Tuning Multipath TCP for Interactive Applications on Smartphones

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May 16, 2018

Université catholique de Louvain IFIP Networking 2018





Smartphones are Multi-Homed Devices





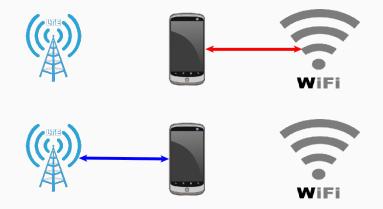


What is wrong today with TCP?

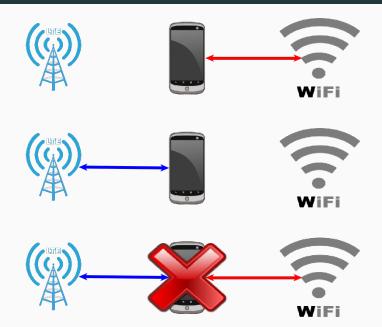




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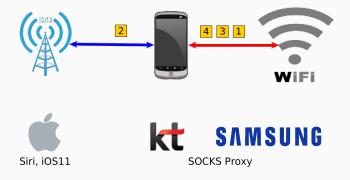
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Imperfections with Multipath TCP on Smartphones

MultiMob

Evaluation with Real Users

Conclusion

Imperfections with Multipath TCP on Smartphones

Smartphone Traffic \neq Bulk Transfers

Linux implementation of Multipath TCP

- Mainly bandwidth aggregation driven
- But most of the connections are (very) short in bytes
 - Most are latency-sensitive
 - Growing importance of interactive applications
 - Apple Siri, Google Now, Amazon Alexa,...
 - Multipath = network resiliency under mobility



















Smartphone = data consumers

- Servers take most of the scheduling decisions
 - Prefer lowest perceived latency path by default
- Ability to define cellular as **backup** path
 - Use non-backup paths unless they all failed

WiFi & cellular path creations as soon as possible

- make-before-break for fast network handover
- But most cellular subflows do not see any data
 - WiFi is often sufficient
 - Short connections, low application push rate, no mobility...
 - Network + energy waste
 - Cellular can consume as much as a screen 100% on

MultiMob

Imperfections with Multipath TCP on Smartphones

MultiMob

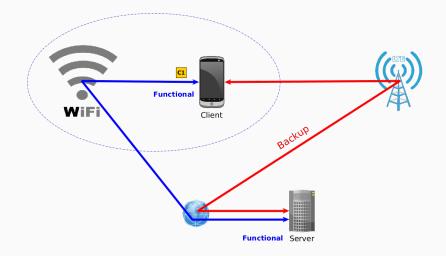
Giving Control to the Client Break-Before-Make Immediate Reinjections

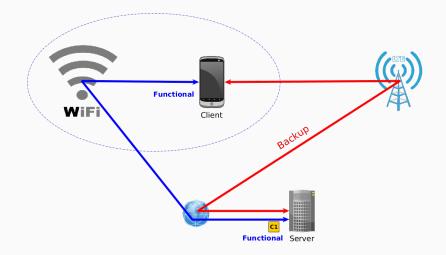
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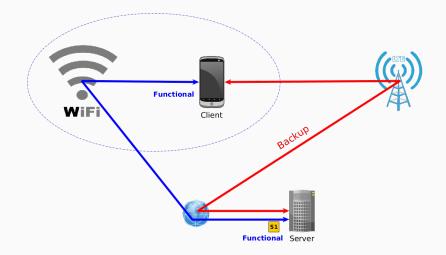
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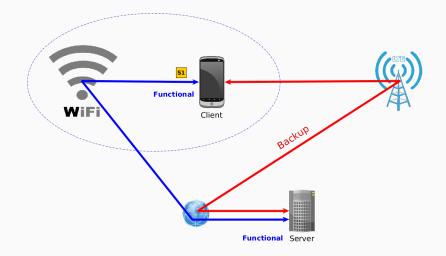
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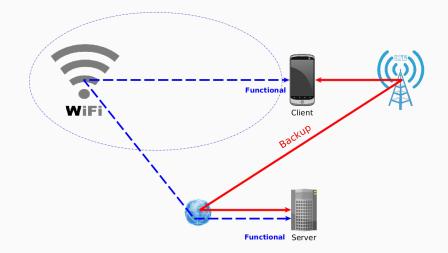
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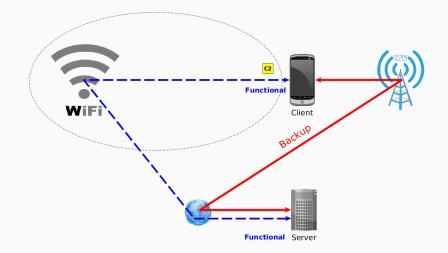


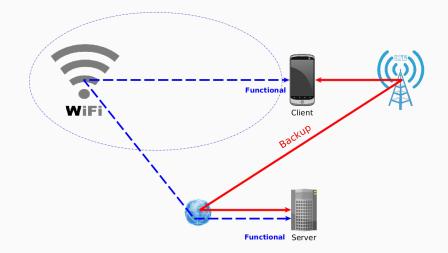


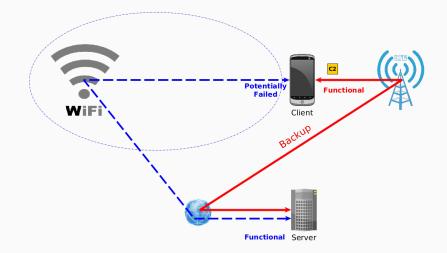


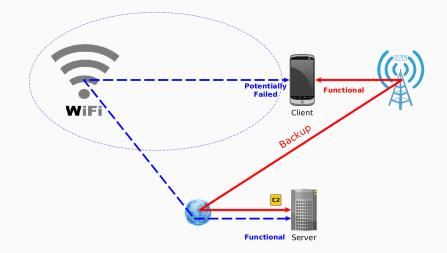


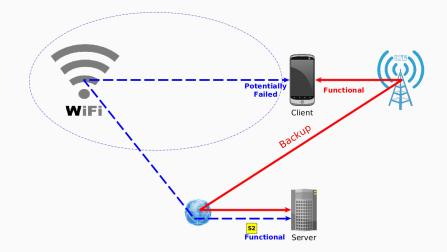


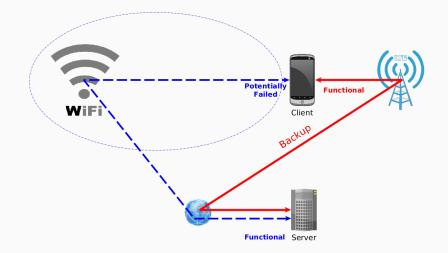


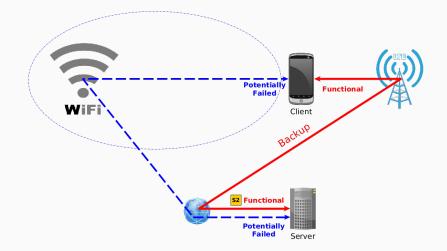


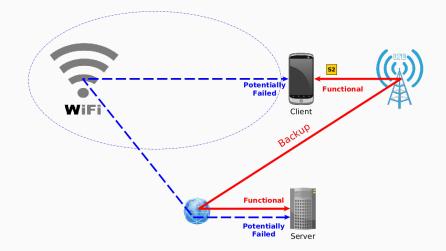






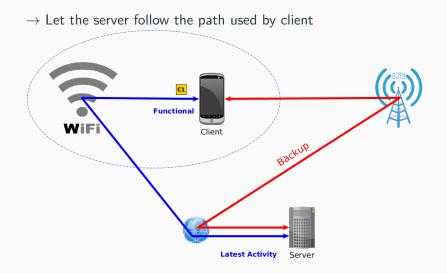


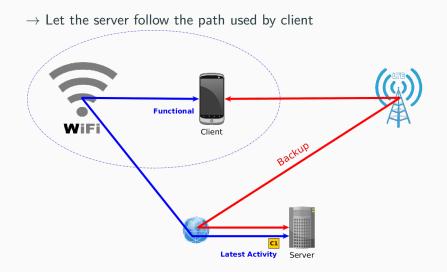


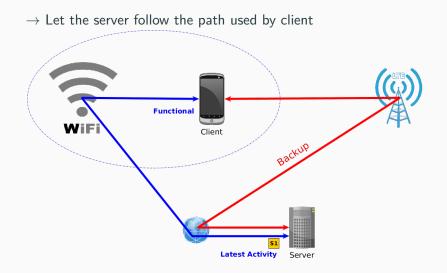


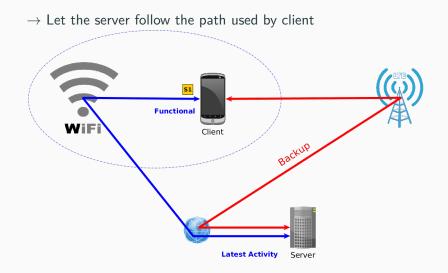
 \rightarrow Let the server follow the path used by client

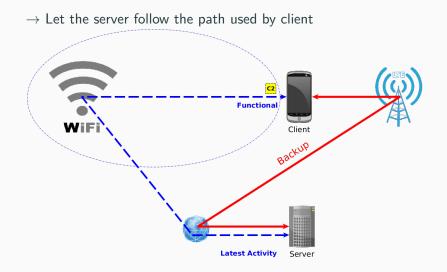
MultiMob Server Scheduler

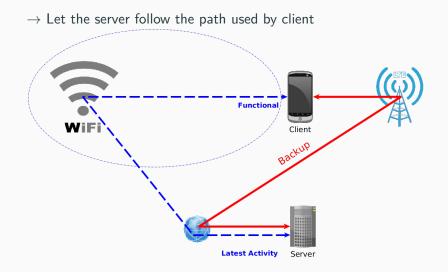


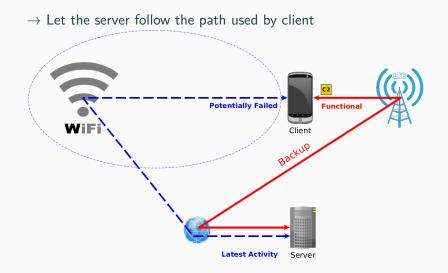


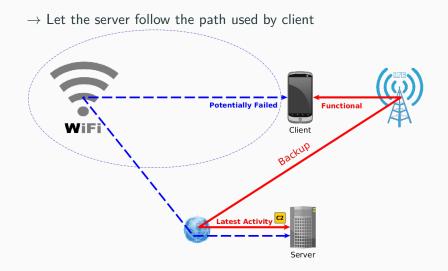


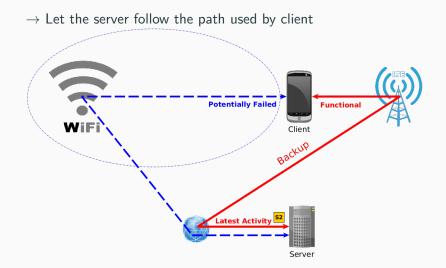


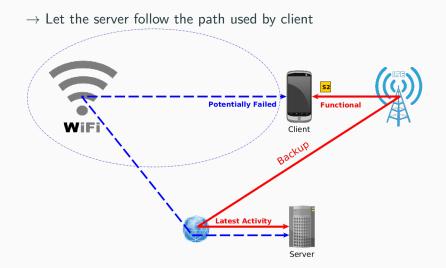












MultiMob

Break-Before-Make









- \rightarrow Minimizes amount of data sent over cellular
 - Yet lot of energy wasted by LTE usage...











- \rightarrow No LTE usage if cellular path not needed!
 - But need to quickly detect bad wireless networks

Quick Failure Detection





Losses / Retransmissions

Quick Failure Detection

?

Losses / Retransmissions

Proposing an in-kernel Multipath TCP oracle

- Periodically compute statistics about *netpaths*
 - Netpath = (*IP*_{src}, *IP*_{dst}, Net interf.)
- Trigger backup creation upon excessive losses/retransmissions
 - \bullet + prevents primary use when bad

IP _{src}	IP _{dst}	Net interf.	TCP sfs	TCP stats
1.2.3.4	4.5.6.7	WiFi	[tp ₁ , tp ₃]	sloss 2%,
2.3.4.5	5.6.7.8	Cellular	[tp ₂]	sloss 0%,
2.3.4.5	4.5.6.7	Cellular	[tp ₄ , tp ₅]	sloss 15%,

 Table 1: Oracle monitoring table example.

MultiMob

Immediate Reinjections

The backup cellular path creation is delayed

- Nice from a energy consumption point of view...
- ...but incurs larger app perceived latency in mobility cases

Furthermore, additional Multipath TCP path creation takes time...

Improving Establishment of Additional Subflows

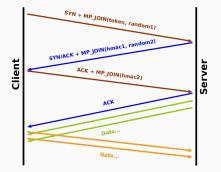


Figure 1: Normal JOIN.

Improving Establishment of Additional Subflows

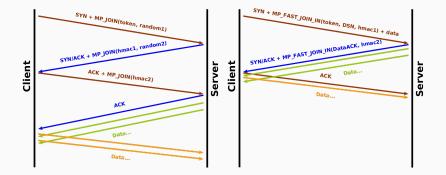


Figure 1: Normal JOIN.

Figure 2: Fast JOIN with data.

Evaluation with Real Users

Half-dozen of smartphone users

- Running Android 6.0.1
- Two sets of users
 - Running vanilla Multipath TCP (with/without backup)
 - Running MultiMob

Performing interactive traffic measurements

- Light bursty request/response traffic, see paper for details
 - Observe delay between request sent and response received
- Running 80 s once motion is detected
- Only consider tests with both WiFi/LTE online at beginning

Interactive Traffic - Maximal Request Delay over Tests

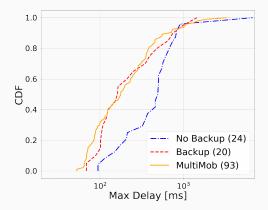


Figure 3: Maximal delays.

- MultiMob does not perform worse than Vanilla Multipath TCP
- Using multiple paths is not always beneficial

Interactive Traffic - Energy Consumption ("static" tests)

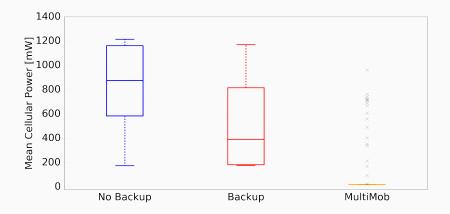


Figure 4: Estimated mean cellular power, WiFi not lost.

• MultiMob consumes much less cellular energy!

Conclusion

- Vanilla Multipath TCP not efficient from energy viewpoint
- MultiMob: Multipath TCP tuned for smartphones
 - Keep similar performances
 - Lower energy consumption
 - Lower LTE radio resource usage
- MultiMob is available

http://multipath-tcp.org/multimob

Thanks for your attention! Feel free to ask questions!

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- Client creates backup paths once bad network detected
- One connection experiences issues, others react

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- \rightarrow What about single bulk download without background traffic?
 - By default, server does not create paths
 - NAT, firewalls,...

Exchange additional information at connection level

- Send *RTO*_{sender} to peer
- MP_IDLE bit in DSS to indicate no more data to send (now)

Allows to setup a receive timer

- Reset to *RTO_{sender}* if activity without MP_IDLE
- Stop if received data with MP_IDLE

If receive timer fires, create backup path

Assessing Receive Timer

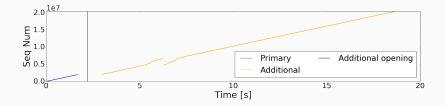


Figure 5: Time-sequence graph of server to client flow from client side perspective for a HTTP GET of 20 MB. Primary has 100% losses at 1.5 s.

- Client opens additional path after inactivity period
- Once established and data detected as lost, server continues the connection on the additional path