Interdomain routing with BGP4
Part 1/5

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General references on BGP include :
Y. Rekhter and T. Li and S. Hares, , A Border Gateway Protocol 4 (BGP-4)},
Internet draft, draft-ietf-idr-bgp4-20.txt, work in progress,2003

A List of the internet drafts related to BGP may be found in
E. Chen, Y. Rekther, List of the Current BGP Documents, Internet draft, draft-chen-bgp-reference-03.txt, work in progress, 2002
See also : http://www.ietf.org/html.charters/idr-charter.html

A more readable textbook description of BGP may be found in
J. Stewart,BGP4 : interdomain routing in the Internet, Addison Wesley,1999
Vendor oriented books include :
• I. van Beijnum, BGP : Building Reliable Networks with the Border Gateway Protocol, O'Reilly, 2002
•S. Halabi, D. McPherson, Internet routing architectures, 2nd Edition, Cisco Press, 2001
D. McPherson, S. Srihari, R. White, Practical BGP, Addison Wesley, 2004

Tim Griffin maintains a very long and up to date list of references on BGP , see http://www.cambridge.intel-research.net/~griffin/interdomain/

Outline

- Organization of the global Internet
  - → Example of domains
  - Intradomain routing

- BGP basics

- BGP in large networks

- Interdomain traffic engineering with BGP

- BGP-based Virtual Private Networks

How to route IP packets in the global Internet ?

A map of the global Internet in 2000 ...

Source http://research.lumeta.com/ches/map/gallery/index.html
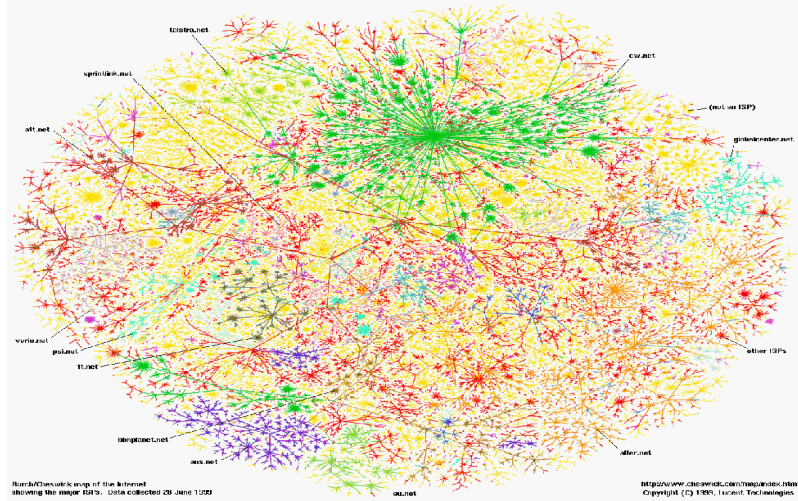
Organization of the Internet

- Internet is composed of more than 10.000
  autonomous routing domains

  - A domain is a set of routers, links, hosts and local
    area networks under the same administrative control
    - A domain can be very large...
      - AS568: SUMNET-AS DISO-UNRRA contains 73154560 IP
        addresses
    - A domain can be very small...
      - AS2111: IST-ATRIUM TE Experiment a single PC running Linux...

  - Domains are interconnected in various ways
    - The interconnection of all domains should in theory allow
      packets to be sent anywhere
    - Usually a packet will need to cross a few ASes to reach its
      destination

Types of domains

- Transit domain
 - A transit domain allows external domains to use its own infrastructure to send packets to other domains
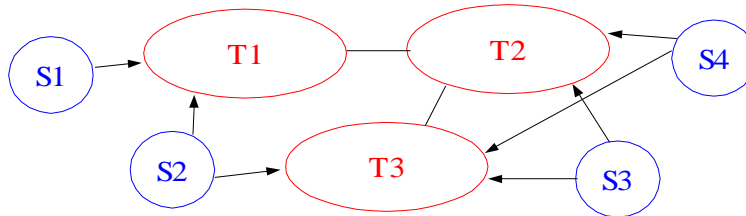
- Examples
  - UUNet, OpenTransit, GEANT, Internet2, RENATER, EQUANT, BT, Telia, Level3,...

Types of domains (2)

- Stub domain
  - A stub domain does not allow external domains to use its infrastructure to send packets to other domains
    - A stub is connected to at least one transit domain
      - Single-homed stub : connected to one transit domain
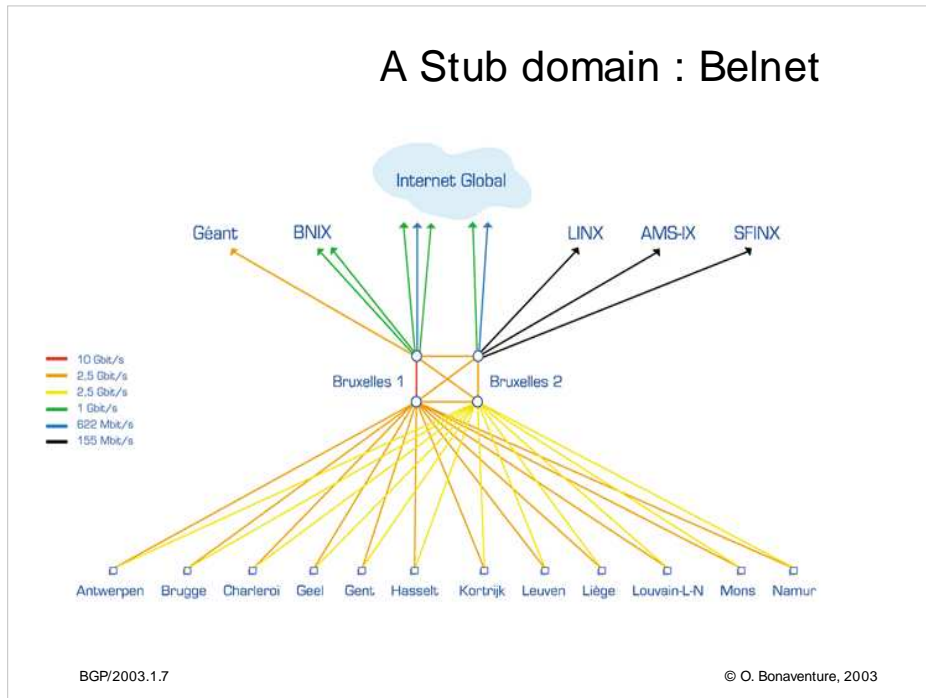   - Dual-homed stub : connected to two transit domains
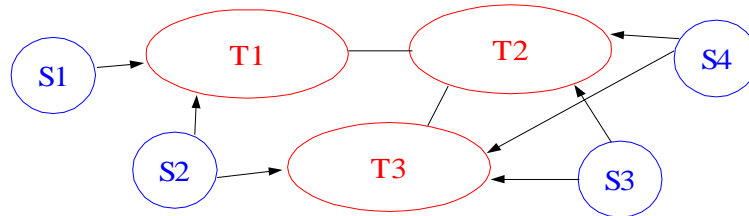
  - Content-rich stub domain
    - Large web servers : Yahoo, Google, MSN, TF1, BBC,...
  - Access-rich stub domain
    - ISPs providing Internet access via CATV, ADSL, ...
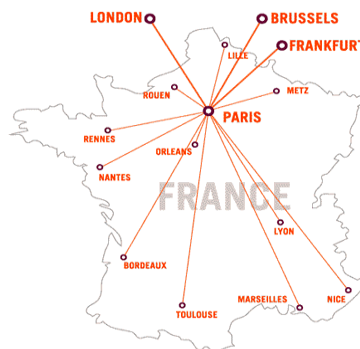
A Stub domain : Belnet

Source : http://www.belnet.be

Other maps of ISPs may be found at :
http://www.cs.washington.edu/research/networking/rocketfuel/interactive/

A transit domain : Easynet

http://www.easynet.be/home/index.cfm?id=15&l=1

A transit domain : GEANT

The Gigabit Research Network

Backbone topology
December 2001

■	10 Gbit/s
■	2.5 Gbit/s
■	1.2 Gbit/s
■	622 Mbit/s
■	310 Mbit/s
■	155 Mbit/s
■	45 Mbit/s
■	34 Mbit/s
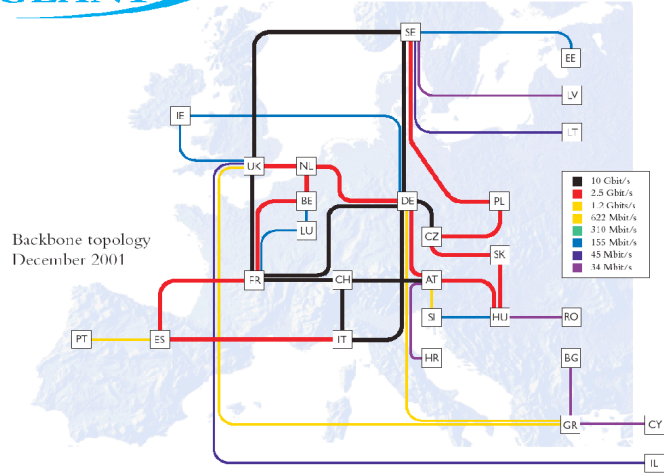
BGP/2003.1.9

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Source http://www.dante.net

A transit domain : BT/IGnite

To Washington

To New York

Stockholm

Dublin

London

Amsterdam

Madley

Frankfurt

Goonhilly

Brussels

Paris

Vienna

Zurich

Budapest

Milan

Madrid

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Source : http://www.ignite.net/info/maps.shtml

A large transit domain : UUNet

— 64 Kbps	— OC12c/STM4 (622 Mbps)
— T1/E1 (1.5 Mbps/2 Mbps)	— OC48c/STM16 (2.5 Gbps)
— E3/T3/DS3 (35 Mbps/45 Mbps)	— OC192c/STM64 (10 Gbps)
— T2 (6 Mbps)	• Single Hub City
— OC3c/STM1 (155 Mbps)	■ Multiple Hubs City
	▣ Data Center Hub
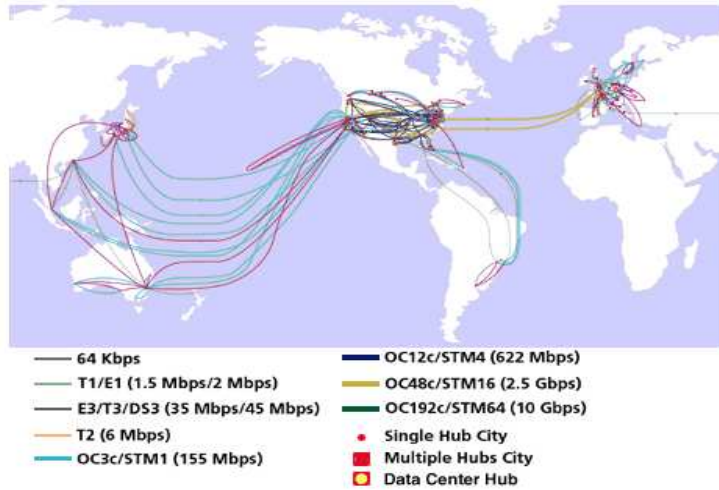
BGP/2003.1.11

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Source http://www.uu.net
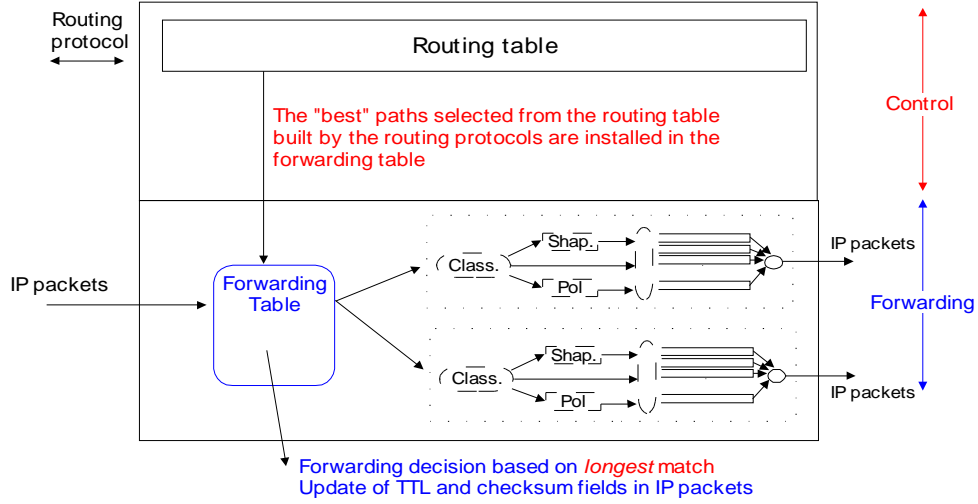
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Architecture of a normal IP router

Routing protocol

Routing table

The "best" paths selected from the routing table built by the routing protocols are installed in the forwarding table

Control

IP packets

Forwarding Table

Class. Shap. Pol

Class. Shap. Pol

IP packets

IP packets

Forwarding

Forwarding decision based on *longest* match
Update of TTL and checksum fields in IP packets

Internet routing

- **Interior Gateway Protocol (IGP)**
  - ◆ Routing of IP packets **inside each domain**
    - ◆ Only knows topology of its domain

Domain2

Domain4
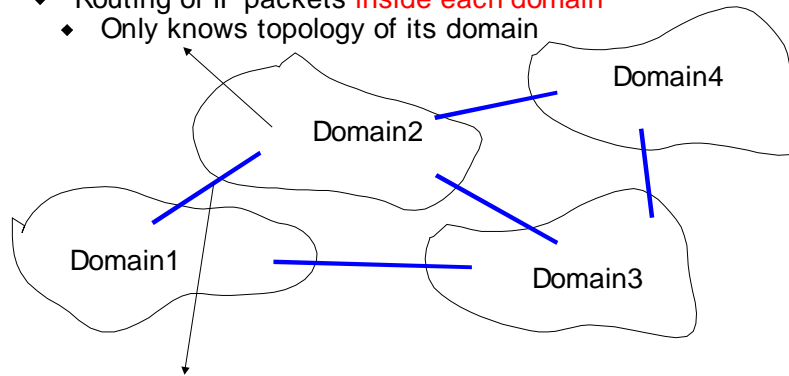
Domain1

Domain3

- **Exterior Gateway Protocol (EGP)**
  - ◆ Routing of IP packets **between domains**
    - ◆ Each domain is considered as a blackbox

Intradomain routing

- Goal
  - Allow routers to transmit IP packets along the best path towards their destination
    - best usually means the shortest path
      - Shortest measured in seconds or as number of hops
    - sometimes best means the less loaded path
  - Allow to find alternate routes in case of failures

- Behavior
  - All routers exchange routing information
    - Each domain router can obtain routing information for the whole domain
    - The network operator or the routing protocol selects the cost of each link

Three types of
Interior Gateway Protocols

- Static routing
  - Only useful in very small domains

- Distance vector routing
  - Routing Information Protocol (RIP)
    - Still widely used in small domains despite its limitations

- Link-state routing
  - Open Shortest Path First (OSPF)
    - Widely used in enterprise networks

  - Intermediate System- Intermediate-System (IS-IS)
   - Widely used by ISPs

Distance vector routing
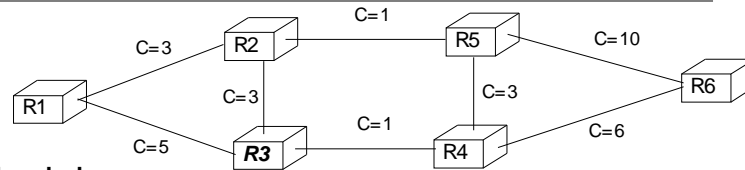
- Principle
  - Router configuration
    - Cost associated with each link
  - Each router sends <u>periodically</u> a distance vector containing, for each known prefix, :
    1. The IP prefix
    2. The distance between itself and the destination
       - The distance vector is a summary of the router's routing table
  - Each router receives its neighbor's distance vectors and builds its routing table based on those vectors

Issues with distance vector routing

- How to deal with link failures ?
  - Routers should send their distance vector when they detect the failure of one of their links

- How to avoid the count-to-infinity problem ?
  - Utilize a non-redundant star shaped network

  - Limit the maximum distance between routers
    - For RIP, $\infty=16$!

  - Split horizon
    - Router A does not advertise to router B the routes for which it sends packets via router B
  - Split horizon with Poison reverse

Link state routing
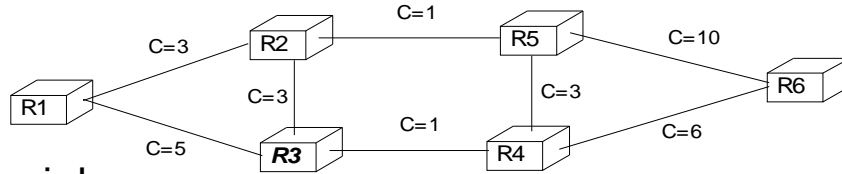
- **Principle**
  - Each router builds link state packet containing its local topology
    - ◆ Link state packets are created at regular intervals and when the local topology changes
  - Link state packet is reliably flooded to all routers inside the domain
  - Each router knows the complete domain topology
  - Computes routing tables by using Dijkstra
    - ◆ The best path is the path with the smallest cost

For a description of OSPF, see J. Moy, OSPF : anatomy of an Internet routing protocol, Addison-Wesley, 1998

ISIS is defined in

R. Callon, Use of OSI IS-IS for Routing in TCP/IP and Dual Environments, RFC1195, Dec. 1990

IP forwarding

- Usually
  - Forwarding table contains, for each prefix
    - The prefix
    - The best path (outgoing interface) to reach this prefix

- Sometimes
  - Forwarding table contains, for each prefix
    - The prefix
    - N equal cost paths to reach this prefix
      - A first path (outgoing interface) to reach this prefix
      - A second path (outgoing interface) to reach this prefix
      - A third path (outgoing interface) to reach this prefix
      - ...
  - A load balancing mechanism is used to send the IP packets over the N available paths

Load balancing algorithms

- Simple solution
  - Round-Robin or variants to dispatch packets on a per packet basis

  - Advantages
    - easy to implement since number of paths is small
    - traffic will be divided over the equal cost paths on a per packet basis
      - each path will carry the same amount of traffic

  - Drawbacks
    - two packets from the same TCP connection may be sent on different paths and thus be reordered
      - TCP performance can be affected by reordering

References to load balancing algorithms include :

C. Hopps, Analysis of an Equal Cost MultiPath algorithm, RFC2992, Nov. 2000

Z. Caro, Z. Wang, E. Zegura, Performance of Hashing-Based Schemes for Internet Load Balancing, INFOCOM2000,
http://www.ieee-infocom.org/2000/papers/650.ps

Load balancing algorithms (2)

- How to perform load balancing without maintaining state for each TCP connection ?
  - Principle
    - concatenate IP src, IP dest, IP protocol, Src port, and Dest port from the IP packet inside a bit string
      - bitstring = [IP src:IP dest:IP protocol:Src port:Dest port]
    - compute path = Hash(bitstring) mod P
      - hash function should be easy to implement and should produce very different numbers for close bitstring values
      - candidate hash functions are CRC, checksum, ...
  - Advantages
    - all packets from TCP connection sent on same path
    - traffic to a server will be divided over the links
  - Drawback
    - does not work well if a few TCP connections carry a large fraction of the total traffic

Summary

- Types of domains
  - Transit domain
  - Stub domain

- Intradomain routing
  - Selects the best route towards each destination based on one metric
    - Static routing
    - Distance vector routing
    - Link-state routing
  - Load balancing methods allow to place several paths in the forwarding table