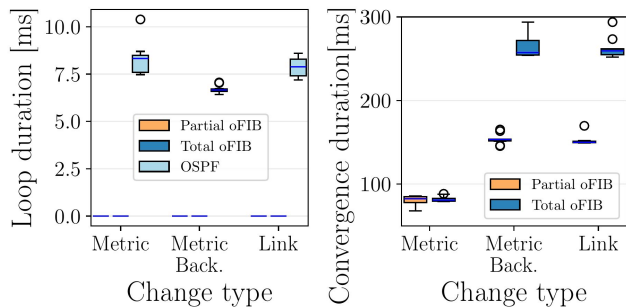
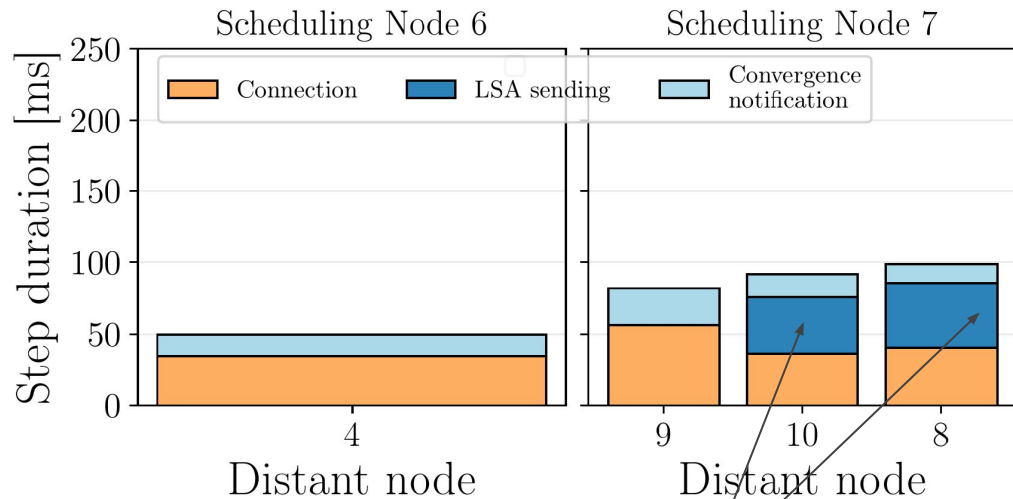
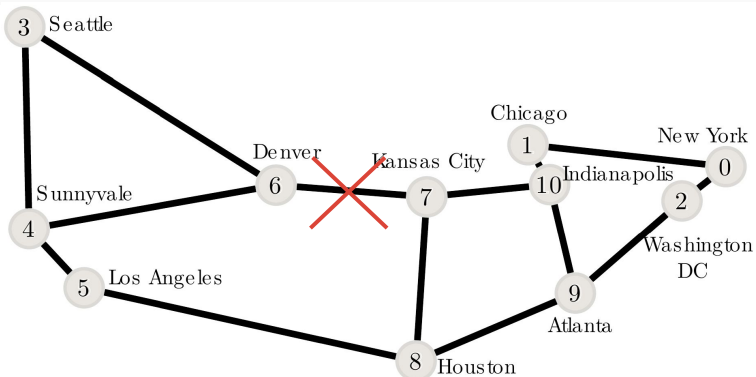
oFIQUIC performances



Related to path length

oFIQUIC contributions in Abilene

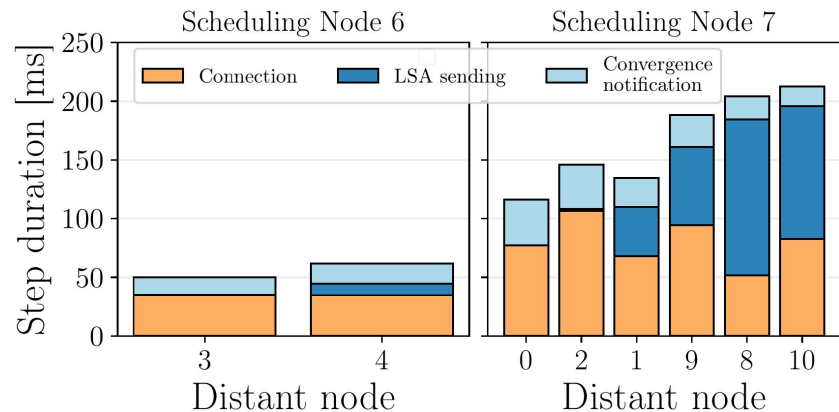
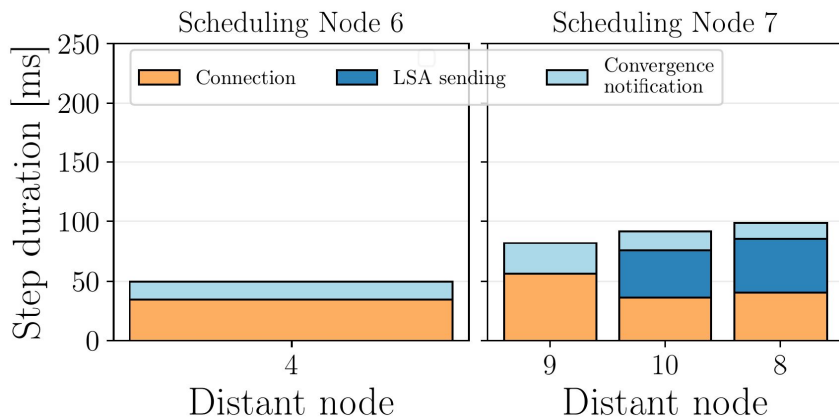
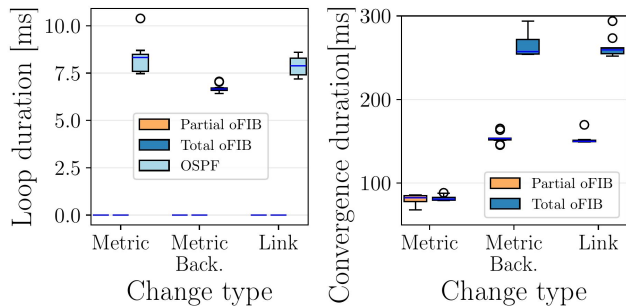
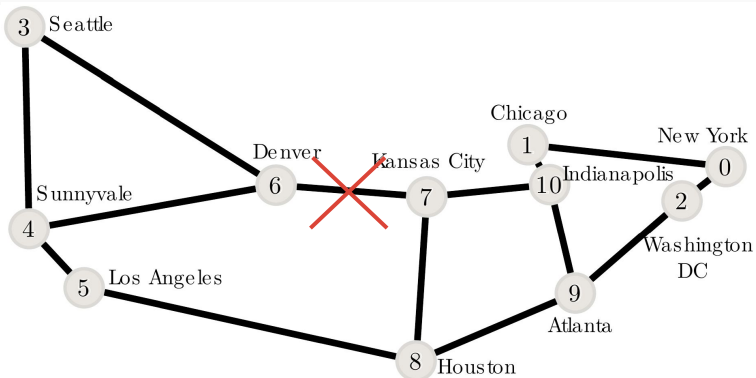
oFIQUIC performances



Related to previous step

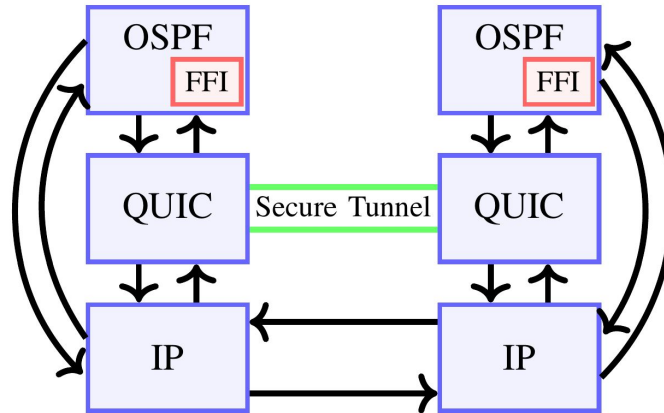
oFIQUIC contributions in Abilene

oFIQUIC performances



Summary

paper



artefacts

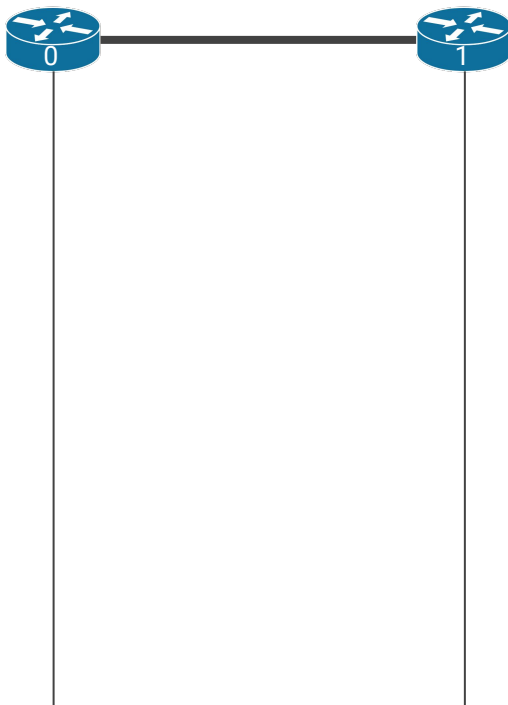
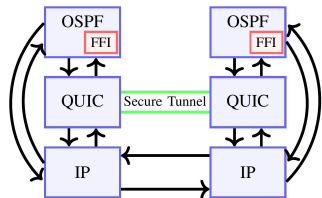


nicolas.rybowski@uclouvain.be

Backup slides

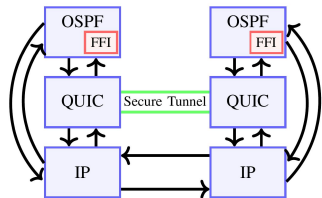
Establishing QUIC tunnels with direct IGP peers

oFIQUIC design

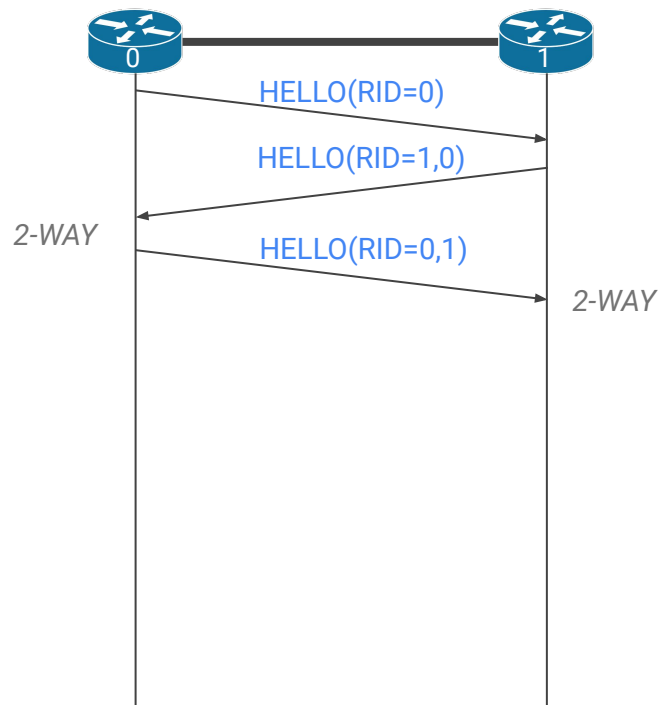


Establishing QUIC tunnels with direct IGP peers

oFIQUIC design

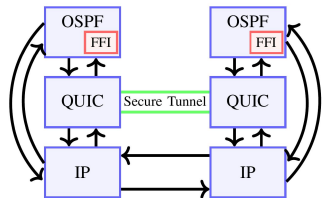


OSPF packet

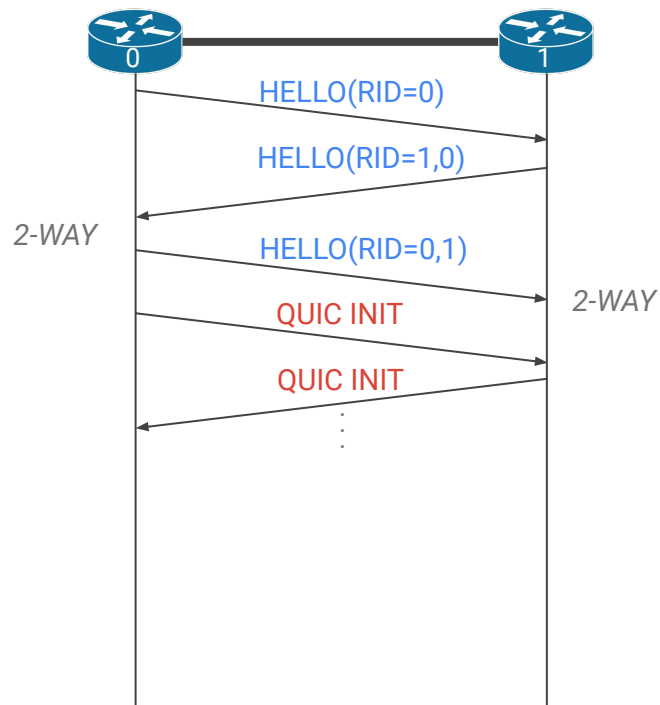


Establishing QUIC tunnels with direct IGP peers

oFIQUIC design

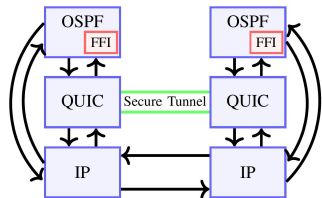


OSPF packet
QUIC packet

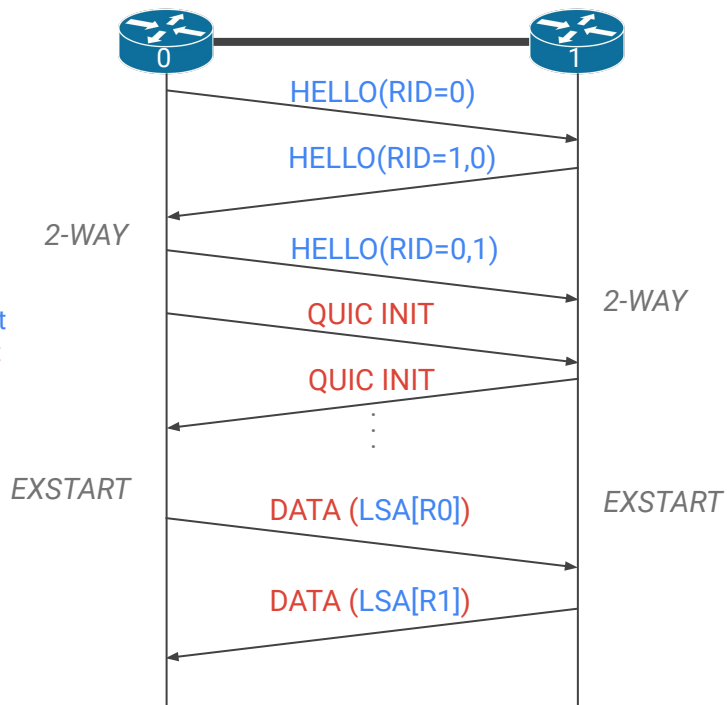


Establishing QUIC tunnels with direct IGP peers

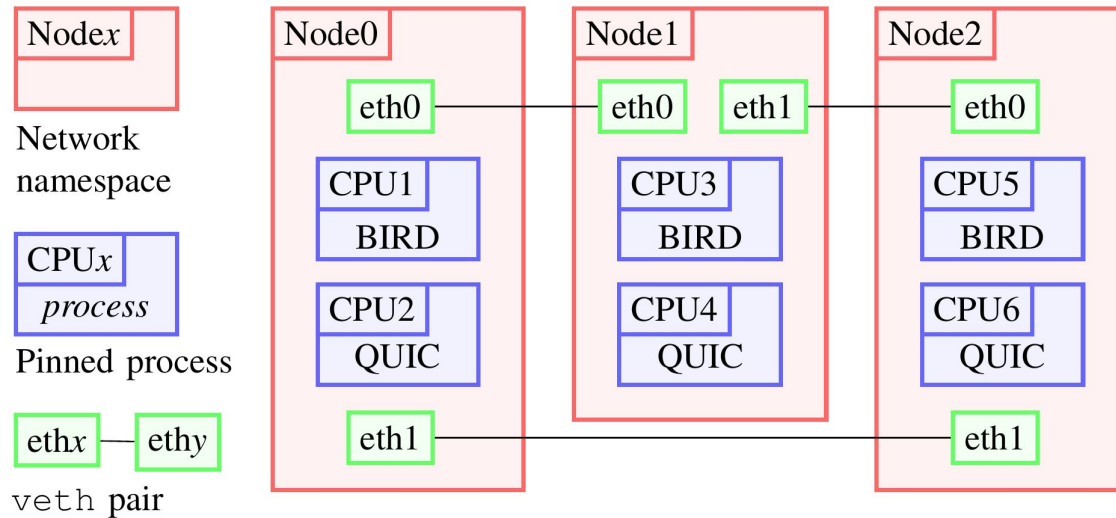
oFIQUIC design



OSPF packet
QUIC packet

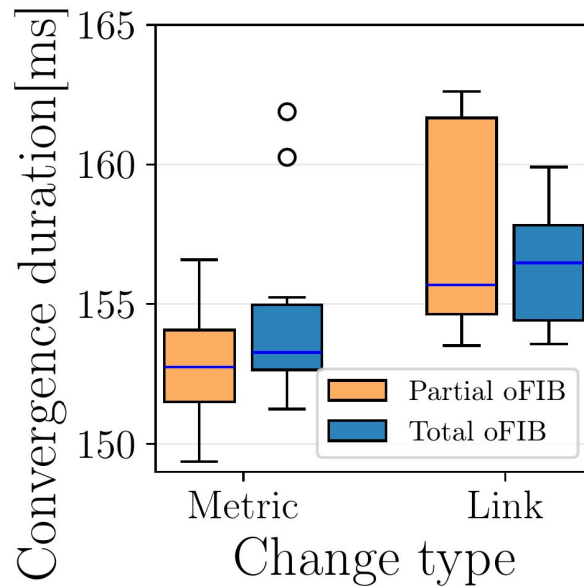
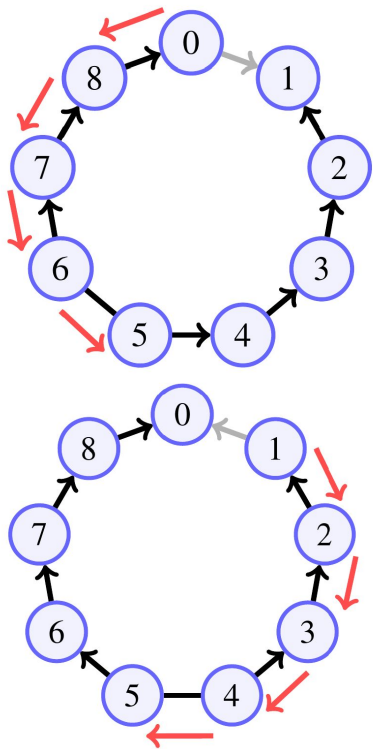


Distributed Micro Network Emulation (DUNE) framework



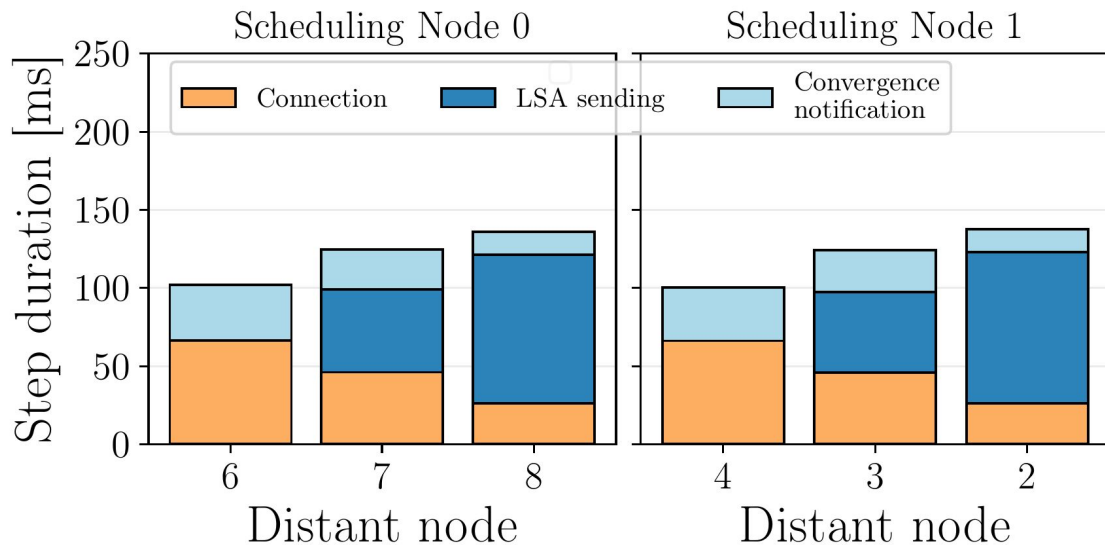
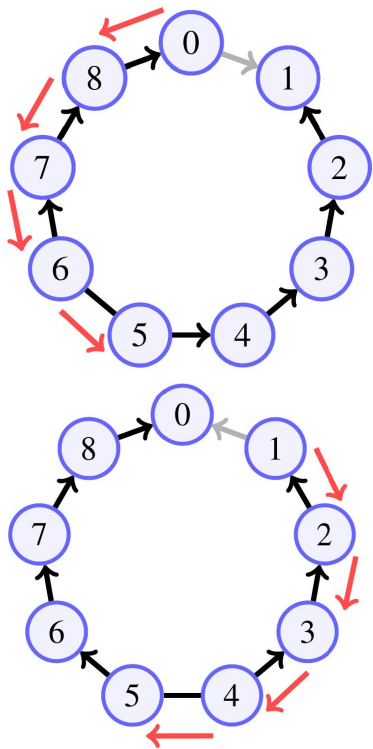
oFIQUIC convergence duration in a ring

oFIQUIC performances



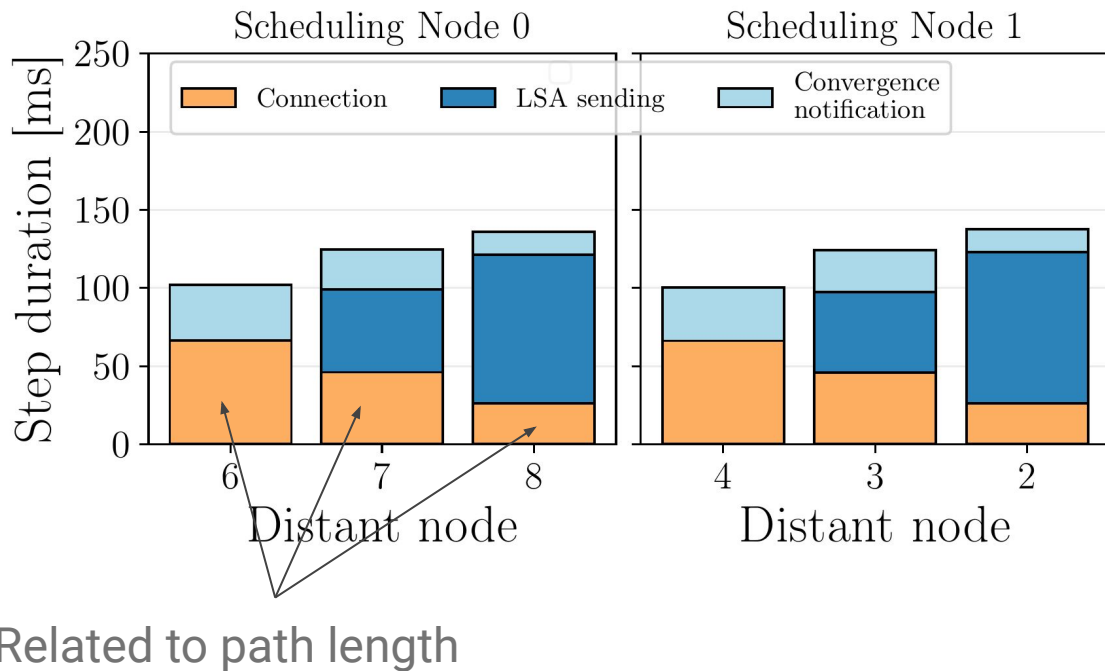
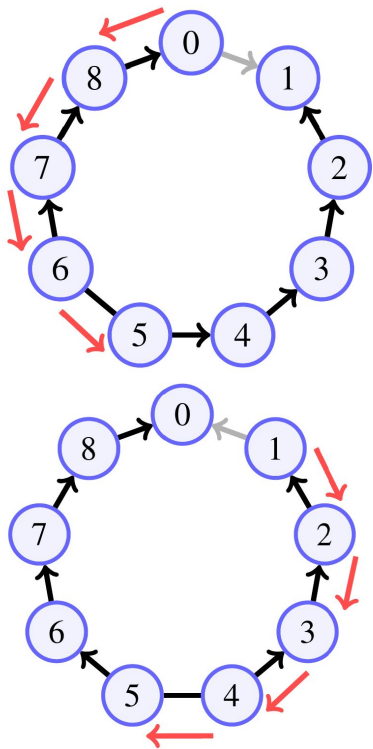
oFIQUIC contributions in a ring

oFIQUIC performances



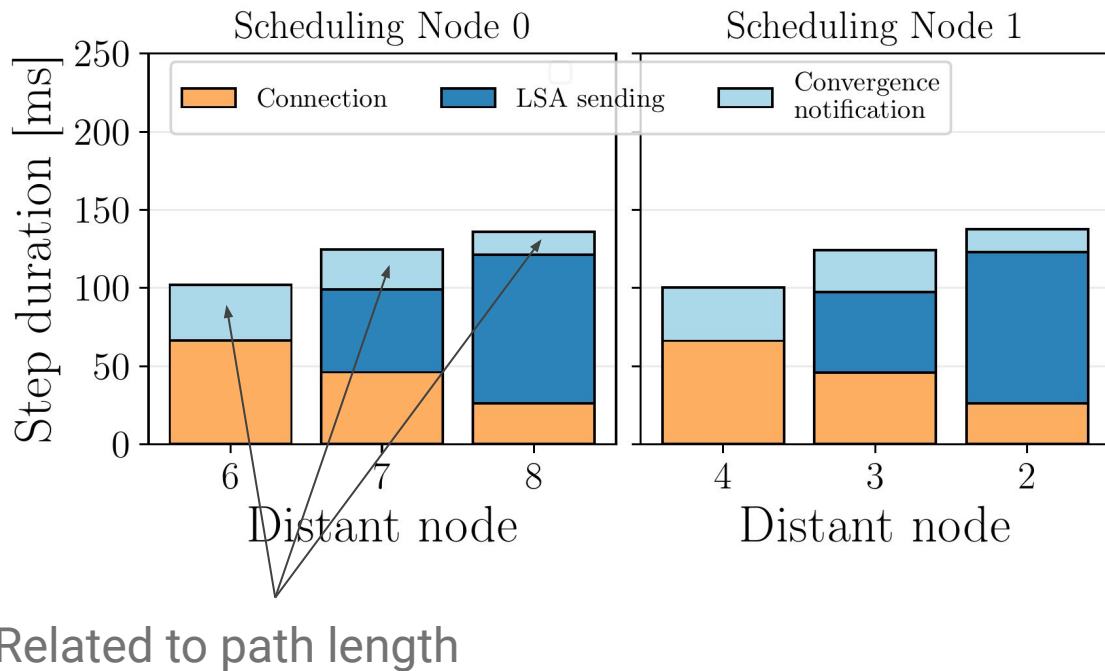
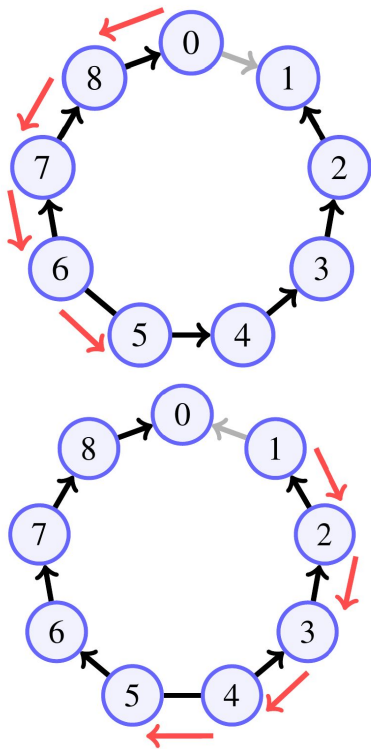
oFIQUIC contributions in a ring

oFIQUIC performances



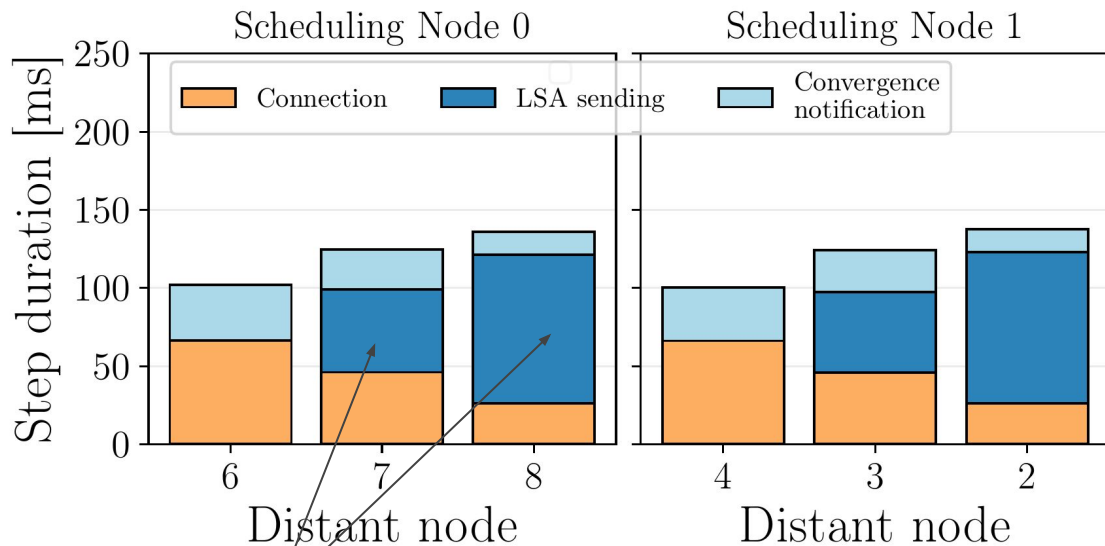
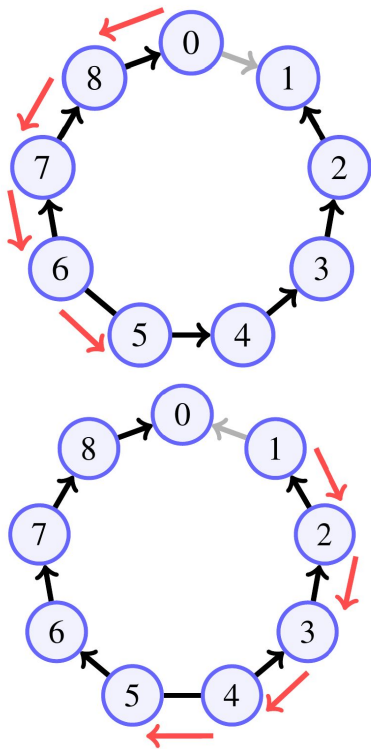
oFIQUIC contributions in a ring

oFIQUIC performances



oFIQUIC contributions in a ring

oFIQUIC performances



Related to previous step